## 8



Been looking for a practical way to input and dump programs to your computer? Well your search is over.

With our new AC-30 Cassette Interface you will be able to store and input program data to any computer system having RS232 serial interfaces and a UART circuit having an accessible 16X clock frequency. Data format is the "Kansas City" standard which was selected for its tolerance of speed variations in the recording device. The AC-30 may be used with any cassette recorder of reasonable quality.

If both your computer and terminal have accessible 16X UART clocks and will operate at 300 baud-as do our 6800 computer and CT-1024 terminal system-the AC-30 may be used between the terminal's serial interface and the computers control interface. This eliminates the need for a separate interface to drive the cassette unit. It also allows you to use the computer system's tape load and dump routines built into Mikbug ${ }^{\circledR}$ or similar ROM software.

Independent control circuits are provided
for two audio cassette recorders (not included in the kit). One recorder's tape may be read while the second is recording a new updated tape; making it possible to generate new program tapes, data tapes and to create program object tapes while reading and assembling program source tapes. The operating mode for each recorder is selected by switches on the front panel and LED indicators show the mode that is selected at any particular time. Computer controlled record, play and motor control commands may be used with this system if they are available from the terminal being used. This feature is available on our CT-1024 terminal if the CT-CA cursor control card is installed.

The AC-30 is housed in a $123 / 4^{\prime \prime} \times 3^{\prime \prime} \times 121 / 2^{\prime \prime}$ aluminum chassis. It is powered by a self contained $115 / 230$ Volt AC $50-60 \mathrm{~Hz}$ power supply. Data is FSK format using 1200 Hz and 2400 Hz at a 300 baud data rate. Recorder speed tolerance need be only $\pm 20 \%$.
© Trademark Motorola

AC-30 Cassette Interface Kit $\$ 79.50$ ppd

## How to use

 your computer with analog signals

Now you have a way to get analog information into and out of your microcomputer. It's an easy, fast, and unbelievably inexpensive way.

It's Cromemco's new $\mathrm{D}+7 \mathrm{~A}^{\mathrm{TM}}$ highperformance $1 / 0$ module which gives you:

- 7 channels of 8 -bit analog-to-digital conversion (to input analog data to the computer)
- 7 channels of digital-to-analog conversion (to output computer data in analog form)
- an 8-bit parallel I/O port to input and output data in digital form.
- a fast conversion time of 5 microseconds.


## A MULTITUDE OF USES

The $\mathrm{D}+7 \mathrm{~A}$ makes it easy to use your computer for the jobs you want it to do-such as process control, digital filtering; games, oscilloscope graphics, speech recognition, speech and music synthesis.

The D+7A lets you input and output analog data with all sorts of devices: joysticks, ham radio gear, measurement instruments, machine tools, transducers, control systems, motors, recorders, and plotters, to name just a few.

## NO FURTHER SOFTWARE NEEDED

The D+7A I/O plugs directly into the Altair 8800 or IMSAI 8080 microcomputers. Analog signal range is from -2.56 to +2.56 volts ( 20 -millivolt increments) on both input and output sides.
Simple "Input" and "Output" instructions initiate A/D conversion and read in or out the ensuing 8 bits of data. No further software is required. During conversion the D+7A holds down the computer "Ready" line.

Addresses of the input and output ports are jumper-wire selectable in blocks of 8 . Sample-and-hold circuitry is used to "latch" the analog outputs.

## LOW-PRICED

The low price of the $D+7 A$ is a result
of Cromemco's design leadership. The D+7A and all Cromemco peripherals are of advanced computer-grade quality. The $\mathrm{D}+7 \mathrm{~A}$ is solder-masked and printed with full legend for easy, error-free assembly.

## AT COMPUTER STORES/MAIL

You can get the D+7A at computer stores in either kit or assembled form.

Or order directly by mail from Cromemco. Delivery is from stock to 30 days. The $D+7 \mathrm{~A}$ is certain to be popular so order now.
D+7ATM I/O kit . . \$145
D+7ATM $1 / 0$ assembled \$245

Each D+7A includes a connector to connect to the 8 input and 8 output ports. Shipped prepaid if fully paid with order. California users add $6 \%$ sales tax.
Mastercharge and BankAmericard accepted with signed order. Please include card number and expiration date.


Have you ever looked through the surplus catalogs and wondered whether those memory core planes and stacks advertised could be used for anything other than tea strainers? For theory and practical information on Coincident Current Ferrite Core Memories turn to James R Jones' article.

Bruce A Anderson describes his experiences Assembling a Sphere in his review of what rolled out of the production facilities in Bountiful UT last fall.

One of the most important questions people ask is "how do I learn about what a computer does?" One way to help out friends who are trying to get into the swing of things with programming is to implement a version of Charles Howerton's Educator-8080 program so that they can interactively Explore an 8080 with Educator- 8080 .

A thorough explanation of the instruction set should accompany any product intended for wide distribution. An example of such an explanation is provided by Nat Wadsworth's

Machine Language Programming for the " 8008 " and Similar Microcomputers, a manual which is sold by Scelbi Computer Consulting Inc. In this issue is the first of three direct reprints from that manual: Chapter 1 which describes the 8008 instruction set.

One of the problems of interfacing unknown electronics is figuring out how to accomplish the match. Ken Barbier built a character generator, went out and bought a TV set, then faced the problem of building a driver for the TV. The result was The "Ignorance is Bliss" Television Drive Circuit.

While not really promising the entire big blue sky, when you Put the "Do Everything" Chip in your Next Design you'll end up with a computer that has five separate programmable real time clocks, standard serial communications data rates from 110 baud to 9600 baud, automatic generation of an 8080's RST $n$ interrupt vectors, an 8 bit parallel output and an 8 bit parallel input port. Turn to Robert Baker's latest article to find out about this nifty chip.

What's it like to be isolated from bountiful US surplus markets? In a sense, it means a relative isolation from modern LSI products, as Dr Michael N Hayes reports on his experiences in Tokyo and Manila in December 1975. Read his report on Surplus Electronics in Tokyo and Manila in this issue.

There are many ways to wire a circuit. The most common manufacturing method is printed wiring. But you can also Make Your Own Printed Circuits at home, using techniques described by James Hogenson in his article.

One of the most interesting applications of computers is in the area of graphic outputs. Using a vector CRT or a plotter, drawing pictures of mathematically generated abstractions or simple cartoons can be the beginning of hours of fun. But A Plot Is Incomplete Without Characters so Richard J Lerseth concocted some software described in his article on the generation of an ASCII character set (or special characters) for a plotter or vector display device.

|  |  | 9U <br> JULY 1976 |
| :---: | :---: | :---: |
| 6 | Foreground <br> COINCIDENT CURRENT FERRITE CORE MEMORIES Recycling - Jones | PUBLISHERS <br> Virginia Peschke Manfred Peschke EDITOR <br> Carl T Helmers, Jr |
| 22 | EXPLORE AN 8080 WITH EDUCATOR-8080 <br> Software - Howerton | GENERAL MANAGER <br> Manfred Peschke <br> PRODUCTION MANAGER |
| 38 | THE "IGNORANCE IS BLISS" TV DRIVE CIRCUIT Hardware - Barbier | Judith Havey PRODUCTION ASSISTANT |
| 46 | WHY WAIT? BUILD A FAST CASSETTE INTERFACE Hardware - Suding | Elizabeth Alpaugh CIRCULATION Deborah R Luhrs |
|  | Background | DEALER CIRCULATION Deena Zealy |
| 18 | ASSEMBLING A SPHERE <br> Review - Anderson | PUBLISHERS ASSISTANTS <br> Cheryl Hurd <br> Carol Nyland |
| 30 | MACHINE LANGUAGE PROGRAMMING FOR THE " 8008 "-Chapter 1 Software - Wadsworth | ADVERTISING Elizabeth Alpaugh Virginia Peschke |
| 40 | PUT THE "DO EVERYTHING" CHIP IN YOUR NEXT DESIGN | TYPOGRAPHY |
|  | Hardware - Baker | Custom Marketing Resources, Inc Goodway Graphics |
| 54 | SURPLUS ELECTRONICS IN TOKYO AND MANILA | Mary Lavoie |
|  | Travelogue - Hayes | Taimi Woodward |
| 58 | MAKE YOUR OWN PRINTED CIRCUITS | PHOTOGRAPHY |
|  | Hardware - Hogenson | Ed Crabtree <br> Custom Marketing Resources, Inc |
| 64 | A PLOT IS INCOMPLETE WITHOUT CHARACTERS <br> Software - Lerseth | ART <br> Mary Jane Frohlich <br> Bill Morello <br> PRINTING |
|  | Nucleus | Custom Marketing Resources, Inc The George Banta Company ASSOCIATES |
| 2 | In This BYTE | Bob Baker |
| 4 | The Trend Toward Hassle Free Products | Dan Fylstra |
| 39 | Tool Box | Don Lancaster Harold A Mauch |
| 44,79 | BYTE's Bits | Chris Ryland |
| 74 | Good Grief! | CONTEST EDITOR |
| 76, 102 | What's New? | Janice D Black |
| 81 | Software Bug of the Month 2 | BYTE magazine is published monthly by BYTE Publica- |
| 83 | What Happened at Trenton | tions, Inc., 70 Main St, Peterborough New Hampshire |
| 86 | Clubs, Newsletters | 03458. Subscription rates are |
| 90 | Letters | Two years, $\$ 22$ yer Three y years. |
| 95 | Ask BYTE | $\$ 30$. Second class postage paid at Peterborough New Hamp- |
| 96 | Reader's Service | shire 03458 and at additional mailing offices. Phone |
| 98 | Classified Ads | 603-924-7217. Entire contents copyright 1976 BYTE Publica- |
| 100 | BYTE's Bugs | tions, Inc, Peterborough NH |
| 104 | Book Reviews | respondence to Editor, BYTE, |
| 112 | BOMB | 70 Main St, Peterborough NH 03458. |

# The Trend Toward Hassle Free Products 

Editorial by Carl Helmers

[^0]One item which backyard entrepeneurs tend to overlook is the fact that the hardware and software engineering of computer systems which are both manufacturable and marketable is a complex process. I had a taste of the problems of working out the details of a system about the time BYTE became a much more compelling personal project in mid 1975. I quickly dropped any illusions of converting the designs I was working upon at that time into a manufactured product when BYTE with its own intricate and unique challenges came into being. But the period of time spent thinking about the problem of engineering a salable combination of hardware and software leaves an impression.

There are a myriad of details which must converge into a well defined, usable product for personal computing applications. These details are handled naturally and implicitly when the advanced amateur or full time engineer assembles a "one of a kind" system as a personal vehicle or engineering test bed. But documenting the critical parameters, converting a one of a kind electronic system into a mass produced design, providing the level of support needed and expected by customers - these are not trivial tasks at all. One of the wonders of the field to this date is that the complicated objects of our computer affections are as free of design and assembly hassles as they are at the present time. Simply look at all the system design principles which have been properly implemented and provided in existing products, and for the moment ignore the obvious residual improvements that would make a complicated product "more perfect" by some standard.

But technology almost by definition is not static. One of the significant indications of personal computing's ever improving
technology is the trend toward minimizing hassles involved in getting a production model computer system up and running. This trend affects both existing product lines and new products which will shortly come to market.

- Item: In talks with the MITS people at the recent World Altair Computer Convention, one point which was made is that an increasing number of customers order assembled and tested Altairs. This delivery of assembled units, both direct from the factory and through computer stores, significantly reduces the time and hassle overhead of getting a system up and running.
- Item: In a recent visit to Sphere, the same story was heard: While initially the kit orders predominated, a growing number of purchasers are opting for assembled systems.
- Item: A new firm (let it be known as "brand A" for now) recently contacted BYTE with a report on its product, the inspiration for this commentary. The product's significance is that it comes in one and only one form: a completely tested board of moderate size which needs only a power supply, monitor with EIA video input, keyboard with parallel TTL interface and audio cassette recorder to complete the system. The price of the board is well under $\$ 1000$ and a BASIC software package is thrown in as part of the deal. The memory capacity of this 6502 based system was quoted as 8 K , with an option to replace the 4 K dynamic memory chips with pin compatible 16 K chips


# 1,000,000 

 BYTES ON-LINE!

# The Digital GroupCassette Storage System 

The Digital Group Cassette Storage System gives you total magnetic tape data storage and retrieval for your microprocessor, capable of operating 1 to 4 computer-controlled PhiDeck cassette transports. Within seconds (20 at most), your system zips to any of over one-quarter million 8 -bit bytes per drive. And that really puts it all on-line!
The Digital Group Cassette Storage System is ideal for:

- Large data files - names, accounts, etc.
- Indexed computer-controlled program files
- Sorts
- Inexpensive mass storage
- Work files
- Indexed random retrieval
- Multi-pass compilers
- System residence

In addition, with a Digital Group System and a Phi-Deck transport, your total load procedure is reduced to a single action - turning on power. Everything else is automatic! Your Digital Group System is completely ready for use in a very few seconds. And you avoid a large investment in single-use PROM memory.

## MAJOR STORAGE SYSTEM COMPONENTS

1. Controlling and Formatting Interface - single card for 1 to 4 drives
2. Software Operating System
3. Computer-controlled Cassette Drive(s)

## Selected Specifications

Data Rate: 800 bytes per second, 8 K loads in 10 seconds Media: High-quality standard audio cassettes
Search Speed: 100 inches per second
Tape Speed: 5 inches per second

Power Requirements: +12 V to +20 V at .7 A peak and +5 V at 1 A plus 60 ma per drive
Port Requirements: One 8 -bit parallel input port plus two 8-bit parallel output ports
Cassette Drive is an enhanced Phi-Deck with a digital head, cast head bar, stronger capstan, and four-foot cabling.

## SOFTWARE OPERATING SYSTEM

## 8080 based -650 bytes

Error Detection: CRC
Retries after soft errors
Automatically bypasses hard errors
Block size $=1$ to 256 bytes or multiple of 256 bytes
Functions supplied:

- Record multiple blocks
- Record 1 block
- Read 1 block
- CRC check
- Fast reverse
- Fast forward
- Search for block

For more information, drop us a line or call . . . but by all means, get on our mailing list.
Prices: Interface - full kit PHI-F . . . . . . . . . . . $\$ 135 \mathrm{ppd}$
Each Drive - assem. PHI-1 . . . . . . . . . \$115 ppd


THEDIGITAL GROUP INC
P.O. BOX 6528

DENVER,CO 80206
(303),861-1686


Figure 1: For a wire carrying current i directed out of the page, the magnetic field H is proportional to $i / r$. The direction of H is tangential to the circle of radius $r$, as shown for points $P_{7}, P_{2}$ and $P_{3}$.

## Coincident Current

## Ferrite Core Memories

James R Jones<br>111 E Jefferson Colorado Springs CO 80907

Have you ever looked through the surplus catalogs and wondered whether those memory core planes and stacks advertised could be used for anything other than tea strainers? How are they supposed to work, and what kind of circuitry is required? Could you troubleshoot and repair that surplus core memory and drive electronics? What good is a core stack without data? Is it a worthwhile project to build a memory system around a core stack?

Some of these questions can be answered directly in what follows. Others can be answered at least partially in terms of my own experience: getting one surplus core stack to work as the main memory of my home brew computer by building the necessary electronics without the benefit of manufacturer's data.

Perhaps the biggest advantage of the core memory over other types of random access memories is its ability to retain stored data when power is removed, and to have it readily available when power is restored. This feature is sometimes used by minicomputer manufacturers to ship their products with preloaded systems software. Also, for some forms of core memory, the storage capacity increases much faster than the amount of driving electronics. In order to understand the functional requirements and the operating restrictions placed on core
memory driving circuitry, it is necessary to take a close look at how cores work, and how they are typically organized to form large capacity stacks.

## Magnetic Fields and Hysteresis

The memory storage element is the ferrite core itself. Its function is to accept, store, and read out a bit of information. It can do this by virtue of its ferromagnetic properties. These allow it to be easily, but strongly, magnetized in a preferred direction by an externally applied magnetic field to signify a 1 bit value, to be easily magnetized in another direction to signify a 0 bit value, and to retain its magnetic direction when the external field is removed. Physically the core is a small doughnut shaped object made of pressed, heat treated, non-conductive, iron oxide powder. A typical dimension is 0.02 to 0.2 inches in diameter ( 0.5 to 5 mm ). The external magnetic fields are applied to the core by means of wires passing through it carrying controlled amounts of current.

A straight wire carrying a current gives rise to a vector field, $H$, in the surrounding space, called the magnetic field strength. $\mathbf{H}$ is proportional to the current $i$, and is a vector directed as shown in figure 1. The arrows indicate the directions the north end of a compass needle would point if placed at the points shown. Another vector quantity, the

magnetic field of induction, $\mathbf{B}$, is also considered to exist, and describes the overall magnetic effects due to H and the presence of matter. In figure 1, B behaves like H. If the wire passes through a ferrite core, the $\mathbf{H}$ field generated by the wire causes molecular and microscopic sized magnetic domains (regions magnetized like the compass needle) in the core to align themselves in the direction of $H$, resulting in a $\mathbf{B}$ field in the core with direction around the circumference of the core as shown in figure 2. This field is much stronger than that due to the wire alone, because the structures in the core add their fields to the wire field when aligned.

Because of the geometry of the cores with respect to the wire, it is possible for us to drop the vector notation and simply refer to the magnitude of $B$ and $H$ in the discussion which follows. The exact way the strength of $\mathbf{B}$ in the core depends on the strength of applied field $\mathbf{H}$ (proportional to current i) is represented graphically in the scalar $B$ versus $H$ curve of figure 3. After manufacture, the core has no magnetization, and $B=0$. If applied field $H$ is increased from 0 to $\mathrm{Hp}, \mathrm{B}$ increases along path $0-\mathrm{l}-\mathrm{c}-\mathrm{e}$. If $H$ decreases to $0, B$ decreases from $e$ to $b$, leaving the core magnetized with no applied field present. Decreasing H from 0 to Hm causes $B$ to move along b-a-g-f, changing direction as it passes through 0 . As $H$ increases to $0, B$ moves to $h$, again leaving the core magnetized with no applied field, but in a direction opposite that above. As H increases again from 0 to $\mathrm{Hp}, \mathrm{B}$ moves along $\mathrm{h}-\mathrm{j}-\mathrm{d}$-e. If H is again varied as described above, B will trace the same counterclockwise path.

This effect of tracing two alternate paths between two fixed states, instead of tracing a single path, is called hysteresis. In the core,


Figure 2: The direction of magnetic induction B in the core results from alignment of core structures with the magnetic field due to current i.


Figure 3: Memory properties are derived from the hysteresis properties of the core material. When the applied $H$ value is 0 , the residual $B$ value in the core material will depend upon the past history of magnetization. This figure and subsequent figures show the magnitude of the vector quantities $\mathbf{H}$ and $\mathbf{B}$, with positive and negative values corresponding to the two directions of rotation about the core axis.

this is due to the alignment of the magnetic domains, a process which requires a minimum applied field to force alignment in a given direction. When the applied field is removed, the aligned structures provide the core with its own field. Path segments a to $g$ and j to d , where this effect occurs, are termed irreversible. Reversible magnetic effects also occur in the core. For example, if $B$ is at position $h, H$ can be varied to move $B$ back and forth along line $\mathrm{f}-\mathrm{j}$ as many times as desired as long as $B$ is not increased past point j . Variations along line a-e are also reversible to point a. This effect is due to elastic deformation of the microscopic structures and alignment of molecular structures, both of which return to normal when the applied field is removed. Both effects are important: Hysteresis is the memory pro-

Figure 4: Sense voltages perceived on a wire strung through the core depend upon the previous $\mathbf{B}$ magnetization state in the core as a current pulse is passed through the wire. A non-switching core produces a simple transient due to reversible B field changes induced in the core by the current edges as at (a). When the core switches as at (b), an extra transient pulse is induced in the sense wire.
perty; reversibility, although the largest noise contributor in the core, allows the core to act as an AND gate, a key principle in selection of individual bits.

## Sensing Core Magnetism Changes

Changes in the core B field are sensed by means of a second wire passed through the core. The voltage difference appearing at the ends of the sense wire is proportional to the change in B , and inversely proportional to the time required for the change. So that an irreversible change can produce a large signal, $H$ is made to change as fast as possible. However, this causes the reversible changes to produce large signals also; but the speed tends to separate the signals: Deformation effects can occur much more rapidly than gross microscopic realignments, eg: for a fast increase of H from 0 to Hp in figure 3, B would not follow h-j-d-e, but rather h-j-s-e, so most reversible changes occur first. Figure 4 illustrates the voltages appearing across the sense wire for a core in initial state $b$, and in initial state $h$, resulting from a fast rising current pulse generating field Hp .

## Using the Core to Remember

If we now interpret the initial direction of $\mathbf{B}$ in the cores to represent a stored 0

INITIAL 8 STATE


Figure 5: Using the physics of core switching. This figure shows the four basic operations needed to manipulate the binary state of a memory core: Read a " 0 ", read a " 7 ", write $a$ " 0 " and write a " 7 ".


Figure 6: How hysteresis makes coincident current bit selection possible. The principle involved is simple: A current of $i$ is sufficient to generate an $H$ field strength which will drive the core from one state to the opposite state, but a current of i/2 is insufficient to cause the core to cross the threshold of the hysteresis curve. The H field of two separate wires add so that if two wires carry currents of $i / 2$ in the same direction, the result is the same as a single wire carrying a current of $l$. The hysteresis curve at the right (d) shows the paths taken by the B field in the case of a single i/2 current (a), additive i/2 currents forcing a change (b), and additive i/2 currents causing no change since the core is already in a state aligned with the $H$ field (c).
(figure 4a) or a stored 1 (figure 4b), and read the core contents by the method of figure 4 , then read and write operations can be performed as indicated in figure 5. Two things should be noted: First, a stored 1 is destroyed in the process of reading it because the core is reset to the 0 state. If it is necessary to retain the 1 in memory for later use, a special refresh write operation is required to restore the 1 . Second, the cases of writing into a core already containing a 1 are not shown, because the memory control logic is usually designed such that a destructive read cycle always occurs before the write cycle, for a given core.

## Memories Larger Than One Bit

Using the method of figure 5 can be an expensive undertaking if one wants to store 4096 bits, for instance. Although the sense line can be made to pass through all the cores to detect switching of a selected core, circuitry must be built which can select any one of 4096 cores and supply the proper read or write current to it alone. To reduce the amount of selection circuitry required

Figure 7: The concept of a coincident current selection is implemented using an array of cores with $X$ and $Y$ lines. In this example, a total of 20 cores is wired with a single sense wire, $4 Y$ selection wires and $5 X$ selection wires. When i/2 (see figure 6) is flowing through one $X$ and one $Y$ line, the matrix intersection point is addressed. In this example, the core at $(2,3)$ will either be read or written. Actual core arrays are much larger than this simple conceptual illustration.
for large memories, the cores' reversibility property is used to assist in the selection process. This is done by passing two independently controlled current lines instead of just one through a given core; each line can pass half the current of the original line. Due to the shape of the B versus H curve, if only one line passes half the required switching current required, B remains in the reversible region and returns to the original state (figure 6a) when the current is removed. If both lines pass current, core switching occurs if the core is in the proper original state (figure 6 b and c ).

## Coincidence Current Selection

Thus by the means of coincident currents, the core acts as an AND gate; both currents need to be present simultaneously


Figure 8: The half-select current i/2 must generate an H field less than Hk if a change of state is to be avoided. Similarly, the full select current of i must generate an $H$ field strength greater than Hs to ensure switching. This leaves a range of field strengths (and corresponding current values) from $\mathrm{Hs} / 2$ to Hk which introduces a critical tolerance for the select currents in a core memory design.

Figure 9: The inhibit line is added to the array assembly of cores so that it is possible to avoid changing the state of a given core with coincident current selection.
to choose the core and allow it to switch if it is in the proper state. A multiple bit memory can now be constructed as shown in figure 7. If currents are applied to lines X2 and Y3, the core at $(2,3)$ is selected. Other cores on the lines X 2 and Y 3 are half selected. Only one pair of lines is allowed to carry current at any one time. To read the core contents at $(2,3)$, current is passed as indicated. To

now write a 1 , the current is reversed; to write a 0 , no current is applied.

## Property and Operation Restrictions

Arrangement of cores in the coincidence current mode imposes some rather severe requirements on core fabrication and memory operating conditions. For a core in B state $x$ in figure 8, the half select current must generate a field less than Hk to prevent B from passing the "knee" of the curve. But the total select current must generate a field greater than Hs to ensure full switching of the core. Thus the half select current can only produce fields in the range shown by the heavy line. As the quantity ( $\mathrm{Hs}-\mathrm{Hk}$ ) gets smaller, the operating range gets larger. So the cores are made with ( $\mathrm{Hs}-\mathrm{Hk}$ ) as small as possible, but even if it were zero, the half select current would have a maximum tolerance of $33 \%$ of its midrange value. The operating range is further restricted by two more factors. First, manufacturing and assembly tolerances allow memories to be constructed of cores with slightly differing B versus $H$ curves, effectively increasing the difference ( $\mathrm{Hs}-\mathrm{Hk}$ ) as all cores are considered. Second, as the memory is operated, the cores heat up due to switching losses, select line heat, and heat from nearby electronics, causing the entire $B$ versus $H$ curve to shrink towards the origin. Newer core memories are made using ferrite with low temperature sensitivity; but this is not true of older memories, and such measures as constant temperature ovens, forced air cooling, and power supply temperature compensation have been used to ensure reliable operation.

## Reading and Writing Words: The 3D Memory

The coincidence current scheme certainly cuts down the required selection circuitry, although at the expense of operating tolerances. The 20 cores of figure 7 are selected by nine lines, and the 4096 bit memory mentioned earlier, if arranged in a $64 \times 64$ array, has only 128 select lines. Further economy is realized if the idea is expanded to word organized memories. For example, consider the 16 bit array of figure 9 which has an added winding - the inhibit line (the arrows on the select lines indicate the read current directions). If a 16 word memory of 3 bits per word is desired, three of these bit planes are interconnected as shown in figure 10.

Sending half select currents through X 1 and Y 1 select lines in the arrow directions will cause the core at $(1,1)$ in each array to read out via the associated array sense line. Reversing the select current directions will


Figure 10: 3D Plane Interconnections. In a 3D memory, three dimensions of addressing are implicit: There is an $X$ and $Y$ dimension for each plane, and one plane is used for each bit of the parallel computer word. In this picture, three planes of a 3 bit wide memory are shown. Only two lines are shown connected in each $X$ and $Y$ direction. Other $X$ and $Y$ lines are connected between planes in a similar fashion.
cause a 1 to be written in each (1,1) position. In order to write a 0 in some of the $(1,1)$ positions, a half select current is also applied to the desired inhibit lines in the direction of the arrows at the time 1 's are to be written. This inhibit current appears as read half select current at all cores of the associated array, cancelling half the write select current at the selected core, and thus preventing the writing of a 1 at that core.

For a 4 K word memory of 16 bit words, the assembly (called a stack since the arrays are often stacked like pancakes for simple interconnection along the edges) would have 128 select lines, 16 sense lines, and 16 inhibit lines. The connection scheme is termed 3D, and is most economical in terms of support circuitry.

## Other Arrangements

Bit oriented memories of, the form of figure 7 are called 2D memories. Another commonly used word oriented connection scheme, $2 \frac{1}{2}$ D, will not be discussed because of the relatively large amount of select circuitry used. Interested readers are referred to the bibliography at the end of this article. A variation of the 3D memory replaces the inhibit and sense lines by a single line which performs both functions, since they occur at separate times. At the cost of additional circuit complexity, the variation allows smaller cores and more closely packed arrays to be used.

## Minimizing the Noise

A core can be oriented two ways at a select line intersection, and sense and inhibit
lines can string the cores of a plane in many patterns. Advantage is taken of these facts to minimize noise from sources which could otherwise mask the sense line signal due to core switching. Major noise contributors are pulses from half selected cores (path a in figure 6 d ). Although small for a single core, the composite signal due to 126 half selected cores in a 4 K array will spread and totally mask a switching signal. A diagonal pattern that results in almost complete cancellation of half select pulses is shown in figure 11a. Noise is also generated through capacitive coupling between select and sense lines, allowing select current edges to induce ringing in the sense line. This effect can be minimized by running the sense line parallel to one select line direction as shown in figure 11b. The parallel select line is turned on first, a large noise pulse is induced which dies down, then the other line is turned on to select the core. Little noise is now generated because capacitive coupling is small.

There are other lesser noise effects which make the 4096 core array about the largest to be practically served by a single sense line. Larger stack word capacities are realized by using multiple sense lines on each plane, thus requiring additional sense circuitry.

## Driving the 3D Memory

The half select current tolerances mentioned in connection with figure 8 are further reduced to less than $20 \%$ per line by the addition of the inhibit line. Unfortunately, the amplitude of the core switch signal depends on the total select current


present, so the select current must be maintained at the higher allowed values to obtain a favorable signal to noise ratio on the sense line. This reduces current tolerances to around $5 \%$ to $10 \%$. Depending on the type of cores used, half select currents are in the range of 150 to 500 mA . Current

Figure 11: Noise pickup is minimized by threading the sense line through the core array in a way which causes induced noise voltages to cancel each other.
rise time of the last select line turned on (to read) should be about half the core design time for a 1 peak (figure 4 b ) to occur, which ranges from 100 nanoseconds to about a microsecond for older large diameter cores. Current should last for at least twice the peaking time to allow the core to switch completely.

## Getting a Memory to Work (Starting from Scratch)

My surplus Spectra 70 stack (built by Electronic Memories, Inc) arrived with no information except that its organization was $128 \times 136$ words $\times 18$ bits with four sense line groups of $64 \times 68$ words. Of that 1 intended to initially use a $64 \times 64$ word $\times 16$ bit segment. Not being able to obtain manufacturer's data, I proceeded to discover the stack characteristics for myself. For readers in similar situations, I would recommend such a course of action only if you have lots of spare time, like surprises, and have access to a 15 MHz dual trace scope.

The physical dimensions of the stack were 5 by 5 by 3 inches ( 12.7 by 12.7 by 7.6 cm ) not counting connectors. Tamper proof, auto destruct construction precluded a view of cores or windings. $X$ and $Y$ lines were labelled; an ohmmeter was used to associate sense and inhibit connector pins to core planes. Memory construction suggested the rectangular winding of figure 11b.

## Tests

Several test circuits were built to determine core switching characteristics. Figure 12 shows the select line driver circuit, two of which are required to drive an $X$ and a $Y$ line. The Texas Instruments Memory Core Driver integrated circuit (SN75325) is used. The external diodes simulate the steering diodes required to select more than one line, and to direct reverse voltages (generated at current turn off time by the inductive load of the select lines) to the sink transistor collector which is internally diode clamped to source voltage V s. Vs is regulated and variable over about a 5 volts range, and R1 is chosen for the current range to be tried. Figure 13 shows the inhibit driver circuit. The dual gates of the integrated circuit are connected in parallel to increase the current sinking capacity. Again the output must be diode clamped, and R2 is chosen as required. Figure 14a shows the sense amplifier with terminating resistors and means for varying the threshold voltage. The sense amplifier operates by generating an internal voltage proportional to the voltage difference appearing on the sense line wire pair. This voltage is compared with the adjustable
threshold reference voltage, and a TTL high level output is generated if the threshold is exceeded. Figure 146 illustrates the outputs produced at two threshold levels for an arbitrary input signal.

Timing signals used are shown in figure 15 along with resulting current directions. The logic diagram illustrates a circuit which can be used to generate this timing. The enveloping of $X$ ENABLE and $Y$ ENABLE eliminates part of the half select noise in the read cycle, and ensures that the inhibit current is at maximum level when full write selection takes place. Repetition rate of the memory cycle during testing is kept under 1 kHz to prevent damage to the SN75325s. Changing the mode, WRITE, with both enable signals low ensures that a glitch will not short current through a directly connected source to sink pair.

Two select line drivers were connected to an arbitrary select line pair, and the sense amplifier was connected to a sense line thought to pass through the $X Y$ intersection. With no signals to the select line drivers, the sense amplifier threshold voltage was advanced from zero until SENSE OUT was a steady logic 0 . Starting with a low (about 150 mA ) half select current, the line drivers were started, and, with the sense amplifier output monitored on the oscilloscope, the threshold voltage was adjusted to show a series of spikes corresponding to select current edges as shown in figure 16 a . As half select currents were increased equally, additional pulses emerged after spikes numbered 3 and 11, as indicated by the arrows in figure 16a. These pulses indicated core switching. Some experimenting was required to obtain proper select current balance and strength, which occurs when the switching pulses completely disappear (and not just shift) when either half select current is temporarily interrupted. Too much drive current was also indicated by a decrease in size of the switching pulses with corresponding pulses appearing behind spikes number 1 and 9.

Choosing inhibit drive resistor $R 2$ to provide half select current in the range determined necessary above, the inhibit line corresponding to the sense line used was found by driving an inhibit current through each inhibit line in turn until the switching pulses were seriously disturbed. Further adjustments of the inhibit, $X$, and $Y$ currents eventually resulted in the switching pulse appearing and disappearing as the inhibit current is turned off and on. In my memory stack, about 380 to 400 mA half select current is required to cause the core to switch, producing a sense amplifier output


Figure 13: Inhibit Line Test Driver. For testing, one of these circuits is used to pick an inhibit line. R2 is a current limiting resistor set as described in the text of the article.

b.


Figure 14: Sense Amplifier. The purpose of the sense amplifier is to create an output signal which canl be strobed into a TTL latch at the appropriate moment, given the signals induced into the sense wire of the memory. For testing, the thresholds are set using the variable resistance of 200 ohms in a voltage divider from the +5 volt supply.
pulse sketched in figure 16b. The smaller pulse generated for a 0 is also shown. The 1 pulse at the sense line terminals is on the order of 40 mV amplitude.

An experiment verified the memory has a rectangular sense winding parallel to the Y select lines, requiring the latter to be turned on first in the read cycle. This was done by swapping the $X$ ENABLE and $Y$ ENABLE signals, and noting which arrangement pro-
duced the most sense noise at clock time 3 of figure 15.

## Selecting the Select Lines

Success in getting the stack to respond to prodding induced me to go ahead with memory circuit construction. The SN75325s were employed in a commonly used scheme of line selection illustrated in figure 17, in which one of four $X$ lines is chosen by four source to sink pairs, determined by a two bit

Figure 15a:


I KHz DURING TESTIN
CONTROLLED BY CPU IN SYSTEM OPERATION.


Figure 15b:


Figure 15: A suggested circuit (a) and timing wave forms (b) for driving core memories in testing and later in the final working version. The cycle time is set by the clock into pin 5 of the counter IC1. A positive transition at pin 11 of IC6 initiates a single memory cycle illustrated in (b). During testing, approximately 1 kHz will provide a good repetition rate which avoids burning up driver circuits; in final system operation, this circuit can also be used to generate the read then write cycle of core memory operation. For an $8 \mu \mathrm{~s}$ full cycle time, the clock should be 2 MHz .
address (low order bits 0 and 1). $X$ READ ENABLE equals $X$ ENABLE AND NOT WRITE; X WRITE ENABLE equals $X$ ENABLE AND WRITE. By extension of this scheme (see the TI catalog listed in the bibliography), 16 source to sink pairs are required to select one of $64 \times$ lines. I found that a single source resistor (R1 of figure 12) could service all source transistors used to select $X$ drive lines, which eliminated the necessity of matching resistor values to provide the same current flow as different lines are selected. Two 3 to 8 decoders use six address bits to enable two source to sink pairs. The above circuitry is duplicated for the 64 Y lines, thus a 12 bit address is decoded into a unique $X Y$ intersection of the core matrix. A total of 16 SN75325 integrated circuits and 4 decoders are required, along with diode arrays. Three independently regulated and variable supply voltages were provided to allow adjustment of X, Y, and inhibit currents. Inhibit drive resistors (R2 of figure 13) must be closely matched to allow operation from a common supply voltage.

## Memory Operation

Figure 18 illustrates control and data flow of the memory logic. Read data is stored, by means of the Read Strobe signal, in a buffer where it is used for restoring the memory word in a central processor read operation. In a central processor write operation, this read information is ignored and data from the CPU controls the inhibit drivers during memory writing. Timing is of the form of figure 15 , except that the

b. SENSE OUT

VOLTS


Figure 16: Sense Amplifier Outputs. Spikes will be seen in the sense amplifier outputs following current changes through the cores. Detection of a " 0 " or "1" is obtained by observing at the "read strobe point" which follows the leading edge of the $X$ enable pulse. This observation point is typically 400 nanoseconds after the leading edge of the $X$ enable pulse which drives one of the $X$ select lines.
memory cycle (read then write) is accelerated to 8 microseconds in my system. Inhibit and sense circuits are similar to those of figures 13 and 14a.

The arrangement of figure 18 is wasteful of central processor time, because data and address information must be held until the memory cycle is completed. Buffering data and address information in the memory logic would allow the processor to continue after

Figure 17: X Line Select Switch. This arrangement is an extension of that illustrated for testing purposes in figure 12. For 75325 pin connections see figure 12, or page 10-21 of The Linear and Interface Circuits Data Book for Design Engineers by Texas Instruments, 1973 edition.



Figure 18: System diagram of a memory built according to the principles described in this article. The basic memory cycle is set up by a circuit such as the one shown in figure 15 as a source of timing control. In a read cycle, data is first read, then rewritten; in a write cycle, old data is read and new data written.
issuing a write command, or to wait for only half the memory cycle for data needed by a read operation.

The memory is coarsely adjusted by electronically forcing repeated manual write cycles through the processor's control panel. This allows adjustment of the half select and inhibit currents, sense amplifier thresholds, and read strobe. The read strobe point is adjusted to occur at the point of sense amplifier output that provides best discrimination between 0 and 1 levels (see figure 16b). Short program loops can then be used to make finer adjustments.

## Conclusion

Recently completed, the $4096 \times 16$ memory subsystem described here was built for less than $\$ 400$, including stack, power supplies, and miscellaneous electronic and mechanical parts; and can be expanded to 16 K for about $\$ 200$. In most of the memory, storage reliability is good; although a few blocks of addresses tend to pick up random bits. Bugs have been eliminated in a straightforward manner, most having been traced to poorly made connections to select lines and to a bad batch of inhibit driver integrated circuits. Building a memory
around a core stack is an interesting but time consuming project that can definitely result in a usable end product. As with many surplus items however, performance depends a lot on how much you can find out about the stack, and on the hardware design concessions that have to be made in the interest of financial solvency.

## BIBLIOGRAPHY

Fundamental Principles of Switching Circuits and Systems. New York: American Telephone and Telegraph Company, 1961.
The Integrated Circuits Catalog for Design Engineers. Dallas: Texas Instruments, Inc (First Edition).
Daly, Edmund B: "Noise Problems in Ferrite Core Coincident Current Memories." Computer Design (November 1967), 46-50.
Hawkins, Joseph K: Circuit Design of Digital Computers. New York: John Wiley and Sons, Inc, 1968.
Kirkpatrick, T, and D Johnson: "Design of Modular, Low-Cost Core Memories for SLT Systems." Computer Design (April 1968), 60-71.
Smith, James E: "A 650 Nanosecond 2½D Memory." Computer Design (October 1966), 80-87.
Zinschlag, Howard P: "A 2½D Integrated Circuit Memory." Computer Design (September 1966), 26-39.

Introducing SCELBAL, the new microcomputer language that's simpler than machine language.

A complete language for "8008"/'"8080" systems including source listings, routines, flow charts and more!

SCELBAL. SCientific ELementary BAsic Language for " 8008 "/'" 8080 ' systems. A complete, illustrated program book. Routines. Techniques. Source Listings. Flow Charts. And more. Took several years to develop. Now yours for many years to come. First time that intimate details of higher level language has been offered for such a low price. Only $\$ 49$ ! You get 5 Commands: SCR, LIST, RUN, SAVE, LOAD. 14 Statements: REM, LET, IF ...THEN, GOTO, FOR with STEP, END, INPUT, PRINT, NEXT, GOSUB, RETURN and optional DIM. 7 Functions: INT, SGN, ABS, SQR, RND, CHR, TAB. And, it runs in 8 K and more. Here's all the data needed to customize a high level language for your system....at a fraction of the cost!

Order your copy today! Get $\$ 4.80$ ppd.


Here's SCELBl's First Book of Computer Games for the 8008/8080. Action-packed. And fun.
Try to beat the computer at its own game. Here's the first complete machine language computer manual for computer games to include source listings, flow charts, routines and more. Space Capture - You against the computer using "search and destroy" strategy to shoot down roaming alien spaceships in outerspace. Hexpawn - a mini-chess game that lets the computer make mistakes ... but only once. Hangman - an updated version of the great kid game. Computer selects words at random from long, expandable list. Try to beat it in 8 moves or less. lilustrated. Fun extras to put your computer to challenging, competitive, fun use. Order yours today! $\mathbf{1 4}_{\text {ppd. }}^{95}$

## SCELBI's new microcomputer Game Book of Outer Space... C/

Captain your own crusading starship
against the logic of your " $8008^{\prime \prime}$ or "8080'. You must destroy a random number of allen ships without running out of stardates, out of fuel, out of ammunition or out of the galaxy. Plan your mission. How much fuel is used for your warp factor? Don't run into roaming stars that can damage your ship. Suddenly! Condition Red! Alien In sight! How big is he? Fire a phasor or torpedo! He's damaged or destroyed. But, you've used valuable fuel. That's Just the beginning. A new game every time. Complete program book in machine language for 4 K memory, including source listings, flow charts, routines, etc. Great intergalactic adventure and fun.


Order "GALAXY" today! $\mathbf{1 4}^{\mathbf{9 5}}$ ppd.

$19^{95}$ ppd.

## MACHINE LANGUAGE <br> PROGRAMMING - Detailed presentation of "8008" codes • Flow Charts <br> - Mapping - Floating-point Package - Debugging <br> - Basic programs: loops, counters, masks <br> - Organizing Tables - Editing/Assembling <br> - Math operations - !/O, Real Time Programming <br> - Maximizing memories - And lots more. <br> Here's the detalled, basic manual you need to

develop today's machine language programs. 170 pages. Illustrated. Easy-to-read, understand. Mos technlques appllcable to other micro-computers, Including "8080". FloatIng-polnt arlthmetic package is worth the price alone!
Order todayl

Master Charge, Postal and Bank Money Orders prelerred. Personal checks delay shipping up to 4 weaks. Pricing, specificatlons, avallability subject to change without notice. Prices for U.S. and Canadian dellvery at book mailing rate. Add $\$ 2.50$ for each publicatlon If Priority Alr Service (U.S.) desired. Foreign orders add $\$ 6.00$ for each publication.

1322 Rear Boston Post Rd., Milford, CT 06460 Telephone: 203/874-1573

## Assembling a Sphere

Bruce A Anderson 4554 Chinook Ct
San Diego CA 92117

Sometime in July an advertisement was being run in several magazines which offered a complete computer system for only $\$ 650$. Up until this time I had been interested in microprocessors but had not even remotely considered buying one. 1 sent off for the brochure anyway. After reading the claims


Photo 1: The Sphere system features a completely separate power supply. All regulation is done in the unit shown here, which eliminates the need for on board regulators and the extra heat dissipation on logic cards.
about the system, I quickly convinced myself that I could justify the expense as an educational expense; so I sent off a check to Bountiful UT where resides Sphere Corporation, the company making these claims.
I ended up ordering a "SYS2/ KIT - Intelligent" which included a Motorola 6800 processor with 4 KB of programmable memory, full keyboard, TV interface, a cassette interface and what was probably the biggest selling point, 1 KB of EROM which had software for what Sphere calls a "mini assembler," text editor and a debug routine. All this (during the special introductory period) for only $\$ 750$. Clearly a bargain!

By mid August $I$ had received confirmation of my order and a promised delivery date of the second week in October. Part way through this 60 day wait I heard a rumor that Sphere had not delivered any systems and that all the company consisted of was two people in a garage. When the second week in October arrived and my kit didn't, I put in my first call to Utah and discovered that it would be about two weeks late because they hadn't been able to get all the parts from their suppliers on time to put the kits together. October 30, as promised, brought a box from Sphere which had most of the electronics (all except the cassette interface) but no metalwork; so the next day I again called, this time to find out about the metalwork and two or three parts which were not exactly as specified on the parts list. Fortunately all the parts were usable substitutions and the metalwork arrived the same day so I commenced to build.

Sockets are provided for all the expensive MOS devices, but I decided to use sockets on
everything so for the next four days 1 soldered in sockets in all my spare time. I was soon ready to plug it in except that the interconnection and power cables had been back ordered and had not come yet. So instead I inspected the boards again, modified a TV for monitor use and waited for a week until they arrived.

Finally I was ready to try out my system; and sweating profusely, I threw the power switch. When no smoke appeared I let out a sigh of relief and then began trying to find out why it wasn't doing what I thought it should be doing.

The television interface seemed to be working, but all that was showing on the screen was a set of random characters and none of the keys would affect anything. Looking at the address lines with a scope, it became obvious that the processor was stuck in a loop. One of the things noticeably lacking from the manual was a set of debugging procedures for the hardware, so again I called Utah. In response to the question of debugging procedures, I was told that it was almost certainly a solder splash on the address lines, although it was possibly the refresh clock notworking.

Even though I had already checked for solder splashes at least twice I went back and looked again just to make sure. When no splashes turned up and the refresh clock appeared to be working correctly, I tried finding out what memory locations were being addressed in the loop (no simple task with only a single trace scope). Finally I pulled out the address drivers (here is where my foresight in using sockets paid off) and manually input the addresses which were in the loop to see what was in these locations in the PROM. Sure enough, with the program which was in the PROM, once the microprocessor got to this point (which it would upon startup), there was no way to get out of the loop.

Armed with what I thought would be convincing data I again called Sphere and this time was asked to send in the PROMs along with a $\$ 5$ shipping and handling charge. Two weeks later the PROMs came back with the explanation that someone must have forgotten to plug in a master when one of the PROMs had been programmed. I plugged in the chips again, turned on the power supply and lo and behold I had a blinking cursor and an otherwise blank screen. Well, an almost blank screen anyway. Several nonblinking cursor blocks would not clear from the screen but would only move around. I decided to ignore these for a while as they seemed to be only an annoyance and nothing disastrous.


Photo 2: Here are the three main boards of the Sphere system. At the upper left is the video interface board with 512 bytes of memory used for the display. At upper right is the main processor board with the $6800,4 K B$ of dynamic memory and $1 K B$ of EROM. The lower board is the keyboard assembly, which includes a numerical keypad and cursor control keys as well as a full alphanumeric keyboard.

A short while later 1 had the editor program working as well as the debug program for opening and changing the contents of memory locations. Unfortunately there was little documentation about the "mini assembler" and I could not figure out how to use this feature based on the original documentation package. So after a few days of inputting simple programs through the debug routine, I went back to trying to fix the TV display.

By this time the problem had become a little more than annoying. On four of the 16 lines of display, several blocks would not clear as noted above, but far worse, if the cursor was placed on any of these lines, more blocks would be created randomly in the two or three positions following the cursor. On four other lines several characters would not be displayed as the correct symbol. By swapping memory ICs, I determined that the problem was in the M6810 chips used for character storage in the TV display. Rather than waiting two weeks for replacements from Sphere, I bought two new ones from an outside vendor. One of these also had a minor defect, so 1 was running a $50 \%$ failure rate on these chips.

1 decided to call Sphere and find out about the "mini assembler." It turned out this call was unnecessary since a manual update covering the "mini assembler" and

Photo 3: It works! Here is my completed Sphere system in operation driving a small video monitor (not purchased from Sphere). The board interconnections are made with flat ribbon cables, and the keyboard is shown mounted in its case.

including a memory map had been mailed the day before I called. I also asked about the problems with the character memories and was told that actually the M6810 was not specified to be fast enough for the application but that because Motorola generally made the chips better than spec, they had decided to use them instead of the slightly more expensive M6810-1 which was guaranteed to be fast enough. Personally, I would have preferred to pay the extra cost (<\$5 extra for the four chips according to Motorola price lists) and have a display that worked the first time.

Finally the manual addition arrived and I could use the "mini assembler" and could start writing programs seriously. Also with the manual update was a notice about the serial interface board which still had not arrived. Shortly after Sphere started shipments, a conference sponsored by BYTE was held to try to determine some standards for cassette interfaces. Sphere then stopped shipment of all their cassette boards so they could redesign the board so it met the standards which were decided on. This meant another wait of about two months, but the eventual cassette interface would meet the provisional industry standards and even included the option of interfacing to two cassettes which the original design did not offer.

Probably the worst fault I found with my Sphere system is its lack of documentation. With the manual addition received in December, a start was made towards correcting this fault. For example, assembly instructions
should include some sort of debugging suggestions. The manual 1 received provided little if any advice on debugging a nonfunctioning system other than: Look for solder splashes. Solder splashes do occur, but other faults (like misprogrammed EROMs) also occur and the method of detecting these faults should not be left entirely up to the customer.

When all things are considered, I think Sphere has done a very good job in the short time the company has been in existence. The hardware supplied in the kits is generally of good quality. In the few cases where the hardware is not up to par, they are working to improve it and usually will retrofit kits already shipped when they do. The problems and delays I faced should no longer face future buyers since the Sphere people are hard at work ironing out bugs and glitches in the product and its documentation.

The programs in the EROM provide an incredible amount of flexibility, especially considering that everything is stuffed into only 1 KB of memory; and certainly the keyboard, TV interface and EROM combination make program entry an order of magnitude simpler than the switch and LED front panel method.

Computer systems can get quite complex. I extend my praise to Sphere for getting all the myriad pieces together in such a short time. If a company which has been in business less than a year has no faults greater than these, it will quite likely become one of the leaders of the industry in a very short period of time.

## the PLEASER...


...we disigned our computers for people like you who are interested in having one, but up 'til now didn't have $\$ 30$ or $\$ 40$ thousand to put into one. It is all together in one package now, thanks to advanced technology and the good ole american spirit of doing it faster, better, and cheaper than the other guy.

Because we are a young company (we just celebrated our first birthday) we have a lot of that "spirit" and some pretty capable men who came up with a computer that others are just now beginning to copy. While copying us is a real compliment, no one has yet equalled or bettered what we have developed.

We have a Central Processing Unit (CPU) that contains the most advanced technology available today. That means to you that the CPU will do what a CPU ought to do. It's reliable, and THAT'S A PLEASER.

We have a Serial Interface Module (SIM) that is a very powerful idea in the computer world today. Let me explain. You can have an inexpensive audio cassette player/recorder as a mass-memory storage device. Understand the power of the SIM board. For just pennies you can store vast amounts of Data on an ordinary audio cassette. NOW THAT'S A PLEASER Sphere also offers the floppy disc as an option when time gets to be a premium.

We have a Program Development System (PDS) that is an aid to programming. What an idea! It remembers its job even when the computer is off. With instant call-up of the program outline when the power is turned on. There is an Editor, an Assembler, and a Debugger to help you do your programming correctly, and quickly, leaving you a lot more time to actually use the computer as a data handler... and THAT'S A PLEASER.

As if that weren't enough ...we have a newly added feature THE 4 K ROM BD. A board containing up to 4000 bytes of Read Only Memory. (The Stuff that's on it doesn't go away when the computer is off.) You put your most used applications and computer games on the ROM board Turn on the power, locate your favorite game, and you're up and hummin immediately. This makes using the computer lots of fun for a change. WHAT A PLEASER.

This is what you get. A computer that does what you want it to do when you want it to do it, at a price that you can finally aford. NOW, THAT'S REALLY A PLEASER!!!
from SPHERE
the COMPUTER PEOPLE who have it ALL TOGETHER.
P.S. There's much much more, so write today for prices and information.

EPHER

## 940 North 400 East

 North Salt Lake, Utah 84054 (801) 292-8466

Photo 1: In this example, Educator-8080 is shown before (left) and after (right) the execution of an $X R A$ A instruction. The effect of this instruction is of course to clear the accumulator $A$, as is shown at the right.

## Explore an 8080

## with Educator-8080

Charles P Howerton<br>President, Digital Group Software Systems Inc PO Box 1086<br>Arvada CO 80001

## What Is an Educator?

Educator-8080 was designed as a classroom instruction aid for a microprocessor programming course. The principal design goals were to develop a system which would illustrate the architecture of the machine and the effect of the execution of various instructions. For example, the reader might ask to what use the logical EXCLUSIVE OR function may be put in an 8080. This function, which operates on each bit, has a value of 1 if either of the two operands or arguments, but not both, has a value of 1 ; otherwise it has a value of 0 . The Edu-cator-8080 can simply illustrate this function. In the example shown in photo 1 (left and right), both arguments of this function are equal: the first argument is the value in the accumulator (A) and the second argument is also the value stored in the accumulator (A). The function value (ie: the result) is placed in the accumulator after
exccution. The left photo shows the accumulator (and other registers, etc, which are not affected) before execution of the instruction XRA $A$, and the right photo shows it after execution. The result is that the accumulator (A) is cleared, ie: it contains eight 0 bits. This result is consistent with the definition of the EXCLUSIVE OR above: Whenever both bit arguments are 0 , or are 1 , a value of 0 is returned. This example shows that the Educator-8080 is a convenient means to illustrate rather complex operations which facilitate learning the instruction set and architecture without the tedium of plowing through books. A subordinate goal was to implement the entire system with the exception of the physical input output routines and the stack in 1024 bytes of memory. All of the design goals were met. In addition, if the 10 devices are ASCII oriented, a reduction in the length of the error messages (perhaps limiting them to the error code number) should provide sufficient space for the inclusion of the physical input output routines and the stack within the 1 K byte memory space.

Educator-8080 is written in a fairly straightforward manner and it should not be particularly difficult to adapt it to any 8080 system with more than 1 K bytes of pro-
grammable memory, a keyboard, and an output device of some kind. It is designed to operate with a television display device and to dynamically show the results of the execution of the input commands.

It would probably be desirable to modify the display function somewhat if the output device is a hard copy device such as a Teletype. The content could be the same but the elimination of blank lines and printing the titles only every 10 or 15 instructions would speed things up considerably on a Teletype device. The input output routines required to adapt Educator- 8080 to almost any system are described functionally but are not given in detail. They should be adapted from routines already in use for a given system.

## The Instruction Set

The Educator-8080 instruction set is a subset of the 8080 instruction set. The commands implemented within Edu-cator- 8080 were selected to provide representative instructions from most of the functional instruction groups. Since the instructions are to be executed one at a time from keyboard input, there was no need to incorporate any of the Jump, Call or Return instructions; however, since the flags are displayed after each operation, it is very easy to determine whether or not a given conditional Jump, Call or Return would cause a transfer of control by simply observing the setting for the flag whose status is being tested. In addition, because of memory limitations none of the instructions which cause memory to be read or written were implemented. Finally, no instruction whose action could not be readily observed was implemented.

To keep the display as uncluttered as possible the registers which could be accessed by Educator-8080 instructions were limited to the accumulator and the B and C registers. It would not be particularly difficult to incorporate the rest of the registers into the display and as operands for the Educator-8080 instruction subset. However, unless the ability to address memory is desirable the only instructions which could be added to the subset would be the DAD and the XCHG.

The instruction subset and the valid operands for each instruction are shown in table 1. Table 2 contains the corresponding information as it is loaded into the computer's memory and used by Educator- 8080 .

## Immediate Operands

Almost half of the instructions supported by Educator-8080 require immediate

Table 1: Command List for Educator-8080. In order to illustrate the operations of the 8080 processor, Educator- 8080 interprets a subset of the 8080's instructions. The subset generally references the accumulator, $A$, and registers $B$ or $C$; it excludes all branching and program control operations. The complete list of available operations is found in this table.



Photo 2: What happens when an 8080 executes an $A D D \quad B$ instruction? $A$ specific example is itlustrated in this set of before and after snapshots.
operands. An immediate operand is a constant value which is part of the instruction being executed and it immediately follows the operation code of the instruction, hence the name immediate.

Whenever a single byte "constant" is required in a program, its inclusion as the immediate value of an appropriate instruction reduces the length of the program because there is no need to address the value

Table 2: Operation Code Table for the Educator-8080 program. Table 1 showed the command list for the program. This table gives the absolute machine codes for the command table beginning at address <2>/122. Each command is represented by a 3 byte ASC/I character string mnemonic followed by the naked (without register values) 8080 operation code and the address of the routine which interprets the command. The routine name is shown symbolically in the right hand column, and can be found in the program of listing 1.

| Address | ASCII <br> Mnemonic |  |  |  |  |  | Routine Name |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Mne | monic |  | Opcode | Routine |  |
| $<2>/ 122$ | 'ACI' | 101 | 103 | 111 | 316 | 212<1> | IMMED |
| $<2>1130$ | 'ADC' | 101 | 104 | 103 | 210 | 152<1> | RG210 |
| $<2>/ 136$ | 'ADD' | 101 | 104 | 104 | 200 | $152<1>$ | RG210 |
| $<2>/ 144$ | 'ADI' | 101 | 104 | 111 | 306 | 212<1> | IMMED |
| $<2>/ 152$ | 'ANA' | 101 | 116 | 101 | 240 | $152<1>$ | RG210 |
| $<2>/ 160$ | 'ANI' | 101 | 116 | 111 | 346 | 212<1> | IMMED |
| <2>/166 | 'CMA' | 103 | 115 | 101 | 057 | 144<1> | DIRCT |
| $<2>/ 174$ | ${ }^{\text {'CMC' }}$ | 103 | 115 | 103 | 077 | $144<1>$ | DIRCT |
| $<2>/ 202$ | 'CMP' | 103 | 115 | 120 | 270 | $152<1>$ | RG210 |
| $<2>1210$ | ${ }^{\text {CPPI' }}$ | 103 | 120 | 111 | 376 | $212<1>$ | IMMED |
| <2>/216 | 'DAA' | 104 | 101 | 101 | 047 | 144<1> | DIRCT |
| $<2>/ 224$ | 'DCR' | 104 | 103 | 122 | 005 | $245<1>$ | RG543 |
| $<2>/ 232$ | 'DCX' | 104 | 103 | 130 | 013 | 264<1> | RG54B |
| $<2>1240$ | 'INR' | 111 | 116 | 122 | 004 | 245<1> | RG543 |
| $<2>/ 246$ | 'INX' | 111 | 116 | 130 | 003 | 264<1> | RG54B |
| $<2>/ 254$ | 'MOV' | 115 | 117 | 126 | 100 | $145<1>$ | MOVRT |
| $<2>/ 262$ | 'MVI' | 115 | 126 | 111 | 006 | 205<1> | MVIRT |
| $<2>1270$ | 'NOP' | 116 | 117 | 120 | 000 | $144<1>$ | DIRCT |
| $<2>/ 276$ | 'ORA' | 117 | 122 | 101 | 260 | 152<1> | RG210 |
| <2>/304 | 'ORI' | 117 | 122 | 111 | 366 | 212<1> | IMMED |
| $<2>/ 312$ | 'RAL' | 122 | 101 | 114 | 027 | $144<1>$ | DIRCT |
| $<2>/ 320$ | 'RAR' | 122 | 101 | 122 | 037 | $144<1>$ | DIRCT |
| $<2>/ 326$ | 'RLC' | 122 | 114 | 103 | 007 | $144<1>$ | DIRCT |
| <2>/334 | 'RRC' | 122 | 122 | 103 | 017 | $144<1>$ | DIRCT |
| <2>/342 | 'SBB' | 123 | 102 | 102 | 230 | 152<1> | RG210 |
| $<2>/ 350$ | 'SBI' | 123 | 102 | 111 | 336 | 212<1> | IMMED |
| <2>/356 | 'STC' | 123 | 124 | 103 | 067 | 144<1> | DIRCT |
| <2>/364 | 'SUB' | 123 | 125 | 102 | 220 | 152<1> | RG210 |
| $<2>/ 372$ | 'SUI' | 123 | 125 | 111 | 326 | $212<1>$ | IMMED |
| $<3>1000$ | 'XRA' | 130 | 122 | 101 | 250 | 152<1> | RG210 |
| $<3>1006$ | 'XRI' | 130 | 122 | 111 | 356 | 212<1> | IMMED |

directly. Immediate values have an implied address which is the address of the byte following the opcode and this address is supplied from the program counter register automatically whenever an immediate type instruction is executed. In the Edu-cator-8080 system the "program counter" is provided by the operator's sequence of commands which are executed one by one.

Educator-8080 has three different types of immediate values as part of the input command and defaults to one of these types if the input command omits type information.

The general form of an immediate operand is as follows:
$T p \vee p$
Where: $T$ is the type code which designates the form of the immediate value and may be any of the following:
$B$ - for a binary immediate value
Q - for an octal immediate value
H - for a hexadecimal immediate value
If the type code is omitted entirely and the first nonpunctuation character encountered is a numeric digit 0 to 7 , then a default type of octal is assumed.
$p$ is any form of punctuation (eg: single or double quotes, parentheses, etc). Punctuation is not required, and provision for its inclusion is solely in the interest of enabling the user to enter commands in a format consistent with that of various advanced assemblers.
V is the value of the immediate operand expressed in a form consistent with the explicit or implied type selected. The form and content of the value field for each type is as follows: $\mathrm{T}=\mathrm{B}: \mathrm{V}$ is a series of eight consecutive numeric characters which have the value zero or one.
Example: B'11000111'. $V$ is 11000111, quotes are optional. $\mathrm{T}=\mathrm{Q}$ or T omitted: V is a series of

Photo 3: To illustrate the use of hexadecimal immediate values, this photo shows the operation of XRI H'C3'.


Listing 1: The Educator-8080 program expressed as an absolute assembly language listing. The notations $\langle 0\rangle,\langle 1\rangle,\langle 2\rangle$ and $\langle 3\rangle$ are used to denote the high order (page) address bytes of four consecutive pages in memory address space. When loading the program into a given system, these notations become bytes with consecutive octal values. Thus to load the program ot location 200/000 in memory address space, the values utilized would be 200, 201, 202 and 203.
three consecutive numeric characters which have octal digit values of from 0 to 7.
Example: $Q^{\prime} 307^{\prime} \mathrm{V}$ is 307 , quotes are optional.
$\mathrm{T}=\mathrm{H}: \mathrm{V}$ is a pair of consecutive characters which have hexadecimal digit values from 0 to $F$.
Example: H'C7' $V$ is C 7 , quotes are optional.

With the exception of the move immediate (MVI) command which requires a destination register, immediate commands are entered as the mnemonic opcode followed by the immediate operand in any of its valid forms.

Some "before and after" examples of Educator-8080 commands are shown in photos 1 through 3. In each case, a command is typed into the keyboard of the computer, then the Educator- 8080 display following the command is depicted.

## Entering Commands

Commands are entered into Educator-8080 as a string of characters (eg: letters, numbers, spaces and punctuation) followed by a command termination character. As written, Educator-8080 assumes that the command termination character will be an ASCII carriage-return (octal 015). However, any other keyboard character code may be used as the command termination character by changing the value of the immediate operand in the instruction located at address $<0>/ 341$ which tests for command termination. (See listing 1.)

Since it is not uncommon to make errors when keying information into a computer, two provisions have been made in Edu-cator-8080 for correcting or eliminating

| ddress |  | tal-co |  | labal | op | operand | commentery |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| -The control routine is the top of the structure and controls the operation of the entire program. |  |  |  |  |  |  |  |
| < $0>1000$ | 061 | xxx | xxx | CNTRL | LXI | SP, STACK | Set stack pointer to programmable memory: |
| <0>1003 | 315 | 026 | <0> | NOTZER | CALL | DSPLY | Display contents of registers; |
| < $0>1006$ | 315 | 316 | <0> |  | CALL | CMDNT | Enter a command; |
| $<0>1011$ | 315 | 063 | <1> |  | CALL | FETCH | Fetch the correct opcode; |
| <0>1014 | 267 |  |  |  | ORA | A | Set zaro flag as per contents: |
| <0>1015 | 302 | 003 | <0> |  | JNZ | NOTZER | Jump if not zero error occurred: |
| <0>1020 | 315 | 030 | <2> |  | CALL | XOTER | Go execute the current command; |
| <0>/023 | 303 | 003 | <0> |  | JMP | NOTZER | Loop forever; |
| -This display routine controls the generation of the dynamic display. |  |  |  |  |  |  |  |
| <0>1026 | 041 | 167 | <3> | DSPLY | LXI | H,TITLS | Losd address of titles into HL: |
| <0>1031 | 315 | 261 | <0> |  | CALL | CHEDT | Displey tities; |
| $<0>1034$ | 041 | 257 | <3> |  | LXI | H,BLINE | Losd addr of BLINE title; |
| $<0>1037$ | 072 | 351 | <3> |  | LDA | BREG | Losd contents of BREG into $A$ : |
| $<0>1042$ | 315 | 132 | <0> |  | CALL | DSPCV | Convert and display: |
| $<0>1045$ | 041 | 271 | <3> |  | LXI | h,Cline | Losd addr of CLINE title; |
| $<0>1050$ | 072 | 350 | <3> |  | LDA | CREG | Losd contents of CREG into A; |
| $<0>1053$ | 315 | 132 | <0> |  | CALL | DSPCV | Convert and display: |
| $<0>1056$ | 041 | 304 | <3> |  | LXI | H,AFHDR | Loed addr of A flogs title; |
| <0>1061 | 315 | 261 | <0> |  | CALL | CHEDT | Display titles; |
| $<0>1064$ | 052 | 346 | <3> |  | LHLD | PSWA | Load flags and A into HL; |
| <0>/067 | 175 |  |  |  | MOV | A,L | Move flags to A; |
| <0>1070 | 346 | 004 |  |  | ANI | $\mathrm{B}^{\prime \prime} 0000010{ }^{\prime \prime}$ | AND off all but parity flag; |
| <0>/072 | 315 | 237 | <0> |  | CALL | DSPFG | Display the flag value: |
| <0>1075 | 175 |  |  |  | MOV | A,L | Move flass to A; |
| <0>/076 | 346 | 100 |  |  | ANI | B01000000' | AND off all but zero fieg; |
| $<0>/ 100$ $<0>/ 103$ | 315 | 237 | <0> |  | CALL | DSPFG | Oisplay the flag value; |
| $<0>/ 103$ $<0>/ 104$ | 175 |  |  |  | MOV | A, | Move flags to A; |
| $<0>1104$ $<0>1106$ | 346 | 200 |  |  | ANI | B'10000000' | AND off all but sign flag: |
| $<0>1106$ $<0>111$ | 315 | 237 | <0> |  | CALL | DSPFG | Display the filag value: |
| $<0>1111$ $<0>1112$ | 175 |  |  |  | MOV | A,L | Move flags to $A_{\text {i }}$ |
| $<0>1112$ | 346 | 020 |  |  | ANI | 8'00010000' | AND off all but auxiliary carry flag: |
| <0>/114 | 315 | 237 | <0> |  | CALL | DSPFG | Display the flag value: |
| <0>/117 | 175 |  |  |  | MOV | A, L | Move flags to $A$; |
| $<0>1120$ | 346 | 001 |  |  | ANI | 8'00000001' | AND off all but carry flag: |
| <0>/122 | 315 | 237 | <0> |  | CALL | DSPFG | Display the flag value: |
| <0>/125 | 174 |  |  |  | MOV | A, H | Move A register value to A; |
| <0>/126 | 315 | 137 | <0> |  | CALL | DSPCN | Display with no title print: |
| <0>1131 | 311 |  |  |  | RET |  | Return to the CNTRL routine; |
| *The display conversion routine prints binary, octal end hexadecimal. |  |  |  |  |  |  |  |
| <0>/132 | 365 |  |  | OSPCV | PUSH | PSW | Save output value for CHEDT: |
| < $<0\rangle / 133$ | 315 | 261 | <0> |  | CALL | CHEDT | Disploy line title addr in HL; |
| < $0>/ 136$ | 361 |  |  |  | POP | PSW | Retrieve seved output value; |
| <0>/137 | 036 | 010 |  | DSPCN | MVI | E,0010' | - Move 8 to E register; |
| $<0>1141$ $<0>/ 142$ | 007 |  |  | DSPBT | RLC |  | Rotate MSB into Carry and LSB; |
| $<0>/ 142$ $<0\rangle / 143$ | 365 |  |  |  | PUSH | PSW | Save current value; |
| $<0>/ 143$ $<0>/ 145$ | 346 | 001 |  |  | ANI | Q'001' | AND off ell but LSE; |
| <0>1145 | 315 | 237 | <0> |  | CALL | DSPFG | Go display bit value: |
| <0>1150 | 361 |  |  |  | POP | PSW | Retrieve saved current value; |
| <0>/151 | 035 |  |  |  | DCR |  | Decrement loop count; |
| <0>/152 | 302 | 141 | <0> |  | JNZ | DSPET | Jump if loop count not zero; |
| <0>/155 | 267 |  |  |  | ORA |  | Reset carry: |
| $<0>1156$ | 036 | 003 |  |  | MVI | E, $\mathbf{Q ' O O S}^{\prime}$ | Move 3 to E register: |
| <0>1160 | 027 |  |  | DSPOT | RAL |  | MSB to Carry, Carry to LSB, |
| <0>/161 | 027 |  |  |  | RAL |  | do it again. |
| <0>/162 | 027 |  |  |  | RAL |  | three times for octal digit thift; |
| $\leq 0>163$ | 365 |  |  |  | PUSH | PSW | Save current value; |
| < $0>1164$ | 346 | 007 |  |  | ANI | 9007' | AND off all but octal LSD: |
| $<0>1166$ | 366 | 060 |  |  | ORI | $\mathrm{Q}^{\prime} 060{ }^{\text {c }}$ | OR on bits to make ASCII numeric character; |
| < $0>1170$ | 315 | x $\times$ x | xxx |  | CALL | CHRPR | Output the character; |
| < $0>1173$ | 361 |  |  |  | POP | PSW | Retrieve saved current value: |
| <0>1174 | 035 |  |  |  | DCR | E | Decrement loop count; |
| <0>1775 | 302 | 160 | <0> |  | JNZ | DSPQT | Jump if loop count not zero; |
| <0>1200 | 315 | 251 | <0> |  | CALL | DSPSP | Output a space; |
| $<0>/ 203$ $<0>1205$ | 036 007 | 002 |  |  | MVI | E.O'002' | Move 2 to E; |
| $\begin{aligned} & <0\rangle / 205 \\ & <0>/ 206 \end{aligned}$ | 007 |  |  | DSPHT | RLC |  | Rotate MSB into Carry and LSB, |
| $<0>1206$ $<0\rangle 1207$ | 007 |  |  |  | RLC |  | do it again. |
| <0>/207 | 007 |  |  |  | RLC |  | four times for. |
| <0> 1210 | 007 |  |  |  | RLC |  | hexadecimal shift; |
| <0>/211 | 365 |  |  |  | PUSH | PSW | Save current value; |
| <0>/212 | 346 | 017 |  |  | ANI | $B^{\prime} 00001111^{\prime}$ | AND off all but hexadecimal LSD; |
| $<0>1214$ | 306 | 060 |  |  | ADI | - $0600^{\prime}$ | Add on bits to make ASCII numeric character: |
| $<0>1216$ | 376 | 072 |  |  | CPI | 9'072' | Compare resulr to one more than 9; |
| <0>1220 | 332 | 225 | <0> |  | JC | DSPHS | If numeric then skip adjustment; |
| <0>/223 | 306 | 007 |  |  | ADI | $0^{\prime} 007{ }^{\prime}$ | Add 7 giving ASCII ' $A$ ' thru ' $F$ ' codes; |

Listing 1, continued:

|  | eddrass | ectal-code |  |  | labal | op | operand | commentary |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | <0>1225 | 315 | xxx | xxx | DSPHS | CALL | CHRPR | Output the character: |
|  | <0>1230 | 361 |  |  |  | POP |  | Retrieve saved current value; |
|  | <0>1231 | 000 |  |  |  | NOP |  |  |
|  | <0>1232 | 035 |  |  |  | OCR |  | Decrement loop count; |
|  | <0>1233 | 302 | 205 | <0> |  | JNZ | OSPHT | Jump if loop count not zero; |
|  | <0>/236 | 311 |  |  |  | RET |  | Return to calling routine: |
|  | - Display flag or binary digit followed by a space. Alternate entry is used to display a space. |  |  |  |  |  |  |  |
|  | <0>1237 | 312 | 244 | <0> | DSPFG | JZ | DSPFZ | Jump if passed value is a zero; |
|  | <0>/242 | 076 | 001 |  |  | MVI | A, O.001' | Otherwise move a 1 into $\mathrm{A}_{\text {: }}$ : |
|  | <0>/244 | 306 | 060 |  | DSPFZ | ADI | $\bigcirc$ | Convert into ASCII numeric character; |
| $\rightarrow$ | <0>1246 | 315 | xxx | xxx |  | CALL | CHRPR | Output the charsctier: |
|  | <0>1251 | 365 |  |  | DSPSP | PUSH |  | Save the flags and value in $A$ : |
|  | <0>/252 | 076 | 040 |  |  | MVI | A. $0^{\prime} 040{ }^{\prime}$ | Move space into $A$ : |
|  | <0>1254 | 315 |  | xxx |  | CALL | CHRPR | Output the space; |
|  | - 0 < ${ }^{\text {c/257 }}$ | 361 311 |  |  |  | POP | PSW | Retrieve the saved flags and $A$ : |
|  | <0>/260 | 311 |  |  |  | RET |  | Return to the calling routine; |
|  | - The character string output edit routine. |  |  |  |  |  |  |  |
|  | <0>/261 | 176 |  |  | CHEDT | mov | A, M | Move next charscter into A: |
|  | <0>/262 | 376 | 200 |  |  | CPI | $0^{\prime 2} 200^{\prime}$ | Compare it to 200 octal: |
|  | <0>1264 | 310 |  |  |  | RZ |  | Return if equal it's end of string; |
| - | <0>/265 | 322 | 277 | <0> |  | JNC | CHSPA | Jump it greater for space routine: |
|  | <0>1270 | 315 | *x* | *x |  | CALL | CHRPR | Else go output the character; |
|  | <0>1274 | 303 | 261 | <0> | CHEND | JMP | CHEDT | Increment the string index; Looo for nexi character: |
|  | <0>/277 | 326 | 200 |  | CHSPA | Sul | $\square^{\circ} 200^{\prime}$ | Subtrect 200 octal from value: |
|  | <0>/301 | 107 |  |  |  | MOV |  | Move space count to $B$; |
|  | <0>1302 | 076 | 040 |  | CHSPL | MVI | A, $0^{\circ} 040^{\prime}$ | Move space to A ; |
| $\rightarrow$ | <0>/304 | 315 | xxx | xx |  | CALL | CHRPR | Output the space; |
|  | <0>/307 | 005 |  |  |  | DCA |  | Decrement space count; |
|  | $<0>1310$ | 302 | 302 | <0> |  | JNZ | CHSPL | Jump it count not zero to start of loop: |
|  | $<0>1313$ | 303 |  | <0> |  | JMP | CHEND | Jump back into CHEDT loop; |
|  | -The command entry routine eccepts input from the keyboard for comriands. |  |  |  |  |  |  |  |
|  | <0>/316 | 041 | 332 | <3> | CMDNT | $\left\llcorner_{\text {L }}{ }_{1}\right.$ | H,CMDMS | Move address of 'COMMAND ? ' to HL: |
|  | <0>/321 | 315 | 261 | <0> |  | GALL | CHEDT | Display the messbge: |
|  | <0>/324 | 041 | 352 | <3> |  | Lxı | H,CMDAR | Move address of command input area HL: |
|  | <0>1327 | 006 | 026 |  |  | MVI |  | Move maximum length to B ; |
| - | <0>/331 | 315 | x×x | xxx | CMDKB | CALL | KEYBD | Get an input character; |
|  | <0>/334 | 376 | 014 |  |  | CPI | $0.014^{\prime}$ | Is it a control- 1 line delate? |
|  | $<0>1336$ | 312 | 000 | <0> |  | 12 | CNTRL | If so then restart program; |
|  | <0>/341 | 376 | 015 |  |  | CPI | $\square^{\circ} 015{ }^{\prime}$ | Is it a carriage return? |
|  | <0>/343 | 312 | 376 | <0> |  | 12 | CMDNO | If so then go compress input; |
|  | <0>/346 | 376 | 177 |  |  | CPI | Q'177 | Is it a delete character? |
|  | <0>1350 | 302 | 355 | <0> |  | JNZ | CMDST | 14 not then go store the character; |
|  | <0>/353 | 076 | 033 |  |  | MVI | A, $\mathrm{Q}^{\prime} 033^{\prime}$ | If so reploce with back arrow: |
|  | <0>/355 | 167 |  |  | CMDST | mov | M, A | Store input charecter in command buffer: |
|  | <0>/356 | 315 | xxx | xxx |  | CALL | CHRPR | Display the input character; |
|  | <0>/361 | 043 |  |  |  | INX | H | Increment command work area index; |
|  | <0>/362 | 005 |  |  |  | DCR | 日 | Decrement command length count; |
|  | <0>/363 | 302 | 331 | <0> |  | JNZ | CMDKB | If not full then reiterate: |
|  | <0>/366 | 076 315 | 001 |  |  | MVI | A, ${ }^{\text {coob }}$ - | If buffer full then select error |
|  | < $0>/ 370$ $<0>/ 373$ | 315 303 | 063 000 | $\begin{aligned} & \langle 2\rangle \\ & \langle 0\rangle \end{aligned}$ |  | ${ }_{\text {chap }}^{\text {CALL }}$ | ERROR | number 1 and print its message: |
|  | <0>/373 | 303 | 000 | <0> |  | JMP | CNTRL | Restart the program; |
|  | -The command compress routine oliminates all but letters and numbers. |  |  |  |  |  |  |  |
|  | <0>/376 | 041 | 352 | <3> | CMDND | LXI | H,CMDAR | Losd HL with eddress of work erea: |
|  | $\langle 1>1001$ | 345 |  |  |  | PUSH |  | Push \& pop move it 10 DE |
|  | $<1>1002$ | 321 | 026 |  |  | POP |  | at the compremion pointer: |
|  | <1>/003 | 076 |  |  |  | MVI | A, $\mathrm{O}^{\prime} 026^{\prime}$ | Losd A with maximum length; |
|  | $<1>1005$ | 220 |  |  |  | SUB |  | Subtract remsining length from $\mathbf{B}$; |
|  | <1>1006 | 107 |  |  | CMDNX | MOV | B.A | Move octual length to B; |
|  | $<1>1007$ $<1>1010$ | 176 376 |  |  |  | MOV |  | Move command charecter to A; Is it a back arrow (charecter delete)? |
|  | <1>1012 | 302 | 027 | <1> |  | JNZ | CMDCH | If not then go to other tests: |
|  | <1>/015 | 076 | 352 |  |  | MVI | A,CMDAR-L | Low addrass byte of CMDAR to A: |
|  | <1>/017 | 273 |  |  |  | CMP |  | Compere to currant low addras byte; |
|  | <1>/020 | 322 | 055 | <1> |  | JNC | CMDNS | If not greater then skip rava: |
|  | $<1>1023$ $<1>1024$ | 033 303 |  | <1> |  | OCX | ${ }_{\text {CMDNS }}$ | Else bock up compression pointer: |
|  | <1>1027 | 376 | 055 060 |  | CMDCH | CPI | O ${ }^{\text {'060' }}$ | Is the charecter less than 00 |
|  | <1>/031 | 332 | 055 | <1> |  | JC | CMDNS | If so then skip saving it; |
|  | <1>1034 | 376 | 072 |  |  | CPI | Q'072' | is the character less than ' 9 ' +1 '? |
|  | <1>1036 | 332 | 053 | <1> |  | JC | cmbsv | If so then save numeric vatue; |
|  | <1>/041 | 376 | 101 |  |  | CPI | Q'101' | Is the character less than ' $A$ '? |
|  | $<1>1043$ $<1>1046$ | 332 376 | 055 133 | <1> |  | ${ }_{\text {J }} \mathrm{Cl}$ | CMDNS |  |
|  | $<1>1046$ $<1>1050$ | 376 | 055 | <1> |  | CP1 | $\mathrm{Q}^{1} 13{ }^{\text {a }}$ | Is the character greater than ' 2 '? |
|  | <1>1053 | 022 |  |  | cmbsv | STAX |  | If so then skip saving it: |
|  | <1>1054 | 023 |  |  | CMDS | INX | D | Incrememt compression pointer index; |
|  | $<1>1055$ $<1>1056$ | 043 |  |  | CmDNS | INX | ${ }_{\text {H }}$ | Increment input erting pointer; |
|  | <1>1056 | 005 | 007 <1> |  |  | DCA | ${ }_{\text {B }}^{\text {B }}$ | Decrement sctual length count; |
|  | <1>1062 | 311 |  |  |  | RET | Cmdnx | Else retum to CNTRL calling point; |


| <1>/063 | 041 | 122 | <2> | FETCH | LXI | H,OPTAB | Loed eddress of opcode table HL: |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $<1>1066$ | 036 | 037 |  |  | MVI | E,0037 | Move table element count to E : |
| <1>/070 | 345 |  |  | FLOOP | PUSH | H | Save currant element sddress; |
| $<1>/ 071$ | 001 | 352 | <3> |  | LXI | B,CMDAR | Load Bddress of CMDAR into BC: |
| <1>/074 | 026 | 003 |  |  | MVI | D.0003' | Move opcode length to D; |
| <1>/076 | 012 |  |  | FCOMP | LDAX | B | Load command character to A indexed by B; |
| <1>1077 | 276 |  |  |  | CMP | M | Compare it to table charscter; |
| <1>/100 | 302 | 125 | <1> |  | JNZ | FNXEL | If not equal then go to next element; |
| $<1>/ 103$ | 003 |  |  |  | INX | B | Increment command charscter index; |
| $<1>/ 104$ | 043 |  |  |  | INX | H | Increment table character index; |
| $<1>1105$ | 025 |  |  |  | DCR | D | Decremant opcode length counter; |
| $<1>1106$ | 302 | 076 | <1> |  | JNZ | FCOMP | If not zero continue test loop; |
| $<1>1111$ | 343 |  |  |  | XTHL |  | Exchenge HL with top of steck: |
| $<1>/ 112$ | 341 |  |  |  | POP | H | Pop HL from stack to clear it; |
| $<1>1113$ | 136 |  |  |  | MOV | E, M | Move naked opcode to E, D is zero; |
| <1>/114 | 325 |  |  |  | PUSH | D | Save naked opcode; |
| $<1>115$ | 043 |  |  |  | INX | H | Increment table pointer; |
| $<1>/ 116$ | 136 |  |  |  | MOV | E,M | Decode routine low address byte to E: |
| $<1>1117$ | 043 |  |  |  | INX | H | Incrament table pointer: |
| <1>120 | 126 |  |  |  | MOV | D,M | Decode routine high address byte to D; |
| $<1>1121$ | 353 |  |  |  | XCHG |  | Move decode routine eddress to HL ; |

errors. The ASCII delete character code (octal 177) is used to delete the last remaining character in the input string. Since a deleted character is not considered to exist, N consecutive delete characters will delete the N preceding characters. For example, if the delete character is shown as a back arrow $(\leftarrow)$, RAX $\leftarrow L$ will be reduced to RAL and CQP $\leftarrow M A$ will be reduced to CMA. Characters which have been keyed in are displayed after they have been tested. The display function uses the octal value 177 as a clear screen control code; therefore, character deletes are transformed into the back arrow before they are displayed and stored. Edu-cator-8080 users with systems which have a back arrow (octal 033) key on their keyboards may use it as a character delete code and it will have the same effect as the delete key assumed in this version. Users who have neither of these keys can designate any keyboard character as the delete character code by changing the immediate operand in the instruction located at $<0>/ 346$ which tests for the delete character. (See listing 1). The other, and somewhat more drastic, method of eliminating keying errors is to delete the entire input line. This is usually done when an error is detected before the command termination character is input but several characters after the error occurred. The procedure for deleting an entire line is to enter an ASCII form feed code (octal 014) which is a "control L." combination on typical ASCII keyboards. This will clear the input line and restart the command entry procedure. Like the command termination and the character delete codes, the line delete code can be made to be any keyboard character by changing the value of the immediate operand in the instruction at location $<0>/ 334$ which tests for the line delete code.

A very useful feature of Educator-8080 permits the user to execute the last command input several times. This is accomplished by simply keying the command termination character when the system calls for the entry of a new command. In order to provide this facility the input buffer is not cleared prior to calling for the entry of a new command, so the last previously entered command is still in the buffer. This feature is especially handy when demonstrating the effect of multiple executions of the rotate, increment, decrement, arithmetic and logical commands.

The general format for entering a command is as follows:

## OPCODE[p OPERAND-1[p OPERAND-2]]t

Where:
OPCODE is the mnemonic opcode for the
command. For example; MOV XRI, etc.
p
is any desired form of punctuation or a space. $p$ is not required and, therefore, may be omitted entirely.
OPERAND-1 is the first or only operand re quired by an instruction. It may be a register identification or an immediate value. See table 1 for the operand requirements.

OPERAND-2 is the second operand where required by a specific instruction. See table 1.
$t$ is the command termination character, an ASCII carriage return in the listing 1 version of Educator 8080.

The brackets ([ ]) shown in the general format are used to indicate that the items within them are optional, since some commands do not require any operands (eg: RAL, STC, CMA, etc), some require one operand only (eg: ADI, CMP, XRA, etc), and some commands require two operands (eg: MOV and MVI).

## Error Messages

In the process of entering and executing commands under Educator- 8080 there are a number of errors which can occur. When this happens an error message is displayed on the output device. For the benefit of users with television displays, a delay of approximately two seconds occurs as the message is being displayed, to provide time to read it. After the two second delay the normal Edu-cator-8080 display is generated and the command entry mode is reentered. Teletype or other hard copy users will probably wish to alter the error display routine slightly by eliminating the extraneous spaces which are used to center the error messages on the TV monitor screen.

The errors which can occur are listed in absolute octal form in table 3. The error numbers and extended explanations of conditions are as follows:

1. INPUT TOO LONG: The input string exceeds 22 characters in length probably because too many characters were deleted since delete character codes count as input characters. Twenty two characters should be sufficient for any normal entry including punctuation and several character deletes.
2. INVALID COMMAND: The input command mnemonic is not one of the ones implemented by Educator-8080.
3. INVALID REGISTER: The operand register is not $A, B$ or $C$ for a command which requires a single register as an operand or it was not $B$

Listing 7, continued:


Listing 1, continued:

| addrass | octal-code |  | label | $\infty$ | operand | commentary |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $<1>/ 317$ | 127 |  |  | MOV | D,A | Move the result beck to O ; |
| $<1>/ 320$ | 341 |  |  | POP | H | Unieve the count; |
| $<1>/ 321$ | 003 |  |  | INX | B | Increment commend character index; |
| <1>/322 | 045 |  |  | DCR | H | Decrement the count; |
| <1>/323 |  | 303 <1> |  | JNZ | BLOOP | If not zero then reiterate: |
| <1>/326 | 257 |  |  | XRA | A | Clear $\mathbf{A}$ indicating no arrors; |
| $<1>/ 327$ | 311 |  |  | RET |  | Return to CNTRL; |
| -The immediate error routine is used vo indicate immedists value errors. |  |  |  |  |  |  |
| $\begin{aligned} & <1>/ 330 \\ & <1>/ 332 \end{aligned}$ | $\begin{aligned} & 076 \\ & 303 \end{aligned}$ | $\begin{aligned} & 006 \\ & 063<2> \end{aligned}$ | IMMER | MVI JMP | A, $\mathbf{Q}^{\prime} 00$ G $^{\prime}$ ERROR | Move erfor code 3 to $A$; <br> Go display arror 3, invalid immediate: |
| - The OCTAD entry point to the OCTAL routine is for the defoult condition. |  |  |  |  |  |  |
| <1>/335 | 013 |  | OCTAD | DCX | $B$ | Decrement command charecter index: |
| - The OCTAL routine converts an octal immediase value into utable form. |  |  |  |  |  |  |
| $<1>/ 336$ | 046 | 003 | OCTAL | MVI | H, $\mathbf{Q}^{\prime} 003{ }^{\prime}$ | Move a 3 into H for count; |
| < $1>/ 340$ | 012 |  | OLOOP | LDAX |  | Losd naxt command charecter into A: |
| $<1>1341$ | 326 | 080 |  | SUI | O'060' | Subtract a '0' from is; |
| $<1>/ 343$ | 376 | 010 |  | CPI | $0^{\prime} 010^{\circ}$ | Is commend charecter leas then 8? |
| <1>/345 |  | 330 <1> |  | JNC | IMMEA | If not then go ditplay immediate error: |
| < $1>/ 350$ | 345 |  |  | PUSH |  | Save the count; |
| <1>/351 | 152 |  |  | MOV | L.D | Move D to L immediate byte; |
| <1>/352 | 051 |  |  | DAD | H | Shift immediate |
| <1>/353 | 051 |  |  | DAD | H | byte left |
| $<1>/ 354$ | 051 |  |  | DAD | H | three bits; |
| <1>/365 | 206 |  |  | ADD | L | Add $L$ to value in $A$; |
| $<1>/ 356$ | 127 |  |  | MOV | D,A | Move retult beck to D; |
| <1>/357 | 341 |  |  | POP | ${ }^{H}$ | Unseve the count; |
| $<1>1360$ | 003 |  |  | INX | B | Incrament command character index: |
| <1>/361 | 045 |  |  | DCR | H | Decrement the count; |
| <1>/362 | 302 | 340 <1> |  | JNZ | OLOOP | If not zero then reiterste; |
| <1>/365 | 257 |  |  | XRA | A | Clear A indicating no errors; |
| $<1>/ 386$ | 311 |  |  | RET |  | Return to CNTRL: |
| -The HEX routine converts a hexadecimal immediate value into usable form. |  |  |  |  |  |  |
| < $1>/ 387$ | 046 | 002 | HEX | MVI | H,0002' | Move a 2 into H for count; |
| $<1>/ 371$ | 012 |  | HLOOP | LDAX |  | Losd next command charscter into $A$ : |
| < $1>/ 372$ | 326 | 060 |  | SUI | $0060 \cdot$ | Subtract a '0' fromit; |
| $<1>1374$ | 376 | 012 |  | CPI | 0012 | Is it less than '9' ${ }^{\text {a }}$ ( 1 ? |
| <1>/376 | 332 | 010 <2> |  | JC | HCHOK | If so then numaric character is OK; |
| $<2>1001$ | 326 | 007 |  | SUI | Q'007' | Else convert alphabetic to numeric; |
| $<2>1003$ | 376 | 020 |  | CPI | Q020* | Is character value greater than 15? |
| $<2>1005$ | 322 | 330 <1> |  | JNC | IMMER | If so then invalid hexadecimal value; |
| <2>/010 | 345 |  | HCHOK | PUSH | H | Save the count; |
| <2>/011 | 152 |  |  | MOV | L, D | Move D to L immediate byte; |
| <2>/012 | 051 |  |  | DAD | H | Shift immediate |
| $<2>1013$ | 051 |  |  | DAD | H | byto left |
| <2>/014 | 051 |  |  | DAD | H | four |
| $<2>1015$ | 051 |  |  | DAD | H | bits |
| $<2>/ 016$ | 205 |  |  | ADD | $\stackrel{L}{L}$ | Add $L$ to value in $A$; |
| <2>/017 | 127 |  |  | MOV | D,A | Move reault back to D: |
| $<2>1020$ | 341 |  |  | POP | H | Unseve the count: |
| <2>1021 | 003 |  |  | INX | 8 | Increment command character index: |
| <2>/022 | 045 |  |  | DCR | H | Decrament the count; |
| <2>1023 | 302 | 371 <1> |  | JNZ | HLOOP | If not zero then reiterste; |
| <2>/026 | 257 |  |  | XRA | A | Clear A indicating no errors; |
| <2>/027 | 311 |  |  | RET |  | Return to CNTRL: |


| $<2>1030$ | 353 |  |  | xater | XCHG |  | Move generated opcode to HL: |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| <2>1031 | 042 | 046 | <2> |  | SHLD | XOTOP | Store it at execution point; |
| <2>1034 | 052 | 346 | <3> |  | LHLD | PSWA | Load working PSW \& A into HL; |
| <2>/037 | 345 |  |  |  | PUSH | H | Push \& pop sets values for |
| <2>1040 | 361 |  |  |  | POP | PSW | working register and flags: |
| <2>/041 | 052 | 350 | <3: |  | LHLO | BANDC | Load working $B$ and C in to HL : |
| <2>1044 | 345 |  |  |  | PUSH | H | Push \& pop sets values for |
| <2 1045 | 301 |  |  |  | POP | B | working 8 and C registers; |
| $<2>1046$ | 000 |  |  | xatop | NOP |  | The command to be executed; |
| <2> 1047 | 000 |  |  |  | NOP |  | Immediate value or NOP: |
| <2>/050 | 305 |  |  |  | PUSH | B | Push B and C working register values; |
| <2>/051 | 341 |  |  |  | POP | H | Pop them into HL ; |
| $<2>1052$ | 042 | 350 | <3> |  | SHLO | BANDC | Store them in save area; |
| <2>1055 | 365 |  |  |  | PUSH | PSW | Push PSW and A working values: |
| <2>1056 | 341 |  |  |  | POP | H | Pop them into HL: |
| <2>1057 | 042 | 346 | <3> |  | SHLD | PSWA | Store them in save area; |
| <2>/062 | 311 |  |  |  | RET |  | Return to CNTRL for next command; ${ }^{\text {; }}$ |
| -The ERROR routine is used to display error messages. |  |  |  |  |  |  |  |
| <2>1063 | 365 |  |  | ERROR | PUSH | PSW | Save error code in A: |
| <2>1064 | 041 | 162 | <3> |  | LXI | H,ERRSP | Load address of error header spaces; |
| <2>1067 | 315 | 261 | <0> |  | CALL | CHEDT | Go output error header spaces |
| <2>1072 | 361 |  |  |  | POP | PSW | Unsave error code: |
| $<2>1073$ | 041 | 014 | <3> |  | LXI | H,ERTAB | Load addpess of error message table: |
| $<2>1076$ | 205 |  |  |  | ADD |  | Add low address byte to error code: |
| <2>1077 | 157 |  |  |  | MOV | L,A | Move result to L, points to offset; |
| <2>/100 | 156 |  |  |  | MOV | L,M | Move offset to L; |
| - Note: HL now contains the address of the error message. |  |  |  |  |  |  |  |
| <2>/101 | 315 | 261 | <0> |  | CALL | CHEDT | Output the error message; |
| <2>/104 | 021 | 000 | 000 | ERTIM | LXI | D, ${ }^{\prime} 000000^{\prime}$ | Load DE with timing loop value; |
| $<2>1107$ | 035 |  |  |  | DCR |  | Decrement value in E 256 times; |
| <2>/110 | 302 | 105 | <2> |  | JNZ | ERTIM +1 | Reiterate loop 256 times; |
| - The above JMP goes to the first 000 in the LXI command which is an effective NOP. |  |  |  |  |  |  |  |
| $<2>1113$ | 025 |  |  |  | DCR |  | Decrement D: |
| $<2>1114$ | 302 | 105 | <2> |  | JNZ | ERTIM+1 | Reiterate outer loop 256 times; |
| <2>1117 | 076 | 377 |  |  | MVI | A, $\mathbf{Q}^{\prime} 377$ | Move a 377 to A indicating errar: |
| <2>/121 | 311 |  |  |  | RET |  | Return to CNTRL; |

for the INX or DCX commands which require register pairs as operands.
4. INVALID IMMED TYPE: The type code for an immediate operand is not $B, Q$, or $H$, or if the default was attempted the first digit of the implied octal value was not a digit from 0 to 7 .
5. INVALID IMMED VALUE: One of the characters in the immediate operand value string was inconsistent with the immediate type code. For example, a digit in a binary input string was not a zero or a one. This can also be caused by not providing the correct quantity of digits for the immediate type specified; too few digits will possibly cause a problem. If too many digits are entered only the first $N$ will be used ( $\mathrm{N}=8$ for binary, $\mathrm{N}=3$ for octal and $\mathrm{N}=2$ for hexadecimal).
6. ERROR! This message should not occur unless a grave internal error occurs in Educator-8080.

## Educator-8080 Program Listing

The Educator-8080 program is presented in an assembly language format as listing 1. It was hand assembled and, therefore, some liberties were taken in the way it was presented. Addresses are shown in a split octal ("Intelese") format of page and address within page. Educator-8080 requires four contiguous 256 byte pages of memory (it just fits); to ease the implementation process all addresses and address sensitive bytes are shown with relative page numbers in the format $\langle\mathrm{P}\rangle$, where P is a $0,1,2$ or 3 . A simple process of substitution as the program is being put into the machine will provide the ability to locate Educator- 8080 in any four contiguous pages provided the program begins on a page boundary.

The assignment of three addresses is left to the user. These three addresses are shown symbolically in both the source and the object code. The first address is for the location of the STACK; insert the address of the stack in the command at location $<0>/ 000$. The stack should be capable of being at least 10 to 12 levels deep to function correctly. The second and third addresses are the addresses of the physical input and output routines which must be provided by the user. These routine addresses are shown symbolically as KEYBD and CHRPR in the source listing. The values are shown as ' XXX XXX' in the object code.

## Input and Output Routines

The Educator-8080 program references two subroutines for the purpose of exe-

Table 3:
Error Messages. This table consists of a list of address offsets (location <3>1014) followed by the ASCII error message strings. The octal values 201 through 377 are used to encode from 1 to 177 spaces (1 to 127 decimal). The strings contain a single space for these codes. The octal value 200 is used to indicate end of string, and is shown symbolically as the character " $\nabla$ ". The octal value 177 is used to indicate the clear screen operation, and is shown symbolically as the character "■".
cuting IO operations. The KEYBD subroutine is used to read a single character of input from an ASCII keyboard device. The CHRPR subroutine is used to display (or print) a single character. These routines are not shown in the listings, but should be adapted from the routines normally used with the particular system in which the program is run. Both KEYBD and CHRPR use the accumulator (A) to pass a single character argument. KEYBD defines a value in A obtained from the input device. CHRPR displays the value in A on a device such as a video display or Teletype. All other registers of the 8080 processor should be left unchanged upon return from either of these routines. Entry to the IO routines is shown using a CALL instruction in these listings. A corresponding RET instruction in the routine should return control when either operation is completed. An alternate method of entry would be to employ the 8080 RST instruction in place of CALL. If the Edu-cator-8080 listings accompanying this article are used without reassembly, then the CALL instructions would be replaced by an RST and two single byte NOP instructions.

The keyboard entry routine KEYBD

Continued on page 75

| Address | Octal Code |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $<3>/ 014$ | 153 | 024 | 043 | 063 | 063 | 104 | 127 | 153 |
| $<3>1024$ | 111 | 116 | 120 | 125 | 124 | 040 | 124 | 117 |
| $<3>1034$ | 117 | 040 | 114 | 117 | 116 | 107 | 200 |  |
| $<3>1043$ | 111 | 116 | +26 | 101 | 114 | 111 | 104 | 040 |
| $<3>/ 053$ | 103 | 117 | 115 | 115 | 101 | 116 | 104 | 200 |
| $<3>1063$ | 111 | 116 | 126 | 101 | 114 | 111 | 104 | 040 |
| $<3>1073$ | 122 | 105 | 107 | 111 | 123 | 124 | 105 | 122 |
| $<3>/ 103$ | 200 |  |  |  |  |  |  |  |
| $<3>/ 104$ | 111 | 116 | 126 | 101 | 114 | 111 | 104 | 040 |
| $<3>/ 114$ | 111 | 115 | 115 | 105 | 104 | 040 | 124 | 131 |
| $<3>/ 124$ | 120 | 105 | 200 |  |  |  |  |  |
| $<3>/ 127$ | 111 | 116 | 126 | 101 | 114 | 111 | 104 | 040 |
| $<3>/ 137$ | 111 | 115 | 115 | 105 | 104 | 040 | 126 | 101 |
| $<3>/ 147$ | 114 | 125 | 105 | 200 |  |  |  |  |
| $<3>/ 153$ | 105 | 122 | 122 | 117 | 122 | 041 | 200 |  |

## ASCII String Value

Address Offsets for messages 0 through 7 'INPUT TOO LONG $\nabla$ ' 'INVALID COMMAND $\nabla$ ' -INVALID REGISTER $\nabla$.
'INVALID IMMED TYPE $\nabla$ ' 'INVALIO IMMED VALUE $\nabla$ ' 'ERROR! $\nabla$ '

The following string is given the name "ERRSP' and is used to clear the screen, then space down to the center prior to displaying an error message.
$<3>/ 162 \quad 177377377211200 \quad$ • $\nabla$.

| Addrest | Octal Code |  |  |  |  | Name |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $<3>/ 167$ | 177 | 211 | 105 | 104 | 125 | 103 |

Table 4: Educator-8080 standard display format messages. This table contains the definitions of several character string messages which are used to format the output display device. As in table 3, the codes from octal 201 to 377 represent from 1 to 177 spaces transmitted. The character " $\nabla$ " is used to indicate an end of text code, octal 200. The character " $\square$ " is used to indicate a clear screen code, octal 177.

| SOLID STATES MUSIC PRODUCTS |
| :---: |
| 4Kx8 Static Memories |
| MB- 1 MK-8 board, 1 usec 2102 s or equivalent. |
| Kit .......................... \$103 |
| MB-2 Altair 8800 compatible, may be piggybacked for $8 \mathrm{~K} \times 8$. |
| Kit (1us 2102s or equiv.) . ....... \$112 |
| Kit (.55us 91 L02As) . . . . . . . . . . . $\$ 132$ |
| Erom Board |
|  |
| IMSAI 8080 plug compatible, on board |
| selection of address \& wait cycles, 2K may be expanded to $4 K$. |
| Kit 2K (81702A's) .............. $\$ 145$ |
| Kit 4K (161702A's) . . . . . . . . . . $\$$ \$225 |

## 1/O Boards

1/0-1 8 bit parallel input \& output ports, common address decoding jumper selected, Altair 8800 plug compatible.

Kit . . . . . . \$42 PC Board only . . . \$25
I/O-2 I/O for 8800, 2 ports committed, pads of 3 more, other pads for EROMS UART, etc.
Kit . . . . . \$47.50 PC Board only . . . \$25
Misc.
Altair compatible mother board . . . . $\$ 45$
$32 \times 32$ Video board Kit . .......... . $\$ 125$

| 2102's | 1usec | .65usec | .5usec |
| :--- | :--- | :--- | :--- |
| ea | $\$ 1.95$ | $\$ 2.25$ | $\$ 2.50$ |
| 32 | $\$ 59.00$ | $\$ 68.00$ | $\$ 76.00$ |


| $1702 A^{*}$ | lus | $\$ 10.00$ | $2112-1$ |
| :--- | :--- | :--- | :--- |
| $1702 A^{*} .5 u s$ | $\$ 13.00$ | $74 \mathrm{CB9}$ | $\$ 4.50$ |
| 2101 | $\$ 4.50$ | $74 L 89$ | $\$ 3.50$ |
| $2111-1$ | $\$ 4.50$ | 74200 | $\$ 5.90$ |
| $4002-1$ | $\$ 7.50$ | $74 L 200$ | $\$ 5.90$ |
| $4002-2$ | $\$ 7.50$ | 8223 | $\$ 3.00$ |
| 7489 | $\$ 2.50$ | 91 L02A | $\$ 2.55$ |
|  | *Programing send | 32 ea | $\$ 2.40$ |
| hex list | $\$ 5.00$ | 2602 | $\$ 2.00$ |

Please send for complete listing of $1 \mathrm{C}^{\prime}$ s and Xistors at competitive prices.
 San Carlos, Calif. 94070

Check or money order only. Calif, residents $6 \%$ tax. All orders postpaid in US. All devices tested prior to sale. Money back 30 day Guarantee. $\$ 10 \mathrm{~min}$. order. Prices subject to change without notice.

# Reprinted from MACHINE LANGUAGE PROGRAMMING FOR THE " 8008 " (and similar microcomputers). 

Author: Nat Wadsworth
Copyright 1975
Copyright 1976 - Revised
Scelbi Computer Consulting Inc With the permission of the
copyright owner.

# MACHINE LANGUAGE 

## Chapter I

## THE ‘8008' CPU INSTRUCTION SET

The '8008' microprocessor has quite a comprehensive instruction set that consists of 48 basic instructions, which, when the possible permutations are considered, result in a total set of about 170 instructions.

The instruction set allows the user to direct the computer to perform operations with memory, with the seven basic registers in the CPU, and with INPUT and OUTPUT ports.

It should be pointed out that the seven basic registers in the CPU consist of one "accumulator," a register that can perform mathematical and logic operations, plus an additional six registers, which, while not having the full capability of the accumulator, can perform various useful operations. These operations include the ability to hold data, serve as an "operator" with the accumulator, and increment or decrement their contents. Two of these six registers have special significance because they may be used to serve as a "pointer" to locations in memory.

The seven CPU registers have arbitrarily been given symbols so that we may refer to them in an abbreviated language. The first register is designated by the symbol ' $A$ ' in the following discussion and will be referred to as the "accumulator" register. The next four registers will be referred to as the ' $B$,' ' $C$,' ' $D$ ' and ' $E$ ' registers. The remaining two special memory pointing registers shall be designated the 'H' (for the HIGH portion of a memory address) and the 'L' (for the LOW portion of a memory address) registers.

The CPU also has several "flip-flops" which shall be referred to as "FLAGS." The flipflops are set as the result of certain operations and are important because they can be "tested" by many of the instructions with the instruction's meaning changing as a consequence of the particular status of a FLAG at the time the instruction is executed. There are four basic flags which will be referred to in this manual. They are defined as follows:

The ' C ' flag refers to the carry bit status. The carry bit is a one unit register which changes state when the accumulator overflows or underflows. This bit can also be set to a known condition by certain types of instructions. This is important to remember when developing a program because quite often a program will have a long string of instructions which do not utilize the carry bit or care about its status, but which will be causing the carry bit to change its state from time-to-time. Thus, when one prepares to do a series of operations that will rely on the carry bit, one often desires to set the carry bit to a known state.

The ' $Z$ ' for zero flag refers to a one unit register that when desired will indicate whether the value of the accumulator is exactly equal to zero. In addition, immediately after an increment of decrement of the $B, C, D, E, H$ or $L$ registers, this flag will also indicate whether the increment or decrement caused that particular register to go to zero.

The ' $S$ ' for sign flag refers to a one unit register that indicates whether the value in the accumulator is a positive or negative value (based on two's complement nomenclature). Essentially, this flag monitors the most significant bit in the accumulator and is "set" when it is a one.

The ' $P$ ' flag refers to the last flag in the group which is for indicating when the accumulator contains a value which has even parity. Parity is useful for a number of reasons and is usually used in conjunction with testing for error conditions on words of data especially when transferring data to and from external devices. Even parity occurs when the number of bits that are a logic one in the accumulator is an even value. Zero is considered an even value for this purpose. Since there are eight bits in the accumulator, even parity will occur when zero, two, four or six bits are in the logic one condition regardless of what order they may appear in within the register.

It is important to note that the $Z, S$, and P flags (as well as the previously mentioned C flag) can all be set to known states by certain instructions. It is also important to note that some instructions do not result in the flags being set so that if the programmer desires to have the program make decisions based on the status of flags, the programmer should ensure that the proper instruction, or sequence of instructions is utilized. It is particularly important to note that load register instructions do not by themselves set the flags. Since it is often desirable to obtain a data word (that is, load it into the accumulator) and test its status for such parameters as whether or not the value is zero, or a negative number. and so forth, the programmer must remember to follow a load instruction by a logical instruction (such as the NDA - "and the accumulator") in order to set the flags before using an instruction that is conditional in regards to a flag's status.

The description of the various types of instructions available using an '8008' CPU which follows will provide both the machine language code for the instruction given as three octal digits, and also a mnemonic name suitable for writing programs in "symbolic" type language which is usually easier than trying to remember octal codes! It may be noted that the symbolic language used is the same as that originally suggested by Intel Corporation which developed the ' 8008 ' CPU-on-a-chip. Hence users who may already be familiar with the suggested mnemonics will not have any relearning problems and those learning the mnemonics for the first time will have plenty of good company. If the programmer is not already aware of it, the use of mnemonics facilitates working with an "assembler" program when it is desired to develop relatively large and complex programs. Thus the programmer is urged to concentrate on learning the mnemonics for the instructions and not waste time memorizing the octal codes. After a program has been written using the mnemonic codes, the programmer can always use a lookup table to convert to the machine code if an assembler program is not available. It's a lot easier technique (and less subject to error) than trying to memorize

# PROGRAMMING FOR THE "8008" 

## and similar microcomputers

the 170 or so three digit combinations which make up the machine instruction code set!

The programmer must also be aware, that in this machine, some instructions require more than one word in memory. "Immediate" type commands require two consecutive words. JUMP and CALL commands require three consecutive words. The remaining types only require one word.

The first group of instructions to be presented are those that are used to load data from one CPU register to another, or from a CPU register to a word in memory, or vice-versa. 'This group of instructions requires just one word of memory. It is important to note that none of the instructions in this group affect the flags.

## LOAD DATA FROM ONE CPU REGISTER TO ANOTHER CPU REGISTER

## MNEMONIC MACHINE CODE

| LAA | 300 |
| :---: | :---: |
| LBA | 310 |
| $\cdot$ | $\cdot$ |
| LAB | 301 |

The load register group of instructions allows the programmer to move the contents of one CPU register into another CPU register. The contents of the originating (from) register is not changed. The contents of the destination (to) register becomes the same as the originating register. Any CPU register can be loaded into any CPU register. Note that loading register A into register A is essentially a NOP (no operation) command. When using mnemonics the load symbol is the letter L followed by the "to" register and then the "from" register. The mnemonic LBA means that the contents of register A (the accumulator) is to be loaded into register $B$. The mnemonic LAB states that register B is to have its contents loaded into register $A$. It may be observed that this basic instruction has many variations. The machine language coding for this instruction is in the same format as the mnemonic code except that the letters used to represent the registers are replaced by numbers that the computer
can use. Using octal code, the seven CPU registers are coded as follows:

Register $\mathrm{A}=0$
Register $B=1$
Register $C=2$
Register $\mathrm{D}=3$
Register $\mathrm{E}=4$
Register $\mathrm{H}=5$
Register L = 6

Also, since the machine can only utilize numbers, the octal number ' 3 ' in the most significant location of a word signifies that the computer is to perform a "load" operation. Thus, in machine coding, the instruction for loading register B with the contents of register A becomes ' 310 ' (in octal form). Or, if one wanted to get very detailed, the actual binary coding for the eight bits of information in the instruction word would be ' 11001000 .' It is important to note that the load instructions do not affect any of the flags.

## LOAD DATA FROM ANY CPU REGISTER TO A LOCATION IN MEMORY

| LMA | 370 |
| :--- | :--- |
| LMB | 371 |
| LMC | 372 |
| LMD | 373 |
| LME | 374 |
| LMH | 375 |
| LML | 376 |

This instruction is very similar to the previous group of instructions except that now the contents of a CPU register will be loaded into a specified memory location. The memory location that will receive the contents of the particular CPU register is that whose address is specified by the contents of the CPU H and L registers at the time the instruction is executed. The H CPU register specifies the HIGH portion of the address desired, and the L CPU register specifies the LOW portion of the address into which data from the selected CPU register is to be loaded. Note that there are seven different instruc-
tions in this group. Any CPU register can have its contents loaded into any location in memory. This group of instructions does not affect any of the flags.
LOAD DATA FROM A MEMORY
LOCATION TO ANY CPU REGISTER

LAM
LBM
LCM
LDM

This group of instructions can be considered the opposite of the previous group. Now, the contents of the word in memory whose address is specified by the H (for HIGH portion of the address) and L (LOW portion of the address) registers will be loaded into the CPU register specified by the instruction. Once again, this group of instructions has no affect on the status of the flags.

## LOAD IMMEDIATE DATA INTO A CPU REGISTER

| LAI | 006 |
| :--- | :--- |
| LBI | 016 |
| LCI | 026 |
| LDI | 036 |
| LEI | 046 |
| LHI | 056 |
| LLI | 066 |

An IMMEDIATE type of instruction requires two words in order to be completely specified. The first word is the instruction itself. The second word, or "immediately following" word, must contain the data upon which "immediate" action is taken. Thus, a load IMMEDIATE instruction in this group means that the contents of the word immediately following the instruction word is to be loaded into the specified register. For example, a typical load immediate instruction would be LAI 001. This would result in the value 001 (octal) being placed in the A register when the instruction was executed. It is important to remember that all IMMEDIATE type in-
structions MUST be followed by a data word An instruction such as LDI by itself would result in improper operation because the computer would assume the next word contained data. If the programmer had mistakenly left out the data word, and in its place had another instruction, the computer would not realize the operator's mistake. Hence the program would be fouled-up! Note too, that the load immediate group of instructions does not affect the flags.

## LOAD IMMEDIATE DATA INTO A MEMORY LOCATION

## LMI

076
This instruction is essentially the same as the load immediate into the CPU register group except that now, using the contents of the H and L registers as "pointers" to the desired address in memory, the contents of the "immediately following word" will be placed in the memory location specified. This instruction does not affect the status of the flags.

The above rather large group of LOAD instructions permits the programmer to direct the computer to move data about. They are used to bring in data from memory where it can be operated on by the CPU. Or, to temporarily store intermediate results in the CPU registers during complicated and extended calculations, and of course allow data, such as results, to be placed back into memory for long term storage. Since none of them will alter the contents of the four CPU flags, these instructions can be called upon to set up data before instructions that may affect or utilize the flag's status are executed. The programmer will use instructions from this set frequently. The mnemonic names for the instructions are easy to remember as they are well ordered. The most important item to remember about the mnemonics is that the TO register is always indicated first in the mnemonic, and then the FROM register. Thus LBA equals "load TO register B FROM register A .

## INCREMENT THE VALUE OF A CPU REGISTER BY ONE

| INB | 010 |
| :--- | :--- |
| INC | 020 |
| IND | 030 |
| INE | 040 |
| INH | 050 |
| INL | 060 |

This group of instructions allows the programmer to add one to the present value of any of the CPU registers except the accumulator. (Note carefully that the accumulator can NOT be incremented by this type of instruction. In order to add one to the accumulator a mathematical addition instruction, described later, must be used.) This instruction for incrementing the defined CPU registers is very valuable in a number of applications. For one thing, it is an easy way to have the L register successively "point" to a string of locations in memory. A feature that makes this type of instruction even more
powerful is that the result of the incremented register will affect the $\mathrm{Z}, \mathrm{S}$, and P flags. (It. will not change the C or "carry"flag.) Thus, after a CPU register has been incremented by this instruction, one can utilize a flag test instruction (such as the conditional JUMP and CALL instructions to be described later) to determine whether that particular register has a value of zero ( Z flag), or if it is a negative number ( S flag), or even parity ( P flag). It is important to note that this group of instructions, and the decrement group (described in the next paragraph) are the only instructions which allow the flags to be manipulated by operations that are not concerned with the accumulator (A) register.

## DECREMENT THE VALUE OF A CPU REGISTER BY ONE

| DCB | 011 |
| :--- | :--- |
| DCC | 021 |
| DCD | 031 |
| DCE | 041 |
| DCH | 051 |
| DCL | 061 |

The DECREMENT group of instructions is similar to the INCREMENT group except that now the value one will be subtracted from the specified CPU register. This instruction will not affect the C flag. But, it does affect the $Z, S$, and $P$ flags. It should also be noted that this group, as with the increment group, does not include the accumulator register. A separate mathematical instruction must be used to subtract one from the accumulator.

## ARITHMETIC INSTRUCTIONS USING THE ACCUMULATOR

The following group of instructions allow the programmer to direct the computer to perform arithmetic operations between other CPU registers and the accumulator, or between the contents of words in memory and the accumulator. All of the operations for the described addition, subtraction, and compare instructions affect the status of the flags.

## ADD THE CONTENTS OF A CPU REGISTER TO THE ACCUMULATOR

| ADA | 200 |
| :--- | :--- |
| ADB | 201 |
| ADC | 202 |
| ADD | 203 |
| ADE | 204 |
| ADH | 205 |
| ADL | 206 |

This group of instructions will simply ADD the present contents of the accumulator register to the present value of the specified CPU register and leave the result in the accumulator. The value of the specified register is unchanged except in the case of the ADA instruction. Note that the ADA instruction essentially allows the programmer to double the value of the accumulator (which is the $A$ register)! If the addition
causes an overflow or underflow then the carry (C flag) will be affected.

## ADD THE CONTENTS OF A CPU

 REGISTER PLUS THE VALUE OF THE CARRY FLAG TO THE ACCUMULATOR| ACA | 210 |
| :--- | :--- |
| ACB | 211 |
| ACC | 212 |
| ACD | 213 |
| ACE | 214 |
| ACH | 215 |
| ACL | 216 |

This group is identical to the previous group except that the content of the carry flag is considered as an additional bit (MSB) in the specified CPU register. The combined value of the carry bit plus the contents of the specified CPU register are added to the value in the accumulator. The results are left in the accumulator. Again, with the exception of the ACA instruction, the contents of the specified CPU register are left unchanged. Again too, the carry bit (C flag) will be affected by the results of the operation.

\section*{SUBTRACT THE CONTENTS OF A CPU REGISTER FROM THE ACCUMULATOR <br> | SUA | 220 |
| :--- | :--- |
| SUB | 221 |
| SUC | 222 |
| SUD | 223 |
| SUE | 224 |
| SUH | 225 |
| SUL | 226 |}

This group of instructions will cause the present value of the specified CPU register to be subtracted from the value in the accumulator. The value of the specified register is not changed except in the case of the SUA instruction. (Note that the SUA instruction is a convenient instruction with which to "clear" the accumulator.) The carry flag will be affected by the results of a SUBTRACT instruction.

\section*{SUBTRACT THE CONTENTS OF A CPU REGISTER AND THE VALUE OF THE CARRY FLAG FROM THE ACCUMULATOR <br> | SBA | 230 |
| :--- | :--- |
| SBB | 231 |
| SBC | 232 |
| SBD | 233 |
| SBE | 234 |
| SBH | 235 |
| SBL | 236 |}

This group is identical to the previous group except that the content of the carry flag is considered as an additional bit (MSB) in the specified CPU register. The combined value of the carry bit plus the contents of the specified CPU register are SUBTRACTED from the value in the accumulator. The re. sults are left in the accumulator. The carry
bit (C flag) is affected by the result of the operation. With the exception of the SBA instruction the content of the specified CPU register is left unchanged.
COMPARE THE VALUE IN THE
ACCUMULATOR AGAINST THE
CONTENTS OF A CPU REGISTER

CPA
CPB
CPC
CPD
CPE
CPH
CPL

The COMPARE group of instructions are a very powerful and somewhat unique set of instructions. They direct the computer to compare the contents of the accumulator against another register and to set the flags as a result of the comparing operation. It is essentially a subtraction operation with the value of the specified register being subtracted from the value of the accumulator except that the value of the accumulator is not actually altered by the operation. However, the flags are set in the same manner as though an actual subtraction operation had occured. Thus, by subsequently testing the status of the various flags after a COMPARE instruction has been executed, the program can determine whether the compare operation resulted in a match or non-match. In the case of a non-match, one may determine if the compared register contained a value greater or less than that in the accumulator. This would be accomplished by testing the Z flag and C flag respectively utilizing a conditional JUMP or CALL instruction (which will be described later).

## ADDITION, SUBTRACTION, AND <br> COMPARE INSTRUCTIONS THAT USE <br> WORDS IN MEMORY AS OPERANDS

The five types of mathematical operations: ADD, ADD with CARRY, SUBTRACT, SUBTRACT with CARRY, and COMPARE which have just been presented for the cases where they operate with the contents of CPU registers, can all be performed with words that are in memory. As with the LOAD in structions that operate with memory, the H and $L$ registers must contain the address of the word in memory that it is desired to ADD, SUBTRACT, or COMPARE to the accumulator. The same conditions for the operations as was detailed when using the CPU registers apply. Thus, for mathematical operations with a word in memory, the following instructions are used.

## ADD THE CONTENTS OF A MEMORY WORD TO THE ACCUMULATOR

ADM
207

ADD THE CONTENTS OF A MEMORY WORD PLUS THE VALIE OF THE CARRY FLAG TO THE ACCUMULATOR

## SUBTRAC'T THE CONTENTS OF A MEMORY WORD FROM THE aCCUMULATOR

```
            SUM 227
SUBTRACT THE CONTENTS OF A
MEMORY WORD AND THE VALUE
OF THE CARRY FLAG FROM THE
            ACCUMULATOR
        SBM
        237
    COMPARE THE VALUE IN THE
    ACCUMULATOR AGAINST THE
CONTENTS OF A MEMORY WORD
CPM
2 7 7
```


## IMMEDIATE TYPE ADDITIONS, SUBTRACTIONS, AND COMPARE INSTRUCTIONS

The five types of mathematical operations discussed above can also be performed with the operand being the word of data immediately after the instruction. This group of instructions is similar in format to the previously described LOAD IMMEDIATE instructions. The same conditions for the mathematical operations as discussed for the operations with the CPU registers apply.

## ADD IMMEDIATE

## ADI <br> 004

ADD WITH CARRY IMMEDIATE
$\mathrm{ACl} \quad 014$

SUBTRACT IMMEDIATE

## SUI <br> 024

SUBTRACT WITH CARRY IMMEDIATE
SBI 034

## COMPARE IMMEDIATE <br> $\mathrm{CPI} \quad 074$

LOGICAL INSTRUCTIONS WITH THE

## ACCUMULATOR

There are several groups of instructions which allow BOOLEAN LOGIC operations to be performed between the contents of the CPU registers and the A (accumulator) register. In addition there are logic IMMEDIATE type instructions. The boolean logic operations are valuable in a number of programming applications. The instruction set allows three basic boolean operations to be performed. These are: the LOGICAL AND, the LOGICAL OR, and the EXCLUSIVE OR
operations. Each type of logic operation is performed on a bit-by-bit basis between the accumulator and the CPU register or memory location specified by the instruction. A detailed explanation of each type of logic operation, and the appropriate instructions for each type is presented below. The logic instruction set is also valuable because all of them will cause the $C$ (carry) flag to be placed in the zero condition. This is important if one is going to perform a sequence of instructions that will eventually use the status of the $C$ flag to arrive at a decision as it allows the programmer to set the C flag to a known state at the start of the sequence. All other flags are set in accordance with the result of the logic operation. Hence, the group often has value when the programmer desires to determine the contents of a register that has just been loaded into a register. (Since the load instructions do not alter the flags.)

## THE BOOLEAN 'AND' OPERATION INSTRUCTION SET

When the boolean AND instruction is executed, each bit of the accumulator will be compared with the corresponding bit in he register or memory focation specified by the instruction. As each bit is compared a logic result will be placed in the accumulator for each bit comparison. The logic result is determined as follows. If both the bit in the accumulator and the bit in the register with which the operation is being performed are a logic one, then the accumulator bit will be left in the logic one condition. For all other possible combinations (A hit equals one, $\boldsymbol{X}$ bit equals zero; $A$ bit equals zero, $X$ bit equals one; or A bit equals zero, $X$ bit equals zero), then the accumulator bit will be cleared to the zero state. An example will illustrate the logical AND operation.

INITIAL STATE OF THE ACCUMLLATOR

## 10101010

## CONTENTS OF OPERAND REGISTER

11001100
FINAL STATE OF THE ACCUMULATOR
10001000

There are seven logical AND instructions that allow any CPU register to be used as the AND operand. They are as follows.



The contents of the operand register is not altered by an AND logical instruction.

There is also a logical AND instruction that allows a word in memory to be used as an operand. The address of the word in memory that will be used is pointed to by the conlents of the H and L ('PL' registers.

## NDM <br> 247

And finally there is also a logical AND IMMEDIATE type of instruction that will use the contents of the word immediately following the instruction as the operand.

$$
\therefore \text { DI } \quad 0.44
$$

The next group of boolean logic instructions direct the computer to perform the logical OR operation on a bit-by-hit basis with the accumulator and the contents of a (PU) register or a word in memory. The logical OR operation will result in the acrumulator having a bit set to a logic one if either that bit in the accumulator, or the corresponding bit in the operand register is a logic one. Since the case where both the accumulator bit and operand bit are a one also satisfies the criteria, that condition will also result in the accumulator bit being left in the one state. If neither register has a logic one in the bit position, then the accumulator bit for that position remains in the zero state. An example illustrates the results of a logical OR operation.

INITIAL STATE OF THE ACCUMULATOR
10101010

CONTENT OF THE OPERAND REGISTER
11001100

## FINAL STATE OF THE ACCUMULATOR

## 11101110

There are seven logical OR instructions that allow any CPU register to be used as the OR operand.

| ORA | 260 |
| :--- | :--- |
| ORB | 261 |
| ORC | 262 |
| ORD | 263 |
| ORE | 264 |
| ORH | 265 |
| ORL | 266 |

By using the H and L registers as pointers one can also use a word in memory as an OR operand.

## ORM

267
There is also the logical OR IMMEDIATE instruction.

$$
\text { ORI } 064
$$

As with the logical AND group of instructions, the logical OR instruction does not alter the contents of the operand register.

The last group of boolean logic instructions is a variation of the logic OR. The variation is termed the EXCLUSIVE OR logical operation. The EXCLUSIVE OR operation is similar to the OR except that when the corresponding bits in both the accumulator and the operand register are a one then the accumulator bit will be cleared to zero. Thus, the accumulator bit will be a one after the operation only if just one of the registers (accumulator register or operand register) has a one in the bit position. (Again, the operation is performed on a bit-by-bit basis.) An example provides clarification.

## INITIAL STATE OF THE ACCUMULATOR

## 10101010

## CONTENTS OF THE OPERAND REGISTER

11001100

FINAL STATE OF THE ACCUMULATOR

$$
01100110
$$

The seven instructions that allow the CPU registers to be used as operands are shown next.

| XRA | 250 |
| :--- | :--- |
| XRB | 251 |
| XRC | 252 |
| XRD | 253 |
| XRE | 254 |
| XRH | 255 |
| XRL | 256 |

The instruction that uses registers $H$ and $L$ as pointers to a memory location is:

## XRM 257

And the EXCLUSIVE OR IMMEDIATE type instruction is:

$$
\text { XRI } 054
$$

As in the case of the logical OR operation, the operand register is not altered except for the special case when the XRA instruction is used. This instruction, which directs the computer to EXCLUSIVE OR the accumulator with itself, will cause the operand register, since it is the accumulator, to have its contents altered (unless it should happen to be zero at the time the instruction is executed).

This is because, regardless of what value is in the accumulator, if it is EXCLUSIVE OR'ed with itself, the result will be zero! The example below illustrates the specific operation.

ORIGINAL VALUE OF ACCUMULATOR
10101010

## EXCLUSIVE OR'ed WITH ITSELF

10101010

## FINAL VALUE OF ACCUMULATOR

00000000

This only occurs when the logical EXCLUSIVE OR is performed on the accumulator itself. It can be shown that the results of performing the logical OR or logical AND between the accumulator and itself will result in the original accumulator value being retained.

## INSTRUCTIONS FOR ROTATING THE CONTENTS OF THE ACCUMULATOR

It is often desirable to be able to shift the contents of the accumulator either right or left. In a fixed length register, a simple shift operation would result in some information being lost because what was in the MSB or LSB (depending on in which direction the shift occured) would be shifted right out of the register! Therefore, instead of just shifting the contents of a register, an operation termed ROTATING is utilized. Now, instead of just shifting a bit off the end of the register, the bit is brought around to the other end of the register. For instance, if the register is rotated to the right, the LSB (least significant bit) would be brought around to the position of the MSB (most significant bit) which would have been vacated by the shifting of its original contents to the right. Or, in the case of a shift to the left, the MSB would be brought around to the position of the LSB.

The carry bit ( $C$ flag) can be considered as an extension of the accumulator register. The instruction set for this machine allows two types of ROTATE instructions. One considers the carry bit to be part of the accumulator register for the rotate operation. The other type does not. In addition, each type of rotate can be done either to the right or to the left.

It should be noted that the rotate operations are particularly valuable when it is desired to multiply a number or divide a number. This is because shifting the contents of a register to the left effectively multiplies a binary number by a power of two. Shifting a binary number to the right provides the inverse operation.

ROTATING THE ACCUMULATOR LEFT
RLC
002

Rotating the accumulator left with the RLC instruction means the MSB of the accumulator will be brought around to the LSB position and all other bits will be shifted one position to the left. While this instruction does not shift through the carry bit, the carry bit will be set by the status of the MSB of the accumulator at the start of the ROTATE LEFT operation. (This feature allows the programmer to determine what the MSB was prior to the shifting operation by testing the C flag after the rotate instruction has been executed.

## ROTATING THE ACCUMULATOR LEFT THROUGH THE CARRY BIT <br> RAL <br> 022

The RAL instruction will cause the MSB of the accumulator to go into the carry bit. The initial value of the carry bit will be shifted around to the LSB of the accumulator. All other bits are shifted one position to the left.

## ROTATING THE ACCUMULATOR RIGHT <br> RRC 012

The RRC instruction is similar to the RLC instruction except that now the LSB of the accumulator is placed in the MSB of the accumulator. All other bits are shifted one position to the right. Also, the carry bit will be set to the initial value of the LSB of the accumulator at the start of the operation.

## ROTATING THE ACCUMULATOR RIGHT THROUGH THE CARRY BIT

## RAR

## 032

Here, the LSB of the accumulator is brought around to the carry bit. The initial value of the carry bit is shifted to the MSB of the accumulator. All other bits are shifted a position to the right.

It should be noted that the C flag is the only flag that is altered by a rotate instruction. All other flags remain unchanged.

## JUMP INSTRUCTIONS

The instructions discussed so far have all been "direct action" instructions. The programmer arranges a sequence of these types of instructions in memory. When the program is started the computer proceeds to execute the instructions in the order in which they are encountered. The computer automatically reads the contents of a memory location, executes the instruction it finds there, and then automatically increments a special address register called a PROGRAM COUNTER that will result in the machine reading the information contained in the next sequential memory location. However, it is often desirable to perform a series of instructions located in one section of memory, and then skip over a group of memory locations and start executing instructions in another section of memory. This action can be accomplished by a group of instructions
that will cause a new address value to be placed in the PROGRAM COUNTER. This will cause the computer to go to a new section of memory and then execute instructions sequentially from the new memory location.

The JUMP instructions in this computer add considerable power to the machine's capabilities because there are a series of "conditional" JUMP instructions available. That is, the computer can be directed to test the status of a particular FLAG (C, Z, $S$ or $P$ ). If the status of the flag is the desired one, then a JUMP will be performed. If it is not, the machine will continue to execute the next instruction in the current sequence. This capability provides a means for the computer to make "decisions" and to modify its operation as a function of the status of the various flags at the time that a program is being executed.

In a manner similar to IMMEDIATE types of instructions, the JUMP instructions require more than one word of memory. A JUMP instruction requires three words to be properly defined. (Remember that IMMEDIATE type instructions required two words.) The JUMP instruction itself is the first word. The second word must contain the LOW ADDRESS portion of the address of the word in memory that the PROGRAM COUNTER is to be set to point to, which is the new location from which the next instruction is to be fetched. The third word must contain the HIGH ADDRESS (sometimes referred to as the PAGE) of the memory address that the program counter will be set to. That is, the high order portion of the address in memory that the computer will JUMP to in order to obtain its next instruction.

## THE UNCONDITIONAL JUMP INSTRUCTION

## JMP $\quad$ X4

Note: The machine code 1X4 indicates that any code for the second octal digit of the machine code is valid. It is recommended as a standard practice that the code ' 0 ' be used. Thus, the typical machine code would be 104.

Remember, the JUMP instruction must be followed by two more words which contain the LOW, and then the HIGH (PAGE) portion of the address that the program is to JUMP to!

JUMP IF THE DESIGNATED FLAG IS TRUE (CONDITIONAL JUMP)

| JTC | 140 |
| :--- | :--- |
| JTZ | 150 |
| JTS | 160 |
| JTP | 170 |

As with the UNCONDITIONAL JUMP instruction, the CONDITIONAL JUMP instructions must be followed by two words of information. The LOW portion, then the HIGH portion, of the address that program execution is to continue from if the jump is
executed. The JUMP IF TRUE group of instructions will only jump to the designated address if the condition of the appropriate flag is TRUE (logical one). Thus, the JTC instruction states that if the carry flag (C) is a logical one (TRUE) then the jump is to be executed. If it is a logical zero (FALSE) then program execution is to continue with the next instruction in the current sequence of instructions. In a similar manner the JTZ instruction states that if the ZERO FLAG is TRUE then the jump is to be performed. Otherwise the next instruction in the present sequence is executed. Likewise for the JTS and JTP instructions.

## JUMP IF THE DESIGNATED FLAG <br> IS FALSE (CONDITIONAL JUMP)

| JFC | 100 |
| :--- | :--- |
| JFZ | 110 |
| JFS | 120 |
| JFP | 130 |

As with all JUMP instructions these instructions must be followed by the LOW address then the HIGH address of the memory location that program execution is to continue from if the jump is executed. This group of instructions is the opposite of the jump if the flag is true group. For instance, the JFC instruction commands the computer to test the status of the carry (C) flag. If the flag is FALSE (a logic zero), then the jump is to be performed. If it is TRUE, then program execution is to continue with the next instruction in the current sequence of instructions. The same procedure holds for the JFZ, JFS and JFP instructions.

## SUBROUTINE CALLING INSTRUCTIONS

Quite often when a programmer is develop. ing computer programs the programmer will find that a particular algorithm (sequence of instructions for performing a function) can be used many times in different parts of the program. Rather than having to keep entering the same sequence of instructions at different locations in memory, which would not only consume the time of the programmer, but would also result in a lot of memory being used to perform one particular function, it is desirable to be able to be able to put an often used sequence of commands in just one location in memory. Then, whenever the particular algorithm is required by another part of the program, it would be convenient to jump to the section that contained the often used algorithm, perform the sequence of instructions, and then return back to the main part of the program. This is a standard practice in computer operations. A frequently used algorithm can be designated a SUBROUTINE. A special set of instructions allows the programmer to CALL a SUBROUTINE. In other words, specify a special type of JUMP command that will eventually allow the program to RETURN to the original "jumping" point in the program. A second type of instruction is used to terminate a SUBROUTINE. This special terminator will cause the program to revert back and pick up the next sequential in-
struction in memory that immediately follows the original CALLING instruction. A great deal of computer power is provided by the instruction set in this machine that allows one to CALL and RETURN from SUBROUTINES. This is because, in a manner similar to that provided for the CONDITIONAL JUMP instructions, there are a number of CONDITIONAL CALL and CONDITIONAL RETURN commands in the instruction set

Like the JUMP instructions, the CALL instructions all require three words in order to be fully specified. The first word is the CALL instruction itself. The next two words must contain the LOW and HIGH portions of the starting address of the subroutine that is being "called."

When a CALL instruction is encountered by the computer, the CPU will actually save the current value of the PROGRAM COUNTER by storing it in a special PROGRAM COUNTER PUSH-DOWN STACK. This stack is capable of holding six addresses plus the current operating address. What this means is that the machine is capable of "nesting" up to seven subroutines at one time. Thus, one can have a subroutine, that in turn calls another subroutine, that in turn calls another one, up to seven levels, and the machine will still be able to return to the initial calling location. The programmer must ensure that subroutines are not nested more than seven levels otherwise the PROGRAM COUNTER PUSH-DOWN STACK will push the original calling address(es) completely out of the push-down stack. The program could then no longer automatically return to the initial calling location.

The RETURN instruction which terminates a SUBROUTINE only requires one word. When the CPU encounters a RETURN instruction it causes the PROGRAM COUNTER PUSH-DOWN STACK to "pop" up one level. This effectively causes the address saved in the stack by the calling routine to be taken as the new program counter. Hence, program execution returns to the calling location.

## THE UNCONDITIONAL CALL INSTRUCTION

## CAL <br> 1X6

This instruction followed by two words containing the LOW and then the HIGH order of the starting address of the SUBROUTINE that is to be executed is an UNCONDITIONAL CALL. The subroutine will be executed regardless of the status of the FLAGS. The next sequential address after the CAL instruction is saved in the PROGRAM COUNT ER PUSH-DOWN STACK.

## THE UNCONDITIONAL RETURN INSTRUCTION

## RET 0X7

This instruction directs the CPU to unconditionally "pop" the program counter push-down stack UP one level.

Program execution will continue from the address saved by the subroutine calling instruction.

## CALL A SUBROUTINE IF THE DESIGNATED FLAG IS TRUE

| CTC | 142 |
| :--- | :--- |
| CTZ | 152 |
| CTS | 162 |
| CTP | 172 |

In a manner similar to the conditional JUMP IF TRUE instructions, these instructions (which must all be followed by the LOW and HIGH portions of the called subroutine's starting address) will only perform the "call" if the designated flag is in the TRUE (logical one) state. If the designated flag is FALSE then the CALL instruction is ignored. Program execution then continues with the next sequential instruction.

## RETURN FROM A SUBROUTINE IF THE DESIGNATED FLAG IS TRUE

| RTC | 043 |
| :--- | :--- |
| RTZ | 053 |
| RTS | 063 |
| RTP | 073 |

These one word instructions will cause a SUBROUTINE to be TERMINATED only if the designated flag is in the logical one (TRUE) state.

## CALL A SUBROUTINE IF THE DESIGNATED FLAG IS FALSE

| CFC | 102 |
| :--- | :--- |
| CFZ | 112 |
| CFS | 122 |
| CFP | 132 |

These instructions are the opposit of the previous group of calling commands. The subroutine is called only if the designated flag is in the FALSE (logical zero) condition. Remember, these instructions must be followed by two words which contain the LOW and HIGH part of the starting address of the SUBROUTINE that is to be executed if the designated flag is FALSE. If the flag is TRUE, the subroutine will not be called and program operation will continue with the next instruction in the current sequence.

## RETURN FROM A SUBROUTINE IF THE DESIGNATED FLAG IS FALSE

| RFC | 003 |
| :--- | :--- |
| RFZ | 013 |
| RFS | 023 |
| RFP | 033 |

These one word instructions will terminate a subroutine ("pop" the program count-
er stack UP one level) if the designated flag is FALSE. Otherwise, the instruction is ignored and program operation is continued with the next instruction in the subroutine.

## THE SPECIAL RESTART SUBROUTINE CALL INSTRUCTIONS

There is a special purpose instruction available that effectively serves as a one word SUBROUTINE CALL. (Remember that it normally requires three words to specify a subroutine call.) This special instruction allows the programmer to call a subroutine that starts at any one of eight specially designated memory locations. The eight special memory locations are at locations: $000,010,020,030,040,050,060$ and 070 on page zero. There are eight variations of the machine code for the RESTART instruction. One for each of the above addresses. Thus, the one word instruction can serve to CALL a SUBROUTINE at the specified starting location (instead of having two additional words to specify the starting address of a subroutine). It is often convenient to utilize a RESTART command as a quick CALL to an often used subroutine. Or, as an easy way to call short "starting" subroutines for large programs. Hence, the name for the type of instruction. The eight RESTART instructions, in their mnemonic and machine code forms, along with the starting address associated with each one is listed below.

| RST 0 | 005 | 00000 |
| :--- | :--- | :--- |
| RST 1 | 015 | 00010 |
| RST 2 | 025 | 00020 |
| RST 3 | 035 | 00030 |
| RST 4 | 045 | 00040 |
| RST 5 | 055 | 00050 |
| RST 6 | 065 | 00060 |
| RST 7 | 075 | 00070 |

## INPUT INSTRUCTIONS

In order to receive information from an external device the computer must utilize a group of special signal lines. The typical ' 8008 ' computer is designed to handle up to eight groups (each group having eight signal lines) of INPUT signals. A group of signals is accepted at the computer by what is referred to as an INPUT PORT. The computer controls the operation of the INPUT PORTS. Under program control, the computer can be directed to obtain the information that is on a group of lines coming in to any INPUT PORT. When this is done the information will be transferred to the accumulator. Various types of external equipment, such as an electronic keyboard or measuring instruments, can be connected to the INPUT PORTS. The INPUT PORTS are typically referred to as having numbers from ' 0 ' to ' 7 .' The typical mnemonics and machine codes for INPUT instructions are shown next.

| INP 0 | 101 |
| ---: | ---: |
| INP 1 | 103 |
| INP 6 | 115 |
| INP 7 | 117 |

It may be interesting to note that the machine codes for input ports increase by a factor of two for each port. Note too, that while the mnemonic for an input instruction has two parts, the machine code only requires one word in memory. It is also important to realize that while an input instruction brings data into the accumulator it does not affect the status of any of the CPU flags!

## OUTPUT INSTRUCTIONS

In order to output information to an external device the computer utilizes another group of signal lines which are referred to as OUTPU' PORTS. A Typical '8008' system may be equipped to service up to twenty-four OUTPUT PORTS. (Each OUTPUT PORT actually consists of eight signal lines.) An OUTPUT instruction causes the contents of the accumulator to be transferred to the signal lines of the designated OUTPUT PORT. The output ports are normally designated by octal numbers in the range 10 to 37 . The list below shows the typical mnemonics used to specify an OUTPUT PORT along with the associated machine code. (It may be interesting to note again that the machine code increases by a factor of two for each port.)

| OUT 10 | 121 |
| :--- | :---: |
| OUT 11 | 123 |
| OUT 21 | 141 |
| OUT 36 | 175 |
| OUT 37 | 177 |

An OUTPUT instruction only requires one machine code word (even though the mnemonic is typically specified in two parts). OUTPUT PORTS are connected to external devices that one desires to have the computer transmit information to, such as a CRT display, or machinery that is to be placed under computer control.

## THE HALT INSTRUCTION

There is one more instruction in the '8008' instruction set. This instruction directs the CPU to stop all operations and to remain in that state until an INTERRUPT signal is received. In a typical '8008' system an INTERRUPT signal may be generated by an operator pressing a switch or by an external piece of equipment sending an electronic signal to the CPU. This instruction is normally used when the programmer desires to terminate a program or when it is desired to have the computer wait for an operator or external device to perform some action. There are three machine codes that may be used for the HALT command.

| HLT | 000 |
| :--- | :--- |
| HLT | 001 |
| HLT | 377 |

The HALT instruction does not affect the status of the CPU flags.

## INFORMATION ON INSTRUCTION EXECUTION TIMES

When programming for "real-time" applications it is important to know how much time each type of instruction requires to be executed. With this information the programmer can develop "timing loops" or determine with substantual accuracy how much time it will take to perform a particular series of instructions. This information is especially valuable when dealing with programs that control the operations of external devices which might require events to occur at specific times.

The following table provides the nominal instruction execution time for each category of instruction used in an '8008' system. The precise time needed for each instruction
depends on how close the master clock has been set to a nominal value of 500 kilohertz. The table shows the number of cycle states required by the type of instruction followed by the nominal time required to perform the entire instruction. Since each state executes in four microseconds, the total time required to perform the instruction as shown in the table was obtained by multiplying the number of states by four microseconds. By knowing the number of states required for each instruction the programmer can often rearrange an algorithm or substitute different types of instructions to provide programs that have events occuring at precisely timed intervals.

## INSTRUCTION EXECUTION TIME TABLE

| LOAD DATA FROM A CPU REGISTER TO ANOTHER CPU REGIST | 5 | 20 Us |
| :---: | :---: | :---: |
| LOAD DATA FROM A CPU REGISTER TO A LOCATION IN MEMORY | 7 | 28 |
| LOAD DATA FROM MEMORY TO A CPU REGISTER | 8 | 32 |
| LOAD IMMEDIATE DATA INTO A CPU REGISTER | 8 | 32 |
| LOAD IMMEDIATE DATA INTO A LOCATION IN MEMORY | 9 | 36 |
| INCREMENT OR DECREMENT A CPU REGISTER | 5 | 20 |
| ARITHMETIC/COMPARE BETWEEN ACCUMULATOR \& A CPU REGISTER | 5 | 20 |
| ARITH/COMPARE BETWEEN ACCUMULATOR \& A WORD IN MEMORY | 8 | 32 |
| IMMEDIATE.ARITHMETIC AND COMPARE | 8 | 32 |
| BOOLEAN OPS BETWEEN ACCUMULATOR AND CPU REGISTERS | 5 | 20 |
| BOOLEAN OPS WITH ACCUMULATOR \& A WORD IN MEMORY | 8 | 32 |
| IMMEDIATE BOOLEAN OPERATIONS | 8 | 20 |
| ROTATE THE ACCUMULATOR | 5 | 20 |
| JUMP AND CALL COMMANDS (UNCONDITIONAL) | 11 | 44 |
| JUMP/CALLS WHEN CONDITION NOT SATISFIED (CONDITIONAL) | 9 | 36 |
| JUMP/CALLS WHEN CONDITION SATISFIED (CONDITIONAL) | 11 | 44 |
| RETURN (UNCONDITIONAL) | 5 | 20 |
| RETURN WHEN CONDITION NOT SATISFIED (CONDITIONAL) | 3 | 12 |
| RETURN WHEN CONDITION SATISFIED (CONDITIONAL) | 5 | 20 |
| RESTART COMMAND | 5 | 20 |
| OUTPUT COMMAND | 6 | 24 |
| INPUT COMMAND | 8 | 32 |
| HALT COMMAND | 4 | 16 |

[^1]
# The "Ignorance Is Bliss" 

## Television Drive Circuit



Figure 1: The "Ignorance Is Bliss" Television Drive Circuit. The components to the left of the dashed line were added as part of the interface. The components to the right of the dashed line are part of the Zenith 12FBI2X chassis which was used for the television display.

## Ken Barbier <br> PO Box 1042

Socorro NM 87801

I had not yet heard of BYTE magazine, or hams building such hardware, when I built my CRT terminal (a computerhead's term for "TV typewriter") in the fall of 1975. I didn't even own a TV set! Two situations resulted: I had to buy a new ail solid state TV, and I didn't have any idea how to interface with it. I knew approximately what it took to create horizontal and vertical sync, but had no idea whether levels, pulse widths, and frequencies would be noncritical. I was delighted, therefore, when my sync generator worked just fine the first time I patched its output across the video driver base resistor using the circuit as shown in figure 1. My big fat TTL level pulses swamp the AGC circuity so effectively that normal signals and noise from the TV IF just disappeared and I had nothing to switch off!

Not having any idea how to mix in my video (character generator output) with the sync, I just hooked up a 2 K pot where the 620 ohm is shown, and started reducing video until it stopped interfering with the
sync, and there I was at 620 ohms. All was fine, until 1 erased my character memory and started typing in one character at a time. Contrast went all to pot! I had provided no DC restoration. And I never did. At least not in the TV set circuitry.

My terminal design produces 24 lines of 64 characters each, with a total of 270 scan lines per frame. Vertical sync is the 10 scan lines that would have been character line 26. To eliminate the need for the type of DC restoration as detailed in "Television Interface" (page 20, BYTE, October 1975), I generated a black level blanking signal covering what would have been character lines 25 thru 27. This signal enters the blanking gate, IC1, at pin 2 in figure 1. Now, when I turn on my system and erase the memory, my TV field shows a nice white area with a black border top and bottom. My character generator output produces black on white characters which I find preferable to the usual white on black.

Simply by turning off the logic power I can be instantly flooded with the inanities emanating from the vast TV wasteland. With this design, I have no need to pull plugs or throw switches. Sometimes ignorance can be bliss.■

## Tool Box

Robert Baker 34 White Pine Dr Littleton MA 01460

How many different tools and fasteners can you find hidden in the matrix? The letters of each tool name are in a straight line, going horizontally, vertically or diagonally, either backwards or forwards. See how many words you can find "open loop." The official list will appear in next month's BYTE.

HS URBWMWRAEHSOR AJPEAII I REMMAHU LDLAWRCRVITAPIL I RI MLEREEHIVTAYE NLRRWRMTNSRETCA MLFSEIERUREDNEU EOOAEMTIPOVPIRQ NORWZPEPRSIABES T T C R E ERPES RRBDC S OEERRUETI DWLLR ETPNSNARNCTEEOI KOSCCFI LESURRSB I MEHTALXCKNIFEE
DNEEDLENOSEWOTR

Components and Parts Answer
Here is the solution to the Components and Parts puzzle which appeared on page 64 of the June 1976 BYTE.


## IMSAI

 announces a unique 4K RAM board for just$\$ 139$.
Nobody has a 4K RAM board that gives you so much for your money. It's fully compatible with the Altair 8800.

Through the front panel or under software control, you can write protect or unprotect any 1 K group of RAM's. Also under software control you can check the status of any 4K RAM board in 1 K blocks to determine whether it's protected or not. The board has LED's that clearly show you the memory protect status of each 1 K block and which block is active. And there's a circuit provided that will let you prevent the loss of data in the memory if there's a power failure. This low power board has a guaranteed 450 ns cycle timeno wait cycle required. There's nothing like the IMSAI 4K RAM board around.

Dealer inquiries invited.


IMS Associates, Inc.


## Put the "Do Everything"

## Chip in Your Next Design

## Robert Baker 34 White Pine Dr Littleton MA 01460

The TMS-5501 is a multifunction, 10 controller available from Texas Instruments in a 40 pin dual in line package. It provides an asynchronous communications interface with programmable data rate, a parallel data 10 buffer, interrupt control logic, and five programmable interval timers in a single chip. Although designed specifically for use with a TMS 8080 central processor, it may be used with almost any other microprocessor as described in this article. Figure 1 shows a functional block diagram of the TMS-5501 and table 1 lists the pin assignments along with a brief description of each pin.

The system data bus is used for data transfers between the TMS-5501 and the microprocessor as controlled by the interrupt, chip enable, sync, and address lines. A
convenient method of addressing the TMS-5501 is to connect the chip enable to the high order address line and the four TMS-5501 address inputs to the four lower order address bits of the address bus, limiting the system to 32 K bytes of memory space (on an 8080 based system). Alternately, a full address decode could be done with additional hardware logic. The four memory address inputs of the TMS-5501 are then used to select one of the 14 possible commands shown in table 2. The individual commands are actually generated by executing memory reference instructions, with the low order hexadecimal memory address being the TMS-5501 command. This provides great flexibility by allowing the use of any memory reference instruction for 10 operations.


Figure 1: Block Diagram of the Texas Instruments TMS-5501 Design. This diagram shows the elements which are contained within this "do everything" IO chip. Notations next to the interconnection lines specify the number of bits involved in the data path.

Table 1: TMS-5501 Pin Assignments and Functions. This information is taken directly from the Texas Instruments Incorporated TMS-5501 Multifunction Input/Output Controller Manual.

## Command Functions

Addressing the "Issue Discrete Command" function (storing data in memory location XXX4) will cause the TMS-5501 to decode the data bus information as shown in figure 2 . Bits 1,2 , and 3 are latched until a new discrete command is received. Setting the RESET bit low has no effect, while setting it high will:

1. Clear the receiver flags in the TMS-5501 but not the receiver buffer.
2. Set the transmitter data output high (marking) and set the XMIT BUFFER EMPTY bit high to indicate the buffer is ready to accept a character from the processor.
3. Clear the interrupt register except for the transmitter buffer interrupt.
4. Inhibit interval timers.

Setting the Interrupt Acknowledge Enable bit high will allow the TMS-5501 to decode the processor status (bit 0 of the data bus at SYNC of each 8080 machine cycle) to determine if an interrupt acknowledge is being issued by the processor and initiate appropriate interrupt operation. Otherwise, with the bit set low, the TMS-5501 will ignore the interrupt acknowledge. Bits 4 and 5 are only used for testing purposes and should always be kept low.

Reading the TMS-5501 status (memory reference to address XXX 3 ) will return a data byte indicating the status of the controller as shown in figure 3. A brief description of each bit follows:

Bit 0 - A high indicates a framing error detected on the last received character.
Bit 1 - A high indicates a new character was loaded into the receiver buffer before the previous character was read.
Bit 2 - Normally high when no data is being received.
Bit 3 - A high indicates a new character is in the receiver buffer.
Bit 4-A high indicates an empty transmitter buffer ready to accept a character.
Bit 5 - A high indicates one or more

| Low Order | Binary Address |  |  |  | Command | Function |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (Hexadecimal) | A3 | A2 | A1 | A0 |  |  |
| 0 | L | L | L | L | Read receiver buffer | $\mathrm{RBn} \rightarrow \mathrm{Dn}$ |
| 1 | L | L | L | H | Read external inputs | XIn $\rightarrow$ Dn |
| 2 | L | L | H | L | Read interrupt address | RST $\rightarrow$ Dn |
| 3 | L | L | H | H | Read TMS 5501 status | (Status) $\rightarrow$ Dn See figure 3 |
| 4 | L | H | L | L | Issue discrete commands | See figure 2 |
| 5 | L | H | L | H | Load rate register | See figure 4 |
| 6 | L | H | H | L | Load transmitter buffer | $\mathrm{Dn} \rightarrow \mathrm{TBn}$ |
| 7 | L | H | H | H | Load output port | $\mathrm{Dn} \rightarrow \overline{\mathrm{XO}} \mathrm{n}$ |
| 8 | H | L | L | L | Load mask register | $\mathrm{Dn} \rightarrow \mathrm{MRn}$ |
| 9 | H | L | L | H | Load interval timer 1 | Dn $\rightarrow$ Timer 1 |
| A | H | L | H | L | Load interval timer 2 | Dn $\rightarrow$ Timer 2 |
| B | H | L | H | H | Load interval timer 3 | Dn $\rightarrow$ Timer 3 |
| C | H | H | L | L | Load interval timer 4 | Dn $\rightarrow$ Timer 4 |
| D | H | H | L | H | Load interval timer 5 | Dn $\rightarrow$ Timer 5 |
| E | H | H | H | L | No function |  |
| F | H | H | H | H | No function |  |

Table 2: TMS-5501 Commands. The commands presented to this device are best implemented by memory references to addresses $X X X O$ through $X X X F$ (hexadecimal) where $X X X$ is an arbitrary location in memory address space. The low order hexadecimal address for each command, along with the binary code presented to the address lines. The notation Dn refers to data bus line " $n$ " where $n$ is 0 to 7 .


Figure 2: Discrete Command Format. The data bus contents stored in the TMS-5507 when address location $X X X 4$ receives processor data is interpreted using this format.
interrupt conditions have occurred and the corresponding interrupt is enabled.
Bit 6 - A high indicates detection of the first data bit of a receive-data character.
Bit 7 - A high indicates detection of the start bit of an incoming data character.

Addressing the load rate register function (storing data in memory location XXX5) will allow the TMS-5501 to load the internal data rate register from the data bus. The bit assignments of the rate register are illustrated in figure 4 . Bit 7 selects the desired number of stop bits while setting a bit between bits 0 and 6 selects the transmitter/ receiver baud rate between 110 and 9600 baud. If more than one rate bit is high, the highest rate will be selected. If bits 0 to 6 are all low, the receiver and transmitter circuitry will be disabled.

Reading the receiver buffer (address XXX0) or external inputs (address $X X X 1$ ) will return the data received from the external devices while loading the transmitter buffer (address XXX6) or output port (address XXX 7 ) will output data to the external devices. The load mask register (address XXX8) function loads the data byte into an interrupt mask register where a high in bit " $n$ " enables interrupt " $n$ ". See table 3 for a list of the possible interrupts, their priorities, and information on the corresponding RST instruction for an 8080 processor. Addressing the read-interrupt address function will return the highest priority interrupt address
(the RST instruction) and clear the corresponding bit in the interrupt register. The TMS-5501 status should be checked for an interrupt pending before trying to read an interrupt address to avoid false data.

Loading an interval timer loads the contents of the data bus into the corresponding interval timer and activates that particular timer. The timer then counts down in increments of 64 microseconds providing a programmable interval of 64 to 16,320 microseconds with longer intervals generated by cascading timers through software. When a timer reaches zero, it generates an interrupt which in turn can control any desired function such as scanning a keyboard or switch matrix, time slicing in a multiprogramming environment, etc. Loading a new value while a timer is counting overrides the previous value and the timer starts counting down the new value. If a zero value is loaded to a timer, an interrupt is generated immediately.

## Interrupt System

The TMS-5501 provides several interrupt control functions by receiving external interrupt signals, generating interrupt signals to the processor, masking out undesired interrupts, establishing priority of interrupts, and generating RST instructions for an 8080 processor. External interrupts are normally received on pin 22, the SENS input, but an additional external interrupt can be received at the X 17 input pin if selected by a discrete command. The TMS-5501 will generate an interrupt to the processor whenever any of the five internal interval timers counts to

BIT:

| $7$ <br> START BIT DETECT | $\begin{gathered} 6 \\ \text { FULL } \\ \text { BIT } \\ \text { DETECT } \end{gathered}$ | 5 <br> INTRPT PENDING | $\begin{gathered} 4 \\ \times \mathrm{MIT} \end{gathered}$ <br> BUFFER EMPTY | 3 RCV BUFFER LOADED | $\begin{gathered} 2 \\ \text { SERIAL } \\ \text { RCVD } \end{gathered}$ | $\begin{gathered} 1 \\ \text { OVERRUN } \\ \text { ERROR } \end{gathered}$ | $0$ <br> FRAME ERROR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |

Figure 3: Status Word Format. The data read into the processor when TMS-5501 address location $X X X 3$ is referenced is in this format.

zero, the receiver buffer is loaded, or the transmitter buffer is empty. When an interrupt signal is generated, it is compared with the mask register. If that particular interrupt is enabled, it is passed on to the priority logic which allows the highest priority interrupt to generate an RST instruction to the 8080 processor only if no higher priority interrupt is still pending.

Two methods of servicing interrupts are provided: an interrupt driven system utilizing the RST instructions or a polled interrupt system utilizing bit 5 of the TMS-5501 status byte (the interrupt pending bit). In an interrupt driven system, the INT output signal of the TMS-5501 is connected to the INT input of an 8080 processor. The normal interrupt sequence would be as follows:

1. TMS-5501 reçeives an external inter-
rupt or generates an internal interrupt signal and sets the appropriate RST instruction.
2. The INT output goes high signaling the processor that an interrupt has occurred.
3. If the 8080 is enabled to accept interrupts, it sets the INTA (interrupt acknowledge) status bit high at SYNC time of the next machine cycle.
4. If the TMS-5501 has previously received an interrupt acknowledge enable command from the processor (see bit 3 of discrete command), the RST instruction will be transferred on the data bus to the CPU.

In a polled interrupt system, the INT output is not used and the interrupt sequence would be as follows:

Figure 4: Data Rate Command Format. The data bus contents stored in the TMS-5501 address location $X X X 5$ has this format, and is used to control the internal data rate generator for serial communications. A high level on the appropriate selection bit sets that data rate unless another bit at a higher rate is also selected.

## MODEL CC-7 SPECIFICATIONS:

A. Recording Mode: Tape saturation binary. This is not an FSK or Home type recorder. No voice capability. No Modem. (NRZ)
B. Two channels (1) Clock, (2) Data. OR, Two data channels providing four (4) tracks on the cassette. Can also be used for Bi-Phase, Manchester codes etc.
C. Inputs: Two (2). Will accept TTY, TTL or RS 232 digital.
D. Outputs: Two (2). Board changeable from RS 232 to TTY or TTL digital.
E. Runs at 2400 baud or less. Synchronous or Asynchronous. Runs at 4800 baud or less. Synchronous or Asynchronous. Runs at 3.1"/sec. Speed regulation $\pm .5 \%$
F. Compatability: Will interface any computer or terminal with a serial I/O. (Altair, Sphere, M6800, PDP8, LSI 11, IMSAI, etc.
G. Other Data: (110-220 V), ( $50-60 \mathrm{~Hz}$ ); 3 Watts total; UL listed 955D; three wire line cord; on/off switch; audio, meter and light operation monitors. Remote control of motor optional. Four foot, seven conductor remoting cable provided. Uses high grade audio cassettes.
H. Warrantee: 90 days. All units tested at 300 and 2400 baud before shipment. Test cassette with 8080 software program included. This cassette was recorded and played back during quality control.
ALSO AVAILABLE: MODEL CC-7A with variable speed motor. Uses electronic speed control at $4^{\prime \prime} / \mathrm{sec}$. or less. Regulation $\pm .2 \%$
Runs at 4800 baud Synchronous or Asynchronous without external circuitry.
Recommended for quantity users who exchange tapes. Comes with speed adjusting tape to set exact speed.

## DIGITAL DATA RECORDER \$149.95

FOR COMPUTER or TELETYPE USE Any baud rate up to 4800


Uses the industry standard tape saturation method to beat all FSK systems ten to one. No modems or FSK decoders required. Loads 8 K of memory in 17 seconds. This recorder, using high grade audio cassettes, enables you to back up your computer by loading and dumping programs and data fast as you go, thus enabling you to get by with less memory. Can be software controlled.

Model CC7. . . $\$ 149.95$
Model CC7A. . $\$ 169.95$
NATIONAL multiplex
CORPORATION

NEW - 8080 I/O BOARD with ROM. Permanent Relief from "Bootstrap Chafing"
This is our new "turnkey" board. Turn on your Altair or Imsai and go (No Bootstrapping). Controls one terminal (CRT or TTY) and one or two cassettes with all programs in ROM. Enables you to turn on and just type in what you want done. Loads, Dumps, Examines, Modifies from the keyboard in Hex. Loads Octal. For the cassettes, it is a fully software controlled Load and Dump at the touch of a key. Even loads MITS Basic. Ends "Bootstrap Chafe" forever. Uses 512 bytes of ROM, one UART for the terminal and one USART for the Cassettes. Our orders are backing up on this one. No. 2SIO (R)

Kit form \$140. - Fully assembled and tested \$170.00

Send Two Dollars for Cassette Operating and Maintenance Manual with Schematics and Software control data for 8080 and 6800. Includes Manual on I/O board above. Postpaid

Master Charge \& BankAmericard accepted.
On orders for Recorders and Kits please add $\$ 2.00$ for Shipping \& Handling.
(N.J. Residents add 5\% Sales Tax)

3474 Rand Avenue, Box 288
South Plainfield, New Jersey 07080
(201) 561-3600

| Data Bus Bit |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :--- |
| Priority | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ | $\mathbf{7}$ | Interrupt Caused By |
| $\mathbf{1}$ | H | H | H | L | L | L | H | H | Interval Timer 1 |
| 2 | H | H | H | H | L | L | H | H | Interval Timer 2 |
| 3 | H | H | H | L | H | L | H | H | External Sensor |
| 4 | H | H | H | H | H | L | H | H | Interval Timer 3 |
| 5 | H | H | H | L | L | H | H | H | Receiver Buffer |
| 6 | H | H | H | H | L | H | H | H | Transmitter Buffer |
| 7 | H | H | H | L | H | H | H | H | Interval Timer 4 |
| 8 | H | H | H | H | H | H | H | H | Interval Timer 5 or X17 |

Table 3: Interrupt Assignments. This table lists the priority level and source of each interrupt. For 8080 systems without previously dedicated restart instructions, the TMS-5501 can be used to automatically generate RST $n$ where $n$ is the priority level. Otherwise, some form of polling is required.

1. The TMS-5501 receives or generates an interrupt and sets the corresponding RST instruction.
2. The TMS-5501 interrupt pending status bit (see bit 5 of status byte) is set high.
3. At some predetermined time, the processor polls the TMS-5501 by reading the status byte and checking the interrupt pending bit.
4. If the bit is high, the processor should then read the interrupt address which will cause the TMS-5501 to transfer the RST instruction to the processor as data.
5. The software can then decode the RST instruction to determine the interrupting device.
With an 8080 processor, the TMS-5501 controller may be used in either interrupt mode. For any other microprocessor based system, the TMS-5501 would be easiest to implement using a polled interrupt system that would not require any extra hardware to interface with the CPU. An alternate method would use part of both interrupt schemes by connecting the INT output to the CPU interrupt input but disabling the TMS-5501 interrupt acknowledge enable bit (bit 3 of discrete commands = low). This would eliminate the need to poll the TMS-5501 status byte but the interrupt address (the RST instruction) would still have to be read and decoded by the CPU.

## Applications

The TMS-5501 appears to be an ideal interface for the typical personal computer system with many different possible applications. For example, suppose we connected a CRT to the external output port and an ASCII keyboard to the external input port. The serial 10 port could then be used with a Teletype for hardcopy, an accoustic coupler for loading from another system (or intersystem communications), or for a cassette drive. You may want to connect a switch to the SENS input to cause an interrupt that would return to a system control or debug routine.

For an ideal cassette interface, the serial 10 could transfer data to and from the cassette deck while the external outputs could be used to control drive functions such as forward, reverse, fast forward, etc. The external inputs might be used for inputting status information from the drive such as end of tape, start of tape, tape jam, etc. With this scheme, true searches and tape scans could be accomplished under full software control with no manual actions required once a tape is inserted in the drive.

As you can clearly see, there's an endless number of possibilities for the TMS-5501 limited only by your imagination. Even if your system is not based on an 8080 processor, it is well worth looking at the TMS-5501 for your IO needs.■

## Election Program

The following is paraphrased from a version submitted by Mark T O'Bryan of Portage MI, who credits it to a friend at MIT.

Question: "What's the biggest problem in writing
$a\left\{\begin{array}{l}\text { Jerry Ford } \\ \text { Scoop Jackson } \\ \text { Jimmy Carter } \\ \text { Ronald Reagan }\end{array}\right\} \begin{aligned} & \text { simula tor for an Altair } \\ & \text { with only 4096 bytes } \\ & \text { of memory? }\end{aligned}$

Answer: Trying to figure out what to do with the other 4095 bytes of memory. $\quad$

## New Hope For Computer Phreaques

According to a report published in Electronic Products magazine, April 1976, page 16, Commander Thomas Orr of Warsash, England, has created a digital electronic watch which reads the pulse rate of its wearer. Just think, now you can check to see if you're dead or not. The watch is marketed in the United States by Pulse Watch of Tiburon CA and the transducer technology on which it is based has much wider applications in medical electronics.


Now, you can buy an Altair 8800 or Altair 680 computer kit right off the shelf. Most all Altair options, software and manuals are also available. The MITS Dealer List below is just the beginning:

RETAIL COMPUTER STORE, INC.
410 N. E. 72nd Seattle, WA 98115 (206) 524-4101

COMPUTER KITS 1044 University Ave. Berkeley, CA 94710 (415) 845-5300

CTI DATA SYSTEMS 3450 East Spring St. Long Beach, CA 90806 (213) 426-7375

2839 W. 44th Ave.
Denver, CO 80211
(303) 458-5444

ARROWHEAD COMPUTER CO. (THE COMPUTER STORE) 820 Broadway
Santa Monica, CA. 90401

GATEWAY ELECTRONICS
8і23-25 Page Blvd. St. Louis, MO 63130 (314) 427-6116


BYTE'TRONICS
5604 Kingston Pike Knoxville, TN 37919 (615) $588-8971$
the COMPUTER STORE, INC. 120 Cambridge Burlington, Mass. 01803 (617) 272-8770

THE COMPUTER SYSTEMCENTER 3330 Piedmont Road Atlanta, GA 30305
(404) 231-1691

MICROSYSTEMS 6605A Backlick Rd. Springfield, VA 22150 (Washington DC area) (703) 569-1110


MARSH DATA SYSTEMS
5405-B Southern Comfort BIvd.
Tampa, FL 33614
(813) 886-9890

THE COMPUTER ROOM 3938 Beau D'Rue Drive Eagan, Minn. 55122
(612) 452-2567

## Why Wait?

## Build a FAST Cassette Interface

Dr Robert Suding<br>Research Director for Digital Group Inc<br>PO Box 6528<br>Denver CO 80206

This cassette interface does not have a $\pm 30 \%$ speed tolerance. The design requires $\pm 12 \mathrm{~V}$ and +5 V to run. A good quality recorder must be used, along with excellent quality tapes. Careful adjustments are required.

So why use it? Well, it works! It's dependable. And it's fast. In contrast, the proposed BYTE standard cassette interface runs at 300 Baud. A Teletype paper tape reads @ 110 Baud. I have 24 K on my system. How long would it take me to completely load my system (not including any Bootstrap Loader operations)?

Teletype @ 110 Baud - 40 minutes 58 seconds
Proposed BYTE standard @ 300 Baud 15 minutes 1 second

The system to be shown in this article has been running for almost a year at 1100 Baud (with an upper limit of 1750 Baud with critical tuning).

Suding system @ 1100 Baud - 4 minutes 6 seconds

Past issues of BYTE have included several articles on cassette interface proposals and
circuits. I would suggest re-reading these articles. You will find one common element. Slow. If you get the impression that I'm impatient, you're right. I'll bet you are too. Imagine reading 300 Baud for 15 minutes to discover a noise pulse had destroyed data, requiring re-reading. Ugh!

Thus the proposed standard of the BYTE Kansas City conference in 1975 has a major disadvantage: The use of a redundant Manchester format with a 1200 Hz low frequency critically restricts the user to slower data rates. A related disadvantage for those who use filters or phase lock loops as an input detection method is the fact that the Manchester code employs harmonically related frequencies; this leads to design problems in detectors based upon frequency discrimination techniques.

The system shown in this article avoids the above pitfalls. It uses the non-harmonically related tones of 2125 Hz - Mark and 2975 Hz - Space. The second harmonic of 2125 Hz occurs at 4250 Hz , well down on the passband of a 2975 Hz detector. Sufficient space exists between the two frequencies to allow for reasonable recorder speed discrepancies. The higher frequencies involved permit increasing the data rate.

Several approaches are possible in cassette interfacing, as seen in past BYTE articles. However, their emphasis on wide cassette speed tolerance made them slower. My


Figure 1: The schematic of the Suding cassette input interface as found in the Digital Group systems. This interface amplifies and clips the cassette output with limiting amplifier IC34, discriminates the two data frequencies (see table 1) with bandpass filters followed by full wave detectors, passes the detected signal through a 3 pole active low pass filter, then converts the result to a TTL level which is read by a single bit input port. One example of software (see listing 1) to drive this input interface uses a programmed simulation of UART input algorithm; an actual UART or ACIA device could be substituted if desired.
approach to "out of specification cassette speed" is - "put it in the specification, or get a good recorder." More of that later.

## Theory of Operation

The 1100 Baud Digital Group system uses the circuits of figures 1 and 2 . The cassette receive circuitry detects the prerecorded frequency shift keying and produces a " 1 " or a " 0 " output as a result of a detected 2125 Hz or 2975 Hz tone at the input. A 741 operational amplifier, IC34, is used as a clamped limiter which prevents variations in cassette amplitude from affecting the detection process. The output of the limiter should be about .6 V peak to peak, roughly a square wave with rounded edges of the incoming frequency, constant in amplitude regardless of tape volume setting or minor tape "dropout" problems.

Two bandpass active filters (IC35) then amplify a tone five times when actually tuned to their respective frequencies of 2975 Hz for the top filter, and 2125 Hz for the lower filter. The further off the tuned frequency the tone is, the less amplification the filter will produce. The gain, bandwidth, and tuned frequency are set by the three resistors and two condensers in each filter. Each filter may be exactly tuned to frequency by carefully setting the variable resistance value (which may be either a potentiometer or selected fixed values).

Full wave active detectors produce rectified full wave pulses at the summing junction, pin 5 of IC37. The 2975 Hz tones are rectified to a positive voltage, and the 2125 Hz tones are rectified to a negative voltage. As received tones depart from either exact frequency, a value less positive or


Figure 2: The schematic of the Suding cassette output interface as found in the Digital Group systems. The output interface is a simple audio frequency shift keyer made up of a 566 voltage controlled oscillator with two frequency states controlled by a single TTL data line. The TTL level which drives the output modulator is a single bit derived from an output port. The software (see listing 2) to drive this output interface is shown as a programmed simulation of a UART output algorithm; an actual UART or ACIA device could be substituted if desired.
negative is produced until approximately midway ( 2550 Hz ) a summed voltage of 0 results.

A three pole lowpass active filter then removes the remaining traces of pulsating DC from the summed signal with almost no effect on the data pulses up to a speed of 1000 bits per second. If lower data rates were to be utilized, an improved signal to noise ratio could be obtained by multiplying the values of C12, C13, and C11 by the reciprocal of the data rate ratio. Table 1 shows some component values for alternative frequency designs.

The final receiver section is a 741 operational amplifier, IC38, connected as a slicer. This operational amplifier detects whether the voltage at its pin 2 is positive or negative with respect to the constant voltage at its pin 3. The output voltage will then swing either to nearly -12 V or to nearly +5 V . Notice that this operational amplifier has +5 as its positive supply voltage, pin 7. A forward biased germanium diode prevents the actual output voltage from going less

## Tune Up Notes

The cassette interface must be carafully tuned to achieve proper performance. Careless tuning has been the most frequent cause of cassette system failure.

1. Plug in the six integrated circuits of the cassette interface.
2. Connect a calibrated audio oscillator between the limiter input and ground. A digital frequency counter driven by the audio oscillator is highly recommended. The oscillator should cover the desired range of $2-3 \mathrm{kHz}$, with a sine wave output of .5 or so, although the precise level is not at all critical.
3. Apply +5 and $\pm 12$ voltages to the circuit. Measure the output at pin 6 of the 741 limiter (IC34) with an oscilloscope. The wave shape should be a rounded square wave of about .6 V peak to peak.
4. Set the audio oscillator to 2125 Hz . Measure the output at pin 1 of the 5558 active bandpass filter. Slowly turn R25 until the signal peaks. Be sure that you are peaking at 2125 Hz , not a harmonic. Vary the oscillator frequency a few decades to insure 2125 Hz is the tuned frequency.
5. Similarly, set the oscillator to 2975 Hz and measure the output at pin 7 of the 5558 (IC35). Slowly turn R26 until the signal peaks. Vary the oscillator to insure a 2975 Hz peak.
6. Measure the detected voltages at pin 5 of IC37. When the oscillator approaches 2125, the voltage should go negative. When approaching 2975, the voltage should go positive. Trouble in this area would most likely be caused by reversed or defective diodes, or shorts between adjacent lines.
7. Measure the voltage at the cathode (bar) end of the output clamping germanium diode
(G1). Sweeping the frequency between 2125 and 2975 Hz should result in a clean voltage jump somewhere between 2125 and 2975 Hz . Measure the output swing to insure that it does not exceed $+5,-.3 \mathrm{~V}$.
8. Remove the audio oscillator and short input connector J1 temporarily to ground. Measure the output at pin 6 of IC34. A stable condition (no oscillation) should be seen. Connect the oscilloscope to the cathode of G1 again. Adjust the balance potentiometer (R42) so that the output voltage is a negative level. Slowly turn the potentiometer until the output voltage jumps to a positive level and leave the setting at this point.
9. Disconnect the temporary jumper from the input connector and reconnect the audio oscillator. Perform step 7 again. The crossover threshold should be close to 2550 now. If all proceeds well at this point, the cassette interface is ready to receive data.
10. Connect the oscilloscope to pin 4 of the 566 voltage controlled oscillator (IC33). A triangular wave output should be seen.
11. Connect a temporary jumper between the TTL input going to DS1 and +5 V . Connect a frequency counter to pin 3 of the VCO (IC33). Adjust potentiometer R41 for a resultant output frequency of 2125 Hz .
12. Remove the jumper from +5 V and connect the jumper from DS1's input to ground. This time adjust R40 for 2975 Hz output.
13. Remove the jumpers, and you are ready for final tune in the driving circuit. Connect the cassette interface to the driving output port, and program the driving processor to send a TTL high level (" 1 ") output to the cassette interface. Adjust R41 to 2125 Hz . Then have the processor send a " 0 " level. This time adjust R40 for 2975 Hz output. The cassette interface is now ready for use.
than $\simeq-.2 \mathrm{~V}$, so that valid TTL levels are not exceeded. An offset adjusting potentiometer allows the output to be placed in a "Mark Hold" condition when no tone input is being detected.

The cassette recording section (figure 2) uses a single integrated circuit, a 566 voltage controlled oscillator, IC33. A logic level from the computer's output port controls the resultant audio frequency output to the cassette recorder microphone input. A high input (" 1 ") produces a 2125 Hz output, and a low input ("0") results in 2975 Hz . The output wave shape is a symmetrical triangular wave. Should the user object to using a triangular wave, a more nearly sine wave can be obtained by connecting a pair of back to back 1 N914 diodes between ground and the output side of the coupling capacitor C5.

Exact values and high quality components will result in a trouble-free voltage controlled oscillator. The $47 \mathrm{~K}(\mathrm{R17})$ resistor in series with the output is a typical value to be used when coupling to the low level, low impedence external microphone inputs of most cassette recorders. Using the "AUX" input of your cassette recorder generally gives better results.

## Construction

The cassette interface is available as a part of a printed circuit board kit from the Digital Group. The printed circuit board is shared by a television display circuit to be described in the next article in this series. A kit of the cassette interface only is also available from the Digital Group for $\$ 30$, which includes all parts and the printed circuit board. The experienced builder can build the circuit in an evening or two by hand wiring components on standard .1 inch grid Vectorboard. All the circuitry can be contained in an area of approximately 3 inch by 5 inch (about 8 cm by 13 cm ).

Be sure to use only high quality components, particularly in the active bandpass filters and voltage controlled oscillator. Some strange "frequency jump" problems have been traced to surplus 566 s which were temperature sensitive. Lay out the receive circuit to avoid feedback paths from output to input, particularly in the limiter, active bandpass filters, and slicer areas. Different op amps could be used, but may result in instability or degradation of final performance due to suboptimization.

## Modifying Your Cassette Recorder

It is very helpful to listen to the data from the cassette so that the beginning of the data burst may be detected, as well as


VIDEO TERMINAL INTERFACE
Connects to standard TV monitor (or modified receiver) to display 16 lines of 64 or 32 characters in a $7 \times 9$ matrix. Character set includes 128 upper and lower case ASCII characters and 64 graphic characters for plotting $48 \times 64$ ( 128 with memory option) array. Text and graphics may be mixed on the same screen. 8 -bit keyboard input port is provided. Characters are stored in the onboard memory and may be read or written by the computer. Cursor control, text editing and graphics sottware is included. Fully compatible with Altair and IMSAI. Sockets included.
$\mathrm{VTI} / 3232$ character line $\$ 185.00 \mathrm{kit}$
VTI/64 64 character line $\$ 210.00$ kit $\$ 285.00$ assd.


## ANALOG INTERFACE

Complete interface for a CRT graphics display or X-Y plotter. Provides 8 channels of software-controlled A/D conversion. Yor 2 channels of analog output with 10 -bit resolution ( $0-10 \mathrm{v}$ or $\pm 5 \mathrm{v}$ out), 6 bits of latched digital output, and 8 analog comparators. Software included for A/D conversion by successive approximation and tracking. Sockets included.
ADA/1 1 analog output $\$ 145.00 \mathrm{kit}$ ADA/2 2 analog outputs $\$ 19500 \mathrm{kit} \$ 255.00$ assod.

## SEE THESE AND OTHER

PRODUCTS AT THE FOLLOWING COMPUTER STORES
THE COMPUTER STORE • (213) 478-3168
Los Angeles, CA
THE COMPUTER MART (714)633-1222
Orange, CA
COMPUTERS \& STUFF (415) 278-4720
San Lorenzo, CA
COMPUTERS \& STUFF (801) 377-1717
Provo, Utah
COMPUTER SYSTEMS CENTER (404) 231-1691
Atlanta, GA
BYTE SHOP (415) 969-5464

Mountain View, CA

## DEALER INQUIRIES INVITED

all prices and specilications subplect to change without notice Prices are USA only POLYMORPHIC SYSTEMS
737 S. Kellogg, Goleta, CA 93017 [805] 967-2351 SankAmericard and Master Charge Accepled

|  | Low Filter |  |  | High Filter |  |  | Low Pass Filter |  |  | Vco |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | R21 | R24 | R25 | R22 | R23 | R26 | C13 | C12 | C14 | R12 | R15 |
| $\begin{aligned} & 2125-2975 \mathrm{~Hz} \\ & 1100 \text { Baud } \end{aligned}$ | 6.8 k | 68 k | 938 | 4.7 k | 47 k | 697 | . $0056 \mu \mathrm{~F}$ | . 01 | . 015 | 2.7 k | 1.3 k |
| $1200-2400 \mathrm{~Hz}$ 300 Baud (Simple) | 6.8 k | 68 k | 4173 | 4.7 k | 47 k | 1162 | . $0056 \mu \mathrm{~F}$ | . 01 | . 015 | 470 k | 2.7 k |
| $1200-2400 \mathrm{~Hz}$ 300 Baud (Correct) | 12 k | 120 k | 1668 | 5.6 k | 56 k | 906 | . $015 \mu \mathrm{~F}$ | . 033 | . 047 | 470 k | 2.7 k |
| $\begin{aligned} & 2125-2295 \\ & 100 \text { Baud } \\ & \text { (Simple) } \end{aligned}$ | 6.8 k | 68 k | 938 | 4.7 k | 47 k | 1301 | . $0056 \mu \mathrm{~F}$ | . 01 | . 015 | 47 k | 2.7 k |
| $\begin{aligned} & 2125-2295 \\ & 100 \text { Baud } \\ & \text { (Correct) } \end{aligned}$ | 36 k | 360 k | 156 | 27 k | 270 k | 179 | . $056 \mu \mathrm{~F}$ | . 1 | . 15 | 47 k | 2.7 k |

* means that the value so indicated is the typical calculated value. The precise value is dependent on component tolerance.

Table 1: Theoretical values of components for alternate frequencies. This table gives values of components to be used with the circuits of figures 1 and 2 in order to make this cassette interface work with several alternate specifications. See the text for a definition of the various comments at the left of the table.

## Potential Troubles

Knowing about potential problem areas is a first step to minimization of their effects. Troubles seem to break down into six classes.

- Cassette recorders and the cassettes used: A marriage betweenyour $\$ 1000$ microprocessor and junior's \$20 cassette recorder, which has been using $30 d$ cassettes for the last five years, will not produce happy offspring! I have been using a Superscope C-104 for the past year, and can report no failures except for defective cassette tapes. The $\mathrm{C}-104$ has several attractive features. Besides the usual conveniences such as index counter, cuing, etc, it has a variable readback speed control, dandy for out of spec cassettes from friends. Inside, another special motor speed control potentiometer is located near the speaker which allows precisely setting the record/write speed. Quality control seems good overall, and the list price of $\$ 120$ (cheaper at discount stores) is worth the investment. Don't waste your money on cheap cassettes. Sony Low Noise C-45s have been generally good. Some \$2 - \$4 Data Certified Cassettes are superior, but not needed.
- Microprocessor caused problems: Some microprocessor designs will not work directly with this interface system. This interface was designed to be connected directly to a single bit 10 port, with the processor handling all of the bit timings through timing loops. If your processor must periodically catch its breath for such things as dynamic memory refreshing, you may be unable to directly use the "Software UART" system. What a shame! However, a hardware UART will permit using the system even with a system of this nature. - Cabling problems: It is possible to connect your cassette recorders with the read and write cables reversed. Enough crosstalk from the write line to the read limiter existed to give the appearance of data being read, but so many errors resulted that the programming would not run. - Tuning problems: Circuit tuning is the most common problem. Carefully tune the active filters! - Cassette Crashes: Cassette damage is frequent
on tapes which have always worked before, but now mysteriously fail. The most common cause of this is removing a cassette from the recorder without completely rewinding. The exposed oxide then gets damaged, and is no longer usable.
- Miscellaneous circuit problems:

Defective level output from cassette read limiter.

1. None at all: Check for $\pm 12 \mathrm{~V}$ to IC34, and IC34.
2. Too high output level: Diodes (DS4 and DS5) open, or one is reversed.
Bandpass active filters don't filter.
3. Off frequency
4. Bad 5558
5. Check for shorts or out of tolerance condensers C8, C9, C10, or C11. Disk ceramics are a "no-no" in tuned circuits.
6. Resistors improperly wired or inserted.

Full wave detector does not work as described:

1. Diodes open, reversed or shorted.
2. Defective IC36.

Low pass active filter fails to work:

1. Shorted or out of tolerance condensers.
2. Defective IC37.

Output slicer (IG38) fails to produce TTL levels:

1. Reversed, open or not Germanium diode at DG1.
2. Too heavily loaded output. This circuit should drive no more than one TTL load (standard for most 10 ports).
VCO won't oscillate.
3. Defective 566 (IC33).
4. Shorted condenser C6.

VCO has parasitic oscillation (high frequency):

1. $C 7$ not connected.
2. Defective 566.
3. C6 is open, producing a very high frequency.
VCO won't tune to frequency or stay there:
4. Out of tolerance or defective C6. You really didn't use a disk ceramic here, did you?
5. Defective 566.
6. Non-TTL levels used to drive VCO.
7. Defective potentiometers R40 or R41.
8. DS1 or DZ2 reversed or defective.
hearing the end of the data. When the cassette read cable is plugged into most cassette recorders' earphone output jack, the speaker output is usually cut off. However, since a closed circuit jack is all that is involved, a quick solution is to connect a jumper on the jack so that the speaker is not disconnected. Even better, use a 100 ohm $1 / 4$ watt resistor instead of the jumper, and the data howl won't be so loud. A 10 ohm, $1 / 4$ watt resistor from the amplifier lead to jack, to the jack frame will prevent potential damage to the output driving transistor(s).

## Alternative Frequencies and Applications

The cassette interface design may be used with the proposed BYTE standard should you so desire. Table 1 has appropriate component values calculated for two alternative possibilities: the simple way (less desirable) and the "right way". The simple way permits using a switch on the bandpass active filters to select the frequency pairs. The right way involves setting the circuit to the optimal values, and using separate interfaces for each frequency pair.

Amateur radio (ham) radioteletype (RTTY) generally uses $2125-2295 \mathrm{~Hz}$ frequency shift keying for 170 Hz shift. The existing cassette interface can be used by "straddle tuning," but improved performance may be obtained by selecting a second R26 which will tune the high filter to 2295. The cassette read cable may then be attached to the short wave receiver and the microprocessor, programmed as a radioteletype video terminal, which can replace the noisy Teletype machine. Of course, a cassette interface specifically designed for this 170 Hz shift at 100 WPM will give superior performance under marginal conditions.

The cassette interface may be used as a stand alone radioteletype terminal unit and audio frequency shift keying if desired, and works quite nicely in this application.

## Software

I would suggest using software for your cassette read and write timings. Sample 8080 software is included as listing 1 . Timings at locations $\langle 0\rangle / 116,<0\rangle / 133,<0\rangle / 241$, and $<0>/ 260$ are based on an 8080 system with a 500 ns T time and no wait states. Slower systems will require proportionately decreased loop timings.

A UART could be used instead of the "software UART" system shown. However, several disadvantages arise. First, a slightly greater cost and complexity. More important, however, is a degradation in total

# If you want a microcomputer with all of these standard features 

- 8080 MPU (The one with growing software support) - 1024 Byte ROM (With maximum capacity of 4 K Bytes) - 1024 Byte RAM (With maximum capacity of 2 K Bytes)
- TTY Serial 1/O
- EIA Serial I/O
- 3 parallel l/O's - ASCII/Baudot terminal comterminal com-- Monitor having load, dump, display, insert and go functions
- Complete with card connectors - Comprehensive User's Manual, plus Intel 8080 User's Manual - Completely factory assembled and tested - not a kit
- Optional accessories: Keyboard/video display, audio to follow. The HAL MCEM-8080. $\mathbf{\$ 3 7 5}$


## ...then let us send you our card.

HAL Communications Corp. has been a leader in digital communications for over half a decade. The MCEM-8080 microcomputer shows just how far this leadership has taken us... and how far it can take you in your applications. That's why we'd like to send you our card-one PC board that we feel is the best-valued, most complete

microcomputer you can buy. For details on the MCEM-8080, write today. We'll also include comprehensive information on the HAL DS-3000 KSR microprocessorbased terminal, the terminal that gives you multi-code compatibility, flexibility for future changes, editing, and a convenient, large video display format.

HAL Communications Corp.
Box 365, 807 E. Green Street, Urbana, Illinois 61801 Telephone (217) 367-7373

Listing 1: Stand Alone Suding Cassette Input Program. This program is a self contained data transfer routine which will transfer a block of data from cassette to split octal memory locations $x x x / x x x$ through yyy/000. This program assumes that MEMTOCAS (see listing 2) was used to create the tape being read. A more generally useful input facility would be modelled on this program and linked to a system monitor as a subroutine.


## Notes:

- Input is assumed to be wired to bit 0 of port 1, from output of IC38 pin 6 via resistor R38 and shunted by diode DG1.
- Loading proceeds from split octal address $x \times x / x \times x$ to address yyy/000. Enter this program by jumping to location $<0>/ 100$ after setting up constants of address.
- "*" indicates a timing constant for the "software UART" inputs.
- " $v$ " indicates the end of transfer comparison mentioned in text.
- $<0>$ indicates an arbitrary page location for this program, to be replaced by a real memory page number when actually loading the program at byte 100 of some page.

Listing 2: Stand Alone Suding Cassette Output Program. This program is a self contained data transfer routine which will transfer a block of data from split octal memory locations $x x x / x x x$ through yyy/000 onto cassette tape after a five second leader output delay. This program assumes that CASTOMEM (see listing 1) will be used to read the tape being created. A more generally useful output facility would be modelled on this program and linked to a system monitor as a subroutine.


## Note:

- Output is assumed to be wired from bit 0 of port 1 to DS1 in figure 2.
- See notes to listing 1 for listing conventions.
system flexibility. The "software UART" allows the timing constants to be dynamically modified (if desired) by detecting the variations in the stop bit timing, thereby compensating for wow and flutter. Digital integration of the incoming data bits is possible by setting a register to octal 200 at the beginning of each bit time. During the bit time, repeated sampling either adds or subtracts from the register (depending on whether 1 or 0 ) and a "branch minus" instruction system effectively eliminates receive problems. This digital integration detection is utilized by the Digital Group Z-80 cassette read software.

Versions of this "software UART" system have been written for $8008,8080, \mathrm{Z}-80$, 6502 , and 6800 . All work satisfactorily.

## Operation

This cassette system is utilized by first turning on the cassette recorder and waiting until the lower tone 5 second leader tone is heard. At this point, restart the system to the beginning address of the "Cassette to Memory" software.

Cassette writing is accomplished by restarting the system to the beginning of the
"Memory to Cassette" programming. Be sure to set the appropriate start and stop addresses prior to beginning the read or write operations. The monitor programs in the various Digital Group systems automatically set the start and stop addresses. The check marks in the listing $(\sqrt{ })$ indicate the points where start and stop addresses may have to be modified.

The software may be adjusted to run at different data rates by changing the values at the addresses mark with an asterisk (*). Note that the constants at $<0>/ 133$ and $<0>/ 241$ are the same. The constant at $<0>/ 116$ is $50 \%$ greater and the constant at $<0>/ 260$ is twice the value of the constant at $<0>/ 241$.

## Summing It Up

This cassette interface represents a simple but fast and dependable way to store programs and data for the serious hobbyist. It does not seek to be all things to all users, but a number of applications can be run using the same basic design. The detail interface design has independence from other components in the system, allowing various processors to use the same cassette system (with appropriate software)..


## Surplus Electronics in Tokyo and Manila

Dr Michael N Hayes PO Box 167

Port Orchard WA 98366

Both Manila and Tokyo are characterized by a large number of very competitive shops which deal primarily in components.

Integrated circuits are a rarity in the surplus electronics markets. A shop in Manila keeps integrated circuits in a glass case similar to those found in a jewelry shop.

This article is a brief summary of some of my experiences, opinions, and observations while searching for and buying surplus electronic components in Tokyo, Japan, during October 1975; and in Manila, Republic of the Philippines, during November 1975. Perhaps this information will be of interest to US buyers and hobbyists for comparison purposes or in case they have the opportunity to visit these places.

In summary, the surplus electronics market in Tokyo is far better than that in Manila, but neither place offers the quantity of sophisticated computer electronics and integrated circuits available in the United States. The primary surplus electronic items sold in both cities are components: resistors, capacitors, switches, panel lights, terminals, connectors, transistors, etc. Component prices in Tokyo are roughly two thirds of US mainland prices, whereas in Manila the component prices are about the same as US prices. An important consideration here is that these components really aren't considered as surplus electronics in these cities, but are sold as retail electronics, subject to considerable bargaining, through hundreds of small shops and stands concentrated into specific areas of each city. Time did not permit investigation of industrial surplus outlets or auctions, but the impression formed is that the large numbers and competitiveness of these small stores preclude easy bargains from auctions or company sales.

Prices quoted in this article are of specific sampled items and are for illustration pur-


Photo 7: The author in front of a surplus outlet in Manila, Philippines.
poses only. All prices are in US dollars or cents with the following approximate exchange rates in effect: Tokyo, 300 yen per US $\$ 1$; Manila, 7.5 pesos per US $\$ 1$. There was a considerable disparity in prices among various stores; and thus, much like a marketplace, significant price savings could be found on selected items by spending the time to shop around. In many cases prices on selected items in one store, especially in Manila, might be half or double what was found in another store.

The real surplus electronics bargains of the Far East are rumored to be found in Taiwan, Hong Kong, and related areas. Perhaps a reader who has visited these areas could write in and supply additional findings.

## Tokyo, Japan

The entire Tokyo surplus electronics market, along with significant portions of the retail electronics market, is concentrated in an area called Akihabara (pronounced ah-kee-ha-ba-rah), northeast of the downtown area. To reach it, one simply takes one of the many trains available from the elaborate Tokyo train and subway system.

What is incredible about Akihabara is not so much what they have, as the quantity in which they have it. There are literally hundreds of tiny outlets here selling everything imaginable in the way of electrical and electronic goods: televisions, stereos, speakers, tape recorders, radios, ham equipment, light bulbs, tools, wire, refrigerators, air conditioners, etc. Much of the electronic market is transistor radios and calculators, but the number of components stores and amount of available "surplus" electronic components to be found far exceeds any comparable area I have seen in any US city.

Each of these shops has components in counters, bins, or stands by the hundreds, with small signs marking the types and prices. Bargaining is expected in most stores but frowned upon in a few. Most of these shop owners are hard bargainers, but in most cases a $10-20$ percent savings can result.

Here is a list of some sample items and prices in Tokyo:

Ribbon cable, 10 conductor, Spectra Strip - $10 \phi$ per foot.
Toggle switches - $75 \phi$ to $\$ 1$.
Small experimenter's PC cards (2 inch square) $-10 \phi$.
MC 7805s - $\$ 2.50$.
44 pin PC board edge connectors - $75 \phi$. Small capacitors - $5 \phi$ to $10 \phi$.
7400 s - $55 \phi$ (there are not really too many ICs available).
PC boards with components - $25 \phi$ to $50 \phi$ to $\$ 1$ (most have Japanese numbered transistors and ICs).
Resistors - $2 \phi$ to $5 \phi$.
General 50 pin connectors - $\$ 2$.
Crystals - small types - $\$ 1$.
Partial calculator PC boards - no cases older styles - \$2.
Partial cassette tape decks, audio, missing or damaged electronics - $\$ 1$ to $\$ 3$.
IC sockets - $30 \phi$ to $60 \phi$.

## Manila, Republic of the Philippines

The main area of Manila for surplus electronics is called Raon (pronounced rahown) Street and is located in a district near the downtown area called Quiapo (pronounced Kee-ah-po). The name of Raon Street, including its" signs, was recently changed to Gonzalo Puyat; but the Jeepney drivers still know it as Raon Street or Raon district. The best way to get to Raon is to take a Jeepney, which is any of the old WW II jeeps which have been painted and elaborately decorated and are now used as taxis throughout the Philippines. The ride will cost about $25 \phi$ from anywhere in the downtown Manila area. English is spoken and used extensively in the Philippines making it very easy for an American to get around and to be understood.

The Raon District has about 30 small shops which sell electronic parts in display cases and parts bins. The setup is similar to that in Tokyo, but Manila has only a small fraction (less than 10 percent) of the number of stores. The quantity of merchandise available in each store was considerably less than in the typical Tokyo store. Bargaining in Manila is expected and a 20 to 30 percent savings is not hard to get.

1 also found several stores in San Fernando, a small town about 40 miles west of Manila. These stores were very similar to the ones in Manila and seemed to support the local TV repair shops.

None of the shops in Manila or San Fernando had any computer equipment, and only a few of them had any integrated
circuits. The integrated circuits were usually only a handful, mounted in styrofoam and kept by the owner in a glass display case, like jewels in a jewelry store.

Here is a list of some sample items and prices in Manila:

Small capacitors - $10 \phi$ to $20 \phi$.
7400 s - $\$ 1$ to $\$ 3$ depending upon the store.
Alligator clips- 20\&.
Panel lights, small - 50 $\phi$ to $60 \notin$; no LEDs available.
Medium sized capacitors - $20 \nless$ to $30 \nless$.
Larger electrolytic capacitors - $\$ 1$ to \$1.50.
Small variable capacitors - 75 \& .
Power transistors (TO3), general types 20d to $30 \neq$.
Small potentiometers - $25 \not \subset$ to $50 \nless$.
Regular metal potentiometers - $50 \&$.
12 Watt transformers - $\$ 1.50$.
Small plastic transistors $-10 \phi$ to $20 \phi$.
7472 - \$3, $7423-\$ 6,7485-$ - $\$ 4,7473$ - \$1.50, UA709 - \$1.50 to \$2, LM311-\$3.25.
44 pin PC board edge connectors \$1.40
Transistor heat sinks - 75 .
Resistors - 5d to 10d.
In conclusion, it's probably not worthwhile to make a special trip to either Tokyo or Manila to purchase surplus electronics from these markets. However, if one happens to be visiting either city, there are some selected bargains to be had by shopping around.■

It's probably not worthwhile to make a special trip to either Tokyo or Manila to buy surplus electronics; but if you happen to pass through either town on other business, the surplus markets are definitely worth a side trip.


Photo 2: Bargain hunting in Akihabara, a district of Tokyo.

## The Great 8

## With 8080 plug-in modules and software fron

 the capabilities you've been dreaming of 1
## Sound Memory Judgment

Processor Technology now offers two options in low-power, fast static memory. Our read/write static memory modules - the 4KRA (4096 words) and the double-density 8KRA ( 8192 words)-are more reliable, less expensive, and require less power than any other available modules of comparable capacity.

Both memories use the same low-power static RAM's (our prime supplier is AMD). All RAM's are manufactured to exacting military specification MIL STD-883C. The 2102A-4 RAM's used by our competitors are not low power-their data sheets show they use more than double the power (worst case) of ours. Both of our modules will retain data when powered by two " D " flashlight cells.
(Connectors for battery power and re-charging are built in.)

The 8KRA has an exclusive address offset switch. This allows it to recognize any 8 K address segment at 1 K intervals, such as from 2 K to 10 K .

Each integrated circuit has its own premium grade, low profile IC socket (optional in the 4KRA). IC sockets make assembly, testing and repair many times easier. You'll quickly grow to appreciate them.

4KRA (without sockets or DIP switch) 4096-word Static Memory Module \$139

4KRA (all sockets and switch incl) \$154
8KRA 8192-word Static Memory Module (all sockets included) \$295

# J60 Evolution <br> Processor Technology, you can now afford to add your Altair or IMSAI minicomputer system. 

## Accessible Software!

We firmly believe that good software should be easily obtainable for the hobbyist. You shouldn't have to resort to searching or copying source listings, just because some software can cost you as much as your entire minicomputer system. Processor Technology now offers four well-tested software packages. Their prices are within reach, and more are on the way!

## Software Package\#1

 Resident Assembler/(Edioor/MonitorSW\# 1 is very easy to use, even if you have little or no software background. It gives you power to develop and debug small to medium programs in 8080 assembly language. It allows program entry and editing by line number. Up to six program files may be stored at once in memory. This sort of Assembly Language programming should be used with any high speed, real time, or complex I/O applications. (Recommended total memory: 8192 bytes or more.) Source Listing: \$5.00—Paper Tape: \$9.50.

## Software Package\#2 BASIC

This conversational language is most useful in scientific, educational and recreational applications. Commands and statements are close to normal English. The floating-point calculations in our 5 K BASIC are considerably faster than most other 8080 versions. (Any BASIC language is, however, best used in programs with fairly low speed requirements.) Ours is the only BASIC available that can store several programs in memory at once. Its formatted output provides for easier financial calculations, with multiple statements per line, and assembly language routine calling capability. (Recommended total memory: 8192 bytes or more, from address zero.) Source Listing only : \$5.00.

## FOCAL <br> (EDigital Єquipment Corporation)

This is a math oriented language designed for scientific and educational applications. All programs written in standard DEC FOCAL for the PDP-8 will run on 8080 based machines, using our version of FOCAL. (Recommended total memory: 8192 bytes or more.) Source Listing: $\$ 5.00$ - Paper Tape: \$9.50.

## Math Pack <br> Scientific Notation Calculator Program

Our Math Pack will provide basic math functions in a small memory system. In essence, it's a calculator that will add, subtract, multiply, divide and square root. You are allowed four variables, and you can set the decimal point (fixed or floating) between 0 and 12. Utilize either Algebraic or Reverse Polish logic-or mix them! (Program Requirements: 4096 total bytes, starting at address zero; 1300 bytes reserved for program space.) Source Listing and Paper Tape (combination only): $\$ 9.50$.

Processor Technology manufactures a full line of plug-in modules compatible with the Altair and IMSAI minicomputers. Write us now for our free brochure or the location of your nearest PTCO dealer.


Photo 1: The Direct Etch Method. In this method, a one of a kind printed circuit is made by putting the pattern directly onto the copper. A Sanford's "Sharpie" pen (available in most stationery stores) can be used to draw patterns directly, and tape resist can be used for more uniform runs. If tape resist is used, care should be taken to avoid gaps in the adhesion of the tape to the copper.

## Make Your Own Printed Circuits

## James Hogenson

Box 295
Halsted MN 56548

Photos accompanying this article are by Ed Crabtree, using materials supplied by the author.

The widespread commercial use of printed circuits in electronic equipment began a few decades back when engineers started looking for more efficient wiring techniques to replace laborious hand-wiring methods. One of the first methods tried was to deposit (in other words, to print) a conductive ink pattern on a base of insulating material. The original method, printing, gave its name to all subsequent methods. Today, the term printed circuit refers to any electrical circuit in which individual wire lead connections have been replaced by a two dimensional conductive pattern bonded to an insulating base material.

Contemporary printed circuits consist of etched copper foil wiring patterns bonded to
any of several insulating substrate materials sturdy enough to serve as a mounting base for the actual electrical components which make up the circuit. Although originally developed for mass production applications, printed circuit fabrication techniques have been refined until they can now be used by almost anyone with average mechanical skills.

Choosing your base material, the board, is a matter of price and purpose. The best is the epoxy glass board while phenolic (bakelite) is the cheapest. Phenolic base material is perfectly adequate for many applications, but since small boards are relatively inexpensive, epoxy glass is usually the optimum choice. The base material often comes


Photo 2: Printing the Circuit. Once a negative of the artwork has been created, the next step is to print the circuit. The negative is placed over a sensitized PC board and held firmly in place by a glass cover plate in the printing frame. The glass guarantees smooth and even contact for accurate transfer of the image. The board is then exposed to a photoflood lamp for one to three minutes.
laminated with copper foil on one or both ${ }^{\prime}$ sides.

The toughest part of making your first printed circuit board is getting started. In other words, the process may not be as difficult as you had thought.

A pattern of etch resist is applied by one of several methods to the copper foil. The board is then immersed in a chemical solution (usually a ferric chloride solution) which etches away all exposed copper. Then the board is washed and the etch resist pattern removed. The copper foil that was covered by etch resist remains on the board to provide you with a printed circuit.

## Plan the Layout

The first step toward making your own printed circuit board is planning the layout.

Draw the circuit pattern on paper as it should appear on the printed circuit board. You will use this as a guide for laying out the actual etch resist pattern. Keep in mind that you are looking at your board from the bottom when looking at the foil side. Be careful not to put the pattern on the printed circuit board upside down. (l've made that mistake more than once!)

## Direct Etch

Direct resist is a method often used when a one of a kind board pattern is needed. Dry transfer etch resistant patterns are applied directly to the copper. The dry transfer patterns form integrated circuit pads, transistor pads, edge connectors, round donut pads, etc. Narrow etch resistant tape is applied to complete the circuit path between

A printed circuit is any electrical circuit in which individual wire leads have been replaced by a two dimensional conductive pattern bonded to an insulating base material.

The toughest part about making your own printed circuits is getting started.


Photo 3: Results of Exposure. After being exposed, the photosensitive layer is developed, using an appropriate solution. An etch resist pattern will then remain on the board as in the example at left. (The dark blotches are oxidation on the copper.) The board is then etched with the usual ferric chloride solution. The finished product (hopefully free of imperfections) is a printed circuit board such as the one at right.


Photo 4: The "Cut-N-Peel" Method. A sheet of red mylar film on clear acetate backing is placed over the pattern to be copied. The negative is made by carefully tracing the pattern with a razor or sharp knife, then removing the red film wherever component pads and connections are to be made. (A trade name for the film used in this method is "Rubylith.") The negative is then transferred to sensitized copper and etched.


Photo 5: The Bishop Graphics " ' $B$ ' Neg" Method. In this method, a negative is made directly, using self adhesive black patterns on a mylar backing. The connections between patterns are made by cutting away the black layer with a sharp knife as in the "Cut-N-Peel" method.
component pads. Etch resistant ink pens and resist paint are also available for direct etching, as illustrated in photo 1.

After etching, copper will remain on the board only where dry transfer patterns or resist ink protected the copper foil from the etching solution. It should be noted that the etch resistant tape must be applied firmly, especially at overlaps, to keep the etching solution from getting under the tape and breaking the conductive copper path.

The direct resist method does not require extra steps for developing, as does the photo etch method. If only one printed circuit board is going to be made from a pattern, the direct etch method may be a time saver. If more than one board is to be made from one pattern, the direct etch method will quickly turn the element of time against you, since the pattern must be reconstructed on each board.

## Photo Etching

The photo resist method is the most efficient method for making more than one printed circuit board of a kind. Photo resist etching is probably the most popular method, and is often preferred even for one of a kind printed circuit boards. The difference between photo resist and direct resist is the way the resist pattern is applied.

The copperclad board to be photo etched is first sprayed with a thin coat of a photo sensitive etch resist. This etch resist is sensitive to ultraviolet light. The sensitized board must be handled in a darkened room using a yellow light for illumination.

After the resist is dry, a negative of the printed circuit pattern is placed over the sensitized board in a print frame, as shown in photo 2 . The board is exposed to the light of a photo flood lamp through the negative for one to three minutes. It is then immersed in a resist developer solution for about one minute. Only the etch resist which was exposed to the bright light will remain on the copper foil, as in photo 3. The resist is no longer light sensitive after developing, but should be allowed to dry for a short time. The board may then be etched. The copper which is protected by the remaining etch resist will not be removed. After the board has been etched, the resist is removed and the board may be cut, drilled, and assembled.

## Making Negatives

It is plain to see that exposing a board through a reusable negative is much simpler than reconstructing the pattern by hand each time the pattern is used. The negative
may be obtained by a number of methods.
If a pattern is not too complex, the "Cut-N-Peel" method of photo 4 can be used. The pattern is simply cut into a red film on a clear acetate backing. The red film is peeled off, leaving a negative of the pattern.

If the pattern involves integrated circuits, the "Cut-N-Peel" method becomes rather difficult. The Bishop Graphics " 'B' Neg" TM method would be more suitable. The ready made negative component patterns are laid out on a mylar sheet according to desired component placement. The areas between these self adhesive patterns are blacked out, using solid black acetate film. The only cutting necessary is for connections between component patterns. Photo 5 illustrates this method. The finished product is a negative of the entire printed circuit pattern.

## Photographic Negatives

Perhaps the easiest and certainly the most popular method of obtaining the necessary negative is to first make a positive pattern, then produce a negative by photographic methods.

Positive artwork is made on a sheet of clear mylar film with matte finish on one side. This film is dimensionally stable and similar to plastic drafting film used by draftsmen. Positive artwork patterns are widely available in a large number of sizes and shapes. Photo 6 shows an example of a circuit being laid out with these patterns. Unless the artwork is going to be photographically reduced, use $1: 1$ artwork patterns. The self adhesive positive artwork patterns are laid out on the mylar sheet according to your pencil layout. Narrow black tape is used to form conductive paths between components. Graph paper or a similar grid should be used as a guide for orderly and uniform positioning of patterns. Since components are normally configured for dimensions which are multiples of 0.1 inch a 0.1 inch grid should be used.

A negative reproduction of your positive pattern can be made by a photographer or (preferably) by you. If you enjoy experimenting with photography, you might try experimenting with lithographic and orthographic films.

## Photography Without a Darkroom?

The most popular negative producing method does not require photographic darkroom facilities. The special reversing film used may be handled in subdued light or in a darkened room using a dim yellow light. The positive pattern is placed directly on top


Photo 6: The Traditional Photo Negative Method. In this method, a positive artwork pattern is created, using preprinted self adhesive patterns and artwork tape. A sharp knife is used to cut the tape as it is being applied to the mylar film backing. A photographic process must be used to invert the image and create the negative form (see photo 5).
of the reversing film. The film is exposed through the positive artwork pattern to a photo flood lamp for one to three minutes and developed by rubbing gently with a cotton swab and a little film developing solution, as shown in photo 7 . The opaque or colored emulsion on the film will rub off areas not exposed to light. The result is a clear pattern on a dark background.

A somewhat more involved but rather unique artwork developing system is made by Datak. With the Datak film and developing solutions, any of the following can be made: (1) negative from film positive or original artwork, (2) film positive from negative, (3) negative from negative, (4) film positive from film positive or original artwork, (5) film positive from black image on white paper, (6) film negative from black image on white paper. The last two methods allow you to copy a printed circuit pattern directly from a magazine page.

The Datak film is developed by methods similar to standard photographic procedures, so this method is more complicated and time consuming. Exposure and developing times are somewhat more critical. Datak film may, however, be handled in subdued tungsten light.

## Advantages of Photo Resist Techniques

One of several advantages in using the photo resist method will become apparent when a modification of an existing board is

For a one shot printed circuit, simply draw the pattern on to copper with a resist pen and dump the board into ferric chloride until done.


Photo 7: Creating a Photo Negative for Etching. A negative is reproduced from the positive artwork pattern, using a reversing film. The film is exposed with a bright light, then developed by rubbing gently with a cotton swab and developing solution. The result is a negative version of the artwork with a 1:1 scaling.

For a unique approach to making jumpers on one layer boards, see Don Lancaster's "How to Build a Memory With One Layer Printed Circuits" in the April 1976 BYTE, page 28.
made. (Like when you need to make a board over because you forgot two or three connections. This does happen!) Rather than reconstructing an entire printed circuit pattern, make only the necessary changes or additions on the original artwork, then make a new negative and a new board. Making a new negative using reversing film requires only a few minutes of your time.

## Double Sided Boards?

Sometimes a circuit will be too complex to fit on one side of a circuit board. Since a printed circuit is only two dimensional, conductor paths cannot cross. Jumper wires can be used to provide some crossovers, but if the circuit requires a large number of crossovers, a double sided circuit board might be considered. A double sided PC board is one which has a copper foil pattern on each side. The major consideration in making a double sided PC board is getting the pattern and terminals lined up. Both sides of the board are developed and etched at the same time.

## Drilling

The step following the fabrication of a PC board is drilling out the holes. A small bench type drill press is ideal for this purpose. A standard hand held drill is unsatisfactory as the small drill bits break at low speeds. Commercially, small holes are drilled in boards at speeds as high as 70,000 RPM. A Dremel "Moto-tool" is a suitable compromise for work on printed circuit boards. This tool runs at 30,000 RPM. Such a tool will not only drill out extremely small holes, but cut and shape printed circuit boards, and lend itself to a host of other uses not related to making boards. A multipurpose tool like this is handy, especially for cutting out things like board edge connectors.

If repairs or small changes are needed on a printed circuit board, a piece of bare wire soldered over the foil is the cheapest and quickest modification. A conductive silver paint is available for printed circuit repairs, but the paint is quite expensive. GC Electronics, Techniques Inc, Kepro, and Datak each manufacture printed supplies for the
hobbyist in addition to their commercial products. Such supplies are distributed through a large number of mail order firms and retailers. The appendix lists the various products and who makes them. Cost of materials will vary depending upon a number of factors, but a figure of 20 cents per square inch of printed circuit board will provide a good rule of thumb to estimate the cost per board.

You will notice that Techniques and Kepro do not manufacture photo resist spray. Instead, they sell printed circuit board panels with the photo resist already applied. Presensitized panels (which come wrapped individually in dark paper) will assure you of a uniform and dustfree coating of photo resist. However, if you make a mistake developing the resist pattern, you will waste the extra cost of presensitized panels. It is a good idea to start with a spray resist, then graduate to presensitized panels once you have refined your circuit fabrication techniques. And the keys to refining your techniques are: Read instructions and familiarize yourself with what you're doing, follow the instructions, take your time, be careful, and practice first, using small sample boards. Follow those hints and you may surprise yourself with the fine boards you can turn out. ${ }^{-}$

## APPENDIX: Sources of Supply

## Direct etch materials

Ink resist is made by GC Electronics, Techniques, and Kepro.

Dry transfer resist patterns are made by Techniques, Datak, and Kepro.

An ordinary "Sanford's Sharpie" marking pen available for about 49 cents at any stationery store can be used as a resist pen.

## Photo etch supplies

"Cut-N-Peel" and " 'B' Neg" supplies are distributed by GC electronics. (The " 'B' Neg" materials are manufactured by Bishop Graphics.)

Rubylith material, available at art supply houses, can also be used for cutting and peeling patterns.

Positive artwork patterns and supplies are made by Datak, Kepro, and Techniques. GC Electronics distributes artwork materials made by Bishop Graphics. Bishop Graphics materials are also distributed by independent distributors.

## Photo etch supplies

Photo resist spray and developer are made by GC and Datak. Presensitized panels are distributed by Techniques and Kepro. Reversing film and developer are made by Techniques, Kepro, and Datak.

All of the above mentioned manufacturers make or distribute plain PC board panels (unsensitized) and etching solutions.

Photo flood lamps are available at photo supply houses. (Look for 375 Watt reflector flood lamp or No. 2 (EVB) Photoflood.)


Photo 8: Close up, a successfully etched printed circuit will have even lines with no hairline cracks or other imperfections. This example shows such a result, prior to drilling out the holes for component leads.


Photo 9: When various imperfections enter the picture, the result is not so clean. Here is a high contrast picture of an imperfect result. The resist layer has separated from the copper during the etch process at several points, resulting in holes in the copper and, in several instances, complete breaks in circuit runs.

# A Plot Is Incomplete Without Characters 

Who would want to miss the opportunity of creating customized graphics for special applications?

The design of a plotting data format can be likened to designing a special purpose computer instruction set; this instruction set is emulated by the plotting software in real time.

## Richard J Lerseth

8245 Mediterranean Way Sacramento CA 95826

As computer hobbyists, a number of us will sooner or later play around with graphics using vector CRTs or XY pen plotters; but very few of us will be willing to pay the high price of a number of copyrighted plotting packages available today through computer graphics houses. Besides, most of us will not want to miss the opportunity of creating our own packages.

So, in the process of interfacing your graphic media to your computer, you will normally have built the software needed to control simple vector moves on your media, as well as be able to window your plottings (that is, confine your moves within a specified area).

But, you will find that one of your major efforts will be building the character generation module. As you will soon realize, computer graphics take large chunks of memory space for the graphic routines and plotting tables describing plotting sequences. Particularly, you will find that character generation will take a large portion of that memory space.

In this article, I will describe some of the basic concepts of character generation, and describe techniques of saving memory space through efficient programming and by maximizing the packing of information in the plotiting tables.

I assume at this point that you have within your basic plotting software: (1) the capability of shifting the relative origin within the plotting frame and (2) the capa-
bility of chain plotting. That is, plotting a vector from the ending point of the last vector move to the new position on the plotting field without explicitly defining the beginning point every time you make a vector move. We make full use of these two capabilities in plotting the character strings.

## Plotting Frame

The easiest way of plotting a character is to define a plotting frame or grid upon which a sequence of vector moves are made from grid point to grid point. To minimize the complications involved with signed vectors, it is best to set the origin in the lower left hand corner of the field on which the character is to be plotted. With this convention, the vector moves are positive upward and to the right in the grid. In this way, we can define the ending point of a vector move with positive integer coordinates.

## Limiting Frame and Plotting Resolution

Next, we have to define the resolution in plotting the characters. That is, we have to decide how many grid points we desire within a character frame. This depends on many factors: How fine you want your plotting; how many different characters you are to plot; how you are to pack the moves into memory; what special effects or options you desire. These considerations are all interrelated and must be considered as a whole.

I will propose an optimum character grid field within a limiting frame which will have a resolution as fine or finer than any used today by the graphic houses in their charac-

> The choice of a character grid should reflect the realities of the common machine designs. For most microcomputers (and minicomputers), a character frame optimized for 8 or 16 bit words is desirable.
ter plotting packages. It will minimize the use of storage, and will also have some capability for special options. However, as a user of this software, you can make appropriate changes in your own system to reduce the resolution or to eliminate some of the special options.

Figure 1 shows the ( $8 \times 16$ ) grid 1 propose. The storage origin $(0,0)$ is defined to be in the lower left hand corner of the grid. The character base origin $(0,5)$ is at the lower third point of the left hand side of the limiting frame such that upper case alphabetic characters will be confined to the upper two-thirds of the grid frame. The lower third will be used for the tails of lower case alphabetic characters. The lower row of the grid will not be used for plotting; this row of 8 points will be reserved for flagging special options, which will be explained later.

## Specifying Moves

Since most of us are using or will be using 8 or 16 bit machines, choosing this grid optimizes the packing of information for a vector move into an 8 bit byte of memory. A move to any point in this grid field (figure 1) could be defined with 3 bits for the horizontal $(\mathrm{H})$ position, 4 bits for the vertical ( $V$ ) position, and 1 bit for the $Z$ function or the status ( P ) of the move (pen up or pen down for pen plotters, or intensity modulation in video graphics).

The 8 bits of $H, V$ and $P$ data for a move can be packed in the six different ways, such as HVP, VPH, PHV, VHP, HPV, or PVH. However, when packing such data into

the byte, one must consider which is the fastest way to unpack the values. This greatly depends on the machine used. In most cases, it simply entails masking and shifting. I am going to use (VHP) as my standard. Why? No reason except that it can be implemented on most of the micros in use today without excessive effort. One procedure of unpacking the byte is given in appendix $A$.

To clarify further discussions on vector moves, the coordinates of the moves within the limiting frame will be written as $(\mathrm{V}, \mathrm{H})$. When the $Z$ function is included, the move will be defined as $(\mathrm{V}, \mathrm{H}, \mathrm{P})$ where
$V$ is the vertical portion of the move
$H$ is the horizontal portion of the move
$P$ is the status of the $Z$ function where,
0 is pen down or display tube electron gun on
1 is pen up or display tube electron gun off.

The lower portion of the grid $(\mathrm{V}=0 ; \mathrm{H}=0$ to 7; $\mathrm{P}=0$ or 1 ) will be reserved for special options which will be defined later.

## Optimization of the Storage of Character Moves

Once a user starts playing around with developing the moves for each and every character, he soon realizes that there are a number of instances where a chain of moves is duplicated in the patterns of several characters. One can take advantage of this by building subchains and referencing them where it is appropriate to combine them in a large sequence. For example, the upper case

Figure 1: 8 by 16 Character Frame. Characters are plotted (or drawn on a vector graphics display) with reference to this local coordinate system. A series of 8 bit codes identifies the successive locations of the pen (or electron beam) and whether or not a line is to be drawn while moving to the location. The codes which reference the bottom row of this grid are treated as special operation codes for the plotting software: subchain reference, half shift right option, and floating subchain operation are defined in this article.

When implementing this software for a graphic display mechanism such as a CRT, pay attention to speed of execution. Flicker will result if your computer and software cannot keep up with your eye's timing characteristics.

With the character defined, the next task is to shift, twist, stretch or squeeze the characters as they are drawn.
alphabetic characters, (G, C, O, and Q) can all be combined together in one single chain. Also, the many lower case alphabetic characters have ( $c, 5$, or o) as part of their chain. Taking advantage of such duplications can significantly lower the storage requirements of character plotting tables.

## Special Options

When I defined the character limiting frame previously, I reserved the lowest horizontal line of grid points for special options. There are 8 grid points on this line. This gives 8 special options that can be used. If one considers the $Z$ function, there are 16 options in all. Whenever ( $\mathrm{O}, \mathrm{H}, \mathrm{P}$ ) is encountered in a plotting chain of moves, then a special option is initiated. The special options can use the following bytes in the plotting sequence and, as such, can involve one, two, three, or more bytes.

The first special option we need is a subchain option. I shall define the code as ( $\mathrm{O}, \mathrm{O}, \mathrm{O}$ ). When this code is encountered, the next byte in sequence is the subchain number. As one can see, there can be 256 subchains. You will probably never need all 256 unless you build a large multi-language or multi-font character set.

The second option needed is a $1 / 2$ shift right option. The code I used is $(0,1,0)$. This option increases the resolution of the plotting in the horizontal direction and comes in handy when plotting upper case alphabetic

## APPENDIX A <br> UNPACKING A VECTOR MOVE FROM AN EIGHT BIT BYTE

Using $V, H$ and $P$ to denote bits, the move is VVVVHHHP in packed form. The unpacking procedure is as follows:

1. An arithmetic shift right will make the $Z$ function of the move available in the carry flag. The user can make use of this information through appropriate compares and jumps. Note that masking all but Bit $P$ will also make the $Z$ function available, but the action of shifting also readies the horizontal position of the move.
2. Temporarily store present value of the accumulator in any other register.
3. Mask the accumulator with octal 7. The horizontal position is now available. Send it out to the graphic device or store it for later use in another register.
4. Bring back the stored value of the accumulator from Step 2, shift right three times and mask the result with octal 17 . Now the vertical portion is available.
The 8008 microprocessor assembly code would look like:

$\mathrm{M}, \mathrm{T}, \mathrm{V}, \mathrm{W}$ and a number of other characters to make them symmetric in the particular grid frame I propose. It is a one byte instruction to shift the horizontal portion of the next move byte one half grid space to the right. That is, if the sequence of bytes $(5,0,0) \quad(0,1,0) \quad(5,3,0) \quad(5,4,0)$ was encountered, then the next two moves would begin at $(5,0)$, move to $(5,3-1 / 2)$ and end at $(5,4)$ with pen down (or gun on).

These last two options ! consider to be the minimum you should have in your system if you are to have the resolution required to plot large character sets.

Another option that could be used is the floating subchain option, $(0,2,0)$. (This option is not shown in figure 6.) It takes three bytes of code to complete the sequence of this option. For instance, a period is used extensively for a number of punctuation characters and lower case i's and $j$ 's. The subchain sequence $(1,0,1)(1,1,0)(2,1,0)$ $(2,0,0)(1,0,0)$ plots a period in the lower left hand corner of the grid. Now, by using the floating subchain option, this period can be floated anywhere on the grid. A threebyte sequence ( $0,2,0$ ), (SV,SH,0), (subchain no.) will move the period to any location desired by using positive offset values (SV,SH). This would save at least two bytes of storage for every different position of the period in the grid field, if there are more than two positions to be plotted. But, it takes some extensive programming to include this option; the advantage is large in large character sets, but minimal in small sets. Also, since timing is important in using CRT graphic systems, one must consider whether the extra computing effort is worth the savings in memory. I will leave it up to you to dream up exotic plotting options of your own for the 13 additional options which remain undefined.

## Pointer and Move Sequence Tables

The pointer and move sequence tables now have to be established. A general schematic of the tables is shown in figure 2, along with the relationship of the tables to one another.

The primary pointer table defines the starting point in the character vector move sequence table, and the number of moves for each particular sequence. The pointer table is two bytes per character and shown in figure 3. Five bits of the first byte gives a maximum number of 31 primary steps per character. This is more than enough for any character contemplated, even if it were script or gothic. It is conceivable that a sequence table can be as large as 8 pages or 2 K bytes long. The remaining 3 bits of the first byte could designate the page number.

The second byte would designate the starting point within that page．

This two byte table will fit into one 256 byte page of memory if there are 128 or less characters in your set．So，the full ASCII character set would fit easily in one page． The 7 bit ASCII code，if it resides in the upper portion of the address byte（bits 7－1） with a zero in the LSB of the byte，can address the location table directly．The location table for the subchains will also use the same format．

In figure 4， 1 give my version of the full ASCII 128 character set．Tables 1－3 give the values needed to plot this set．The tables contain octal 2235 （decimal 1181）bytes of data．The tables are set up so that you can easily reduce the size of the tables to a minimum set containing only 63 upper case alphabetic，numeric，and punctuation char－

## Table 1．PRIMARY POINTER VALUES

| - |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 000 | 000 | 587 | 11 | 132－113 |
| 002 | 001 | 598 | 17 | 212－126 |
| 004 | 002 | 615 | 12 | 142－147 |
| 006 | 003 | 627 | 8 | 102－163 |
| 010 | 004 | 635 | 10 | 122－173 |
| 012 | 005 | 645 | 12 | 142－205 |
| 014 | 006 | 657 | 14 | 162－221 |
| 016 | 007 | 671 | 9 | 112－237 |
| 020 | 010 | 680 | 10 | 122－250 |
| 022 | 011 | 690 | 8 | 102－262 |
| 024 | 012 | 698 | 8 | 102－272 |
| 026 | 013 | 706 | 6 | 062－302 |
| 030 | 014 | 712 | 7 | 072－310 |
| 032 | 015 | 719 | 11 | 132－317 |
| 034 | 016 | 730 | 2 | 022－332 |
| 036 | 017 | 732 | 4 | 042－334 |
| 040 | 020 | 736 | 11 | 132－340 |
| 042 | 021 | 747 | 7 | 072－353 |
| 044 | 022 | 754 | 9 | 112－362 |
| 046 | 023 | 768 | 8 | 103－000 |
| 050 | 024 | 776 | 6 | 063－010 |
| 052 | 025 | 782 | 10 | 123－016 |
| 054 | 026 | 792 | 14 | 163－030 |
| 056 | 027 | 806 | 14 | 163－046 |
| 060 | 030 | 820 | 6 | 063－064 |
| 062 | 031 | 826 | 8 | 103－072 |
| 064 | 032 | 834 | 6 | 063－102 |
| 066 | 033 | 840 | 10 | 123－110 |
| 070 | 034 | 850 | 4 | 043－122 |
| 072 | 035 | 854 | 6 | 063－126 |
| 074 | 036 | 860 | 9 | 113－134 |
| 076 | 037 | 869 | 6 | 063－145 |
| 100 | 040 | 875 | 1 | 013－153 |
| 102 | 041 | 245 | 11 | 130－365 |
| 104 | 042 | 262 | 12 | 141－006 |
| 106 | 043 | 274 | 8 | 101－022 |
| 110 | 044 | 278 | 14 | 161－026 |
| 112 | 045 | 292 | 9 | 111－044 |
| 114 | 046 | 301 | 12 | 141－055 |
| 116 | 047 | 256 | 6 | 061－000 |
| 120 | 050 | 323 | 6 | 061－103 |
| 122 | 051 | 329 | 6 | 061－111 |
| 124 | 052 | 313 | 6 | 061－071 |
| 126 | 053 | 317 | 6 | 061－075 |
| 130 | 054 | 243 | 7 | 070－363 |
| 132 | 055 | 317 | 2 | 021－075 |
| 134 | 056 | 245 | 5 | 050－365 |
| 136 | 057 | 335 | 2 | 021－117 |
| 140 | 060 | 147 | 13 | 150－223 |

Table 1 （continued）：

| Table 1 |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | 142 | 061 | 160 | 5 | 050－240 |
|  | 144 | 062 | 165 | 9 | 110－245 |
|  | 146 | 063 | 174 | 11 | 130－256 |
|  | 150 | 064 | 185 | 4 | 040－271 |
|  | 152 | 065 | 189 | 9 | 110－275 |
|  | 154 | 066 | 206 | 11 | 130－316 |
|  | 156 | 067 | 217 | 5 | 050－331 |
|  | 160 | 070 | 198 | 17 | 210－306 |
|  | 162 | 071 | 222 | 11 | 130－336 |
|  | 164 | 072 | 233 | 10 | 120－351 |
|  | 166 | 073 | 238 | 12 | 140－356 |
|  | 170 | 074 | 344 | 3 | 031－130 |
|  | 172 | 075 | 337 | 4 | 041－121 |
|  | 174 | 076 | 341 | 3 | 031－125 |
|  | 176 | 077 | 347 | 12 | 141－133 |
|  | 200 | 100 | 359 | 19 | 231－147 |
|  | 202 | 101 | 0 | 8 | 100－000 |
|  | 204 | 102 | 8 | 12 | 140－010 |
|  | 206 | 103 | 23 | 8 | 100－027 |
|  | 210 | 104 | 34 | 7 | 070－042 |
|  | 212 | 105 | 41 | 6 | 060－051 |
|  | 214 | 106 | 41 | 5 | 050－051 |
|  | 216 | 107 | 20 | 11 | 130－024 |
|  | 220 | 110 | 47 | 6 | 060－057 |
| м | 222 | 111 | 53 | 8 | 100－065 |
| $\stackrel{0}{\square}$ | 224 | 112 | 61 | 6 | 060－075 |
|  | 226 | 113 | 67 | 6 | 060－103 |
|  | 230 | 114 | 73 | 3 | 030－111 |
| 茄荷茄 | 232 | 115 | 76 | 6 | 060－114 |
| $\bigcirc$－${ }^{\text {No }}$ | 234 | 116 | 82 | 4 | 040－122 |
|  | 236 | 117 | 23 | 9 | 110－027 |
| $212-126$ | 240 | 120 | 86 | 7 | 070－126 |
| 142－147 | 242 | 121 | 23 | 11 | 130－027 |
| 102－163 | 244 | 122 | 86 | 9 | 110－126 |
| 122－173 | 250 | 124 | ＋95 | 12 6 | $140-137$ $060-153$ |
| 162－221 | 252 | 125 | 113 | 6 | 060－161 |
| 112－237 | 254 | 126 | 119 | 4 | 040－167 |
| 122－250 | 256 | 127 | 123 | 6 | 060－173 |
| 102－262 | 260 | 130 | 129 | 4 | 040－201 |
| 102－272 | 262 | 131 | 133 | 8 | 100－205 |
| 062－302 | 264 | 132 | 141 | 6 | 060－215 |
| 072－310 | 266 | 133 | 378 | 4 | 041－172 |
| 132－317 | 270 | 134 | 382 | 2 | 021－176 |
| 022－332 | 272 | 135 | 384 | 4 | 041－200 |
| 042－334 | 274 | 136 | 388 | 3 | 031－204 |
| 132－340 | 276 | 137 | 391 | 2 | 021－207 |
| 072－353 | 300 | 140 | 393 | 2 | 021－211 |
| 112－362 | 302 | 141 | 395 | 10 | 121－213 |
| 103－000 | 304 | 142 | 407 | 10 | 121－227 |
| 063－010 | 306 | 143 | 397 | 8 | 101－215 |
| 123－016 | 310 | 144 | 418 | 4 | 041－242 |
| 163－030 | 312 | 145 | 397 | 10 | 121－215 |
| 163－046 | 314 | 146 | 422 | 8 | 101－246 |
| 063－064 | 316 | 147 | 430 | 7 | 071－256 |
| 103－072 | 320 | 150 | 437 | 7 | 071－265 |
| 063－102 | 322 | 151 | 444 | 10 | 121．274 |
| 123－110 | 324 | 152 | 446 | 10 | 121－276 |
| 043－122 | 326 | 153 | 456 | 6 | 061－310 |
| 063－126 | 330 | 154 | 462 | 5 | $051 \cdot 316$ |
| 113－134 | 332 | 155 | 467 | 12 | 141－323 |
| 063－145 | 334 | 156 | 479 | 4 | 041－337 |
| 013－153 | 336 | 157 | 409 | 9 | 111－231 |
| 130－365 | 340 | 160 | 483 | 4 | 041－343 |
| 141－006 | 342 | 161 | 487 | 4 | 041－347 |
| 101.022 | 344 | 162 | 491 | 6 | 061－353 |
| 161－026 | 346 | 163 | 497 | 12 | 141－361 |
| 111.044 | 350 | 164 | 512 | 6 | 062－000 |
| 141－055 | 352 | 165 | 518 | 7 | 072－006 |
| 061－000 | 354 | 166 | 525 | 4 | 042－015 |
| 061－103 | 356 | 167 | 529 | 5 | 052－021 |
| 061－111 | 360 | 170 | 534 | 4 | 042－026 |
| 061－071 | 362 | 171 | 538 | 7 | 072－032 |
| 061－075 | 364 | 172 | 545 | 6 | 062－041 |
| 070－363 | 366 | 173 | 551 | 7 | 072－047 |
| 021－075 | 370 | 174 | 565 | 2 | 022－065 |
| 050－365 | 372 | 175 | 558 | 7 | 072－056 |
| 021－117 | 374 | 176 | 567 | 4 | 042－067 |
| 150－223 | 376 | 177 | 571 | 16 | 202－073 |

acters. Appendix $B$ explains how to reduce the size of the tables to the minimum set. But, I encourage you to go in the opposite direction and build up other subsets to add to this basic set. For example: Greek alphabet and mathematical sets, or centered symbol sets for line graphs.

Table 2. SUBCHAIN POINTER VALUES

|  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 000 | , | 397 | 8 | 101-215 |
| 002 | 2 | 409 | 8 | 101-231 |
| 004 | 3 | 438 | 6 | 061-266 |
| 006 | 4 | 598 | 6 | 062-126 |
| 010 | 5 | 617 | 10 | 122-151 |
| 012 | 6 | 627 | 6 | 062-163 |
| 014 | 7 | 604 | 5 | 052-134 |
| 016 | 8 | 578 | 4 | 042-102 |
| 020 | 9 | 578 | 9 | 112-102 |
| 022 | 10 | 671 | 5 | 052-237 |
| 024 | 11 | 617 | 6 | 062-151 |
| 026 | 12 | 701 | 5 | 052-275 |
| 030 | 13 | 598 | 11 | 132-126 |
| 032 | 14 | 571 | 7 | 072-073 |
| 034 | 15 | 571 | 11 | 132-073 |
| 036 | 16 | 587 | 4 | 042-113 |
| 040 | 17 | 665 | 6 | 062-231 |
| 042 | 18 | 629 | 4 | 042-165 |
| 044 | 19 | 784 | 6 | 063-020 |
| 046 | 20 | 802 | 4 | 043-042 |
| 050 | 21 | 591 | 4 | 042-117 |
| 052 | 22 | 810 | 10 | 123-052 |
| 054 | 23 | 810 | 4 | 043-052 |
| 056 | 24 | 627 | 5 | 052-163 |
| 060 | 25 | 842 | 6 | 063-112 |

## Position, Orientation, and Scale

Now that we have the ability to pull out the coordinates for a sequence of moves, we have just begun the job of plotting a character chain. We must translate each character into its appropriate position on the plotting media, then scale it up or down, rotate it into the proper position, and if desired, slant the character. What usually is done is to build conversion coefficients prior to plotting the desired character string. While going through the process of plotting, these coefficients transform the move coordinates residing in the move sequence table to the appropriate coordinates on the plotting media.

This requires that you have the capability of multiplying and dividing floating point numbers in your system. I assume you will either have a calculator chip interface or a floating point software package to draw on. Additionally, you will need the capability of obtaining sines and cosines if you want the ability to rotate the character string out of a horizontal position or to define the slant of a character with an angle.

Before we get into the procedure of shifting, twisting, stretching or squeezing the characters onto the plotting media, we must define a few parameters which are required prior to plotting the character string. In

Table 3. MOVE SEQUENCE VALUES

## An Aside:

The techniques used in this article can be directly applied to any repeatable set of plotting sequences for display on a vector graphics device. For example, the chess pieces and chess board of a chess game display are one possible data display; similarly, a Space War game's space ship symbol output to a graphic display device could use techniques of vector generation and rotation.

| 000 | 040 | 100 | 140 | 200 | 240 | 300 | 340 | 000 | 040 | 100 | 140 | 200 | 240 | 300 | 340 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 121 | 167 | 132 | 124 | 376 | 325 | 176 | 224 | 307 | 234 | 347 | 226 | 365 | 242 | 346 | 262 |
| 320 | 076 | 176 | 132 | 121 | 370 | 256 | 260 | 350 | 174 | 002 | 206 | 370 | 142 | 326 | 000 |
| 364 | 121 | 376 | 176 | 376 | 130 | 312 | 320 | 370 | 152 | 146 | 147 | 130 | 355 | 330 | 003 |
| 372 | 360 | 121 | 216 | 361 | 125 | 300 | 364 | 366 | 142 | 371 | 144 | 124 | 134 | 271 | 263 |
| 336 | 372 | 360 | 252 | 136 | 134 | 360 | 372 | 346 | 365 | 346 | 124 | 241 | 000 | 266 | 022 |
| 136 | 336 | 377 | 244 | 361 | 321 | 374 | 336 | 350 | 320 | 264 | 126 | 304 | 001 | 126 | 000 |
| 221 | 176 | 200 | 300 | 002 | 364 | 273 | 276 | 305 | 360 | 224 | 146 | 236 | 353 | 026 | 002 |
| 236 | 132 | 245 | 320 | 226 | 372 | 314 | 172 | 346 | 376 | 146 | 273 | 101 | 372 | 062 | 000 |
| 121 | 120 | 136 | 364 | 376 | 336 | 354 | 124 | 366 | 120 | 130 | 310 | 116 | 364 | 363 | 001 |
| 360 | 241 | 361 | 372 | 002 | 276 | 372 | 311 | 364 | 137 | 367 | 306 | 323 | 124 | 122 | 275 |
| 372 | 250 | 120 | 336 | 227 | 254 | 364 | 330 | 344 | 132 | 350 | 264 | 214 | 127 | 143 | 034 |
| 354 | 377 | 136 | 361 | 002 | 202 | 342 | 326 | 346 | 176 | 272 | 224 | 275 | 122 | 274 | 263 |
| 314 | 360 | 121 | 376 | 126 | 120 | 302 | 306 | 311 | 136 | 232 | 206 | 134 | 243 | 207 | 264 |
| 272 | 120 | 360 | 002 | 213 | 136 | 264 | 310 | 352 | 255 | 150 | 210 | 155 | 250 | 134 | 124 |
| 260 | 136 | 002 | 367 | 304 | 363 | 221 | 171 | 372 | 130 | 126 | 232 | 132 | 000 | 371 | 205 |
| 273 | 121 | 126 | 002 | 361 | 376 | 26.4 | 210 | 370 | 124 | 125 | 313 | 124 | 001 | 366 | 272 |
| 236 | 360 | 376 | 126 | 376 | 314 | 272 | 206 | 350 | 160 | 372 | 212 | 142 | 275 | 126 | 234 |
| 176 | 241 | 136 | 361 | 120 | 272 | 236 | 166 | 352 | 220 | 301 | 174 | 242 | 054 | 125 | 143 |
| 132 | 256 | 121 | 160 | 136 | 266 | 176 | 170 | 221 | 326 | 316 | 216 | 264 | 032 | 130 | 124 |
| 120 | 377 | 360 | 124 | 125 | 273 | 132 | 067 | 236 | 346 | 201 | 336 | 272 | 024 | 121 | 132 |
| 251 | 136 | 136 | 132 | 142 | 236 | 124 | 130 | 277 | 364 | 216 | 372 | 254 | 042 | 260 | 154 |
| 256 | 365 | 376 | 176 | 200 | 176 | 160 | 131 | 260 | 362 | 161 | 364 | 214 | 363 | 241 | 174 |
| 136 | 372 | 121 | 376 | 300 | 132 | 220 | 150 | 367 | 340 | 256 | 320 | 202 | 122 | 262 | 212 |
| 337 | 002 | 360 | 361 | 342 | 124 | 324 | 146 | 124 | 320 | 320 | 160 | 363 | 243 | 264 | 204 |
| 372 | 367 | 372 | 002 | 364 | 160 | 372 | 126 | 131 | 136 | 337 | 124 | 122 | 264 | 246 | 222 |
| 364 | 002 | 336 | 126 | 372 | 133 | 321 | 130 | 372 | 143 | 240 | 132 | 143 | 272 | 126 | 242 |
| 320 | 126 | 276 | 376 | 354 | 372 | 360 | 002 | 355 | 354 | 176 | 373 | 124 | 254 | 247 | 264 |
| 160 | 125 | 232 | 361 | 316 | 160 | 376 | 167 | 344 | 343 | 341 | 366 | 132 | 134 | 270 | 272 |
| 124 | 132 | 220 | 122 | 216 | 176 | 210 | 370 | 322 | 154 | 364 | 126 | 154 | 125 | 272 | 254 |
| 132 | 221 | 231 | 002 | 154 | 161 | 130 | 366 | 262 | 241 | 372 | 132 | 254 | 130 | 254 | 000 |
| 176 | 160 | 136 | 246 | 132 | 124 | 277 | 002 | 244 | 256 | 336 | 365 | 272 | 331 | 134 | 000 |
| 336 | 124 | 161 | 134 | 124 | 132 | 232 | 166 | 252 | 002 | 254 | 132 | 264 | 350 | 123 | 000 |
| 037 | 077 | 137 | 177 | 237 | 277 | 337 | 377 | 037 | 077 | 137 | 177 | 237 | 277 | 337 | 377 |
| UPPER CASE ALPHA |  |  |  | NUMERALS |  |  | PUNCTUATION |  |  |  | LOWER CASE ALPHA |  |  |  |  |

figure 5 , we see that we need the standard height (S) and width ( $W$ ) of each character, the gap ( G ) between each character, the starting coordinate position ( $\mathrm{X}_{0}, \mathrm{Y}_{0}$ ) of the character string defined as the baseline origin (identical to the relative origin), and the angles ( $\theta \& \beta$ ) defining the orientation and slant of the character string. These parameters must be made available prior to plotting the character string.

Now, let's list the formulae you will use in your plotting routine.

1. Scale Equations

SS $=\mathrm{S} / 10.0$ vertical scale
(1)

SW $=$ W/7.0 horizontal scale
(2)

SG $=$ G/SW width-gap ratio
(3)
2. Rotation Equations
a. Horizontal (H) portion of move

$$
\begin{align*}
& H X=\operatorname{Cos} \theta  \tag{4}\\
& H Y=\operatorname{Sin} \theta \tag{5}
\end{align*}
$$

b. Vertical $(\mathrm{V})$ portion of move

$$
\begin{align*}
& V X=-\operatorname{Sin} \theta=-H Y  \tag{6}\\
& V Y=\operatorname{Cos} \theta=H X \tag{7}
\end{align*}
$$

c. Vertical (V) portion of move corrected for the slant $V X=-H Y+H X * \operatorname{Sin} \beta$

$$
\begin{equation*}
V Y=H X+H Y * \operatorname{Sin} \beta \tag{8}
\end{equation*}
$$

3. Final Coefficients for Rotation and Scale
$D H X=H X * S W$
$D H Y=H Y * S W$

| 000 | 040 | 100 | 140 | 200 | 240 | 300 | 340 | 000 | 040 | 100 | 140 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 263 | 274 | 264 | 264 | 176 | 346 | 225 | 000 | 000 | 002 | 272 | 320 |
| 272 | 205 | 260 | 171 | 002 | 324 | 230 | 016 | 017 | 166 | 172 | 325 |
| 331 | 212 | 273 | 070 | 172 | 306 | 361 | 265 | 171 | 071 | 000 | 266 |
| 346 | 263 | 264 | 131 | 002 | 266 | 002 | 164 | 176 | 170 | 004 | 000 |
| 126 | 274 | 164 | 136 | 072 | 000 | 262 | 172 | 076 | 076 | 000 | 031 |
| 132 | 122 | 172 | 177 | 000 | 006 | 366 | 177 | 070 | 176 | 025 | 361 |
| 263 | 134 | 225 | 076 | 006 | 000 | 000 | 170 | 133 | 000 | 000 | 260 |
| 142 | 373 | 230 | 000 | 165 | 011 | 013 | 070 | 136 | 006 | 026 | 266 |
| 124 | 366 | 171 | 004 | 264 | 000 | 367 | 076 | 000 | 000 | 000 | 366 |
| 132 | 270 | 070 | 265 | 172 | 012 | 360 | 131 | 017 | 013 | 006 | 000 |
| 154 | 246 | 076 | 272 | 272 | 000 | 260 | 134 | 117 | 177 | 273 | 031 |
| 275 | 230 | 261 | 002 | 115 | 006 | 321 | 000 | 110 | 170 | 264 | 137 |
| 134 | 126 | 360 | 267 | 076 | 273 | 324 | 017 | 174 | 070 | 224 |  |
| 263 | 132 | 266 | 002 | 070 | 264 | 000 | 153 | 074 | 076 | 232 |  |
| 002 | 365 | 366 | 166 | 170 | 224 | 014 | 174 | 000 | 116 | 172 |  |
| 126 | 370 | 265 | 171 | 176 | 232 | 367 | 074 | 020 | 134 | 164 |  |
| 274 | 266 | 164 | 076 | 076 | 172 | 360 | 073 | 165 | 130 | 000 |  |
| 261 | 250 | 172 | 071 | 261 | 164 | 260 | 076 | 264 | 135 | 027 |  |
| 122 | 226 | 272 | 176 | 360 | 261 | 266 | 000 | 272 | 156 | 000 |  |
| 206 | 130 | 171 | 321 | 366 | 360 | 165 | 017 | 172 | 176 | 030 |  |
| 132 | 124 | 070 | 324 | 266 | 321 | 264 | 151 | 225 | 000 | 000 |  |
| 274 | 371 | 076 | 367 | 321 | 326 | 272 | 172 | 232 | 022 | 031 |  |
| 123 | 130 | 367 | 360 | 326 | 367 | 232 | 176 | 000 | 000 | 000 |  |
| 274 | 221 | 360 | 260 | 000 | 266 | 224 | 136 | 021 | 023 | 022 |  |
| 263 | 264 | 320 | 266 | 010 | 000 | 231 | 110 | 000 | 000 | 326 |  |
| 134 | 232 | 326 | 000 | 171 | 013 | 172 | 070 | 004 | 024 | 324 |  |
| 263 | 276 | 266 | 005 | 070 | 361 | 000 | 076 | 265 | 000 | 000 |  |
| 124 | 261 | 260 | 000 | 111 | 260 | 015 | 000 | 224 | 006 | 031 |  |
| 132 | 360 | 265 | 006 | 176 | 266 | 000 | 000 | 232 | 165 | 261 |  |
| 025 | 364 | 164 | 000 | 133 | 273 | 004 | 000 | 272 | 264 | 360 |  |
| 030 | 346 | 172 | 007 | 076 | 264 | 271 | 000 | 002 | 002 | 366 |  |
| 132 | 306 | 272 | 171 | 367 | 164 | 170 | 000 | 227 | 226 | 326 |  |
| 037 | 077 | 137 | 177 | 237 | 277 | 337 | 377 | 037 | 077 | 137 | 177 |
| LOWER CASE ALPHA | ADDITIONAL PUNCTUATION |  |  | ASCII CONTROL CHARACTERS |  |  |  |  |  |  |  |



Figure 2: Relationship of the Character Generation Tables. The selected character code is rotated left by one bit to define a number from 0 to 254. This number accesses a 16 bit quantity in the primary pointer table. The primary pointer table in turn locates the beginning of a series of pen locations in the move sequence table which define the character's plot representation. Within that series, there might be a pointer to the subchain table, which in turn points to an often used fragment of the graphics representation located at a different place in the move sequence table. Note that to minimize retrieval effort on machines such as the 8008 and 8080 , sequences of moves should be restricted to single pages of memory.


Figure 3: Primary and Subchain Pointer Formats. The pointer tables are composed of two byte elements which contain information on the number of moves required, and the address of the first move of the sequence.

$$
\begin{align*}
& D V X=V X * S S  \tag{12}\\
& D V Y=V Y * S S \tag{13}
\end{align*}
$$

4. Shift Coefficients Between Character Baseline Origins

$$
\begin{align*}
& \text { DSX }=\mathrm{DHX}(7.0+\mathrm{SG})  \tag{14}\\
& \mathrm{DSY}=\mathrm{DHY}(7.0+\mathrm{SG}) \tag{15}
\end{align*}
$$ highest mountain rather than be content with a mole hill.

5. Final Transformation Equations to be Applied to Each Move

$$
\begin{align*}
& X=X O+H * D V X+(V-5.0) * D H X \\
& Y=Y O+H * D V Y+(V-5.0) * D H Y \tag{16}
\end{align*}
$$



Figure 4: an ASCI/ Graphic Character Set. The plotting tables 1-3 are used to define this set of characters when displayed or drawn on an XY plotter. This flgure was prepared by the author, using a commercial plotter as the output device.
6. Shift Relative Origin from Character

$$
\begin{align*}
& X O=X O+D S X  \tag{18}\\
& Y O=Y O+D S Y \tag{19}
\end{align*}
$$

Formulas (1) through (15) are calculated prior to plotting the first character of a line. The coefficients thus derived will not change throughout the plotting of the character chain. Note that the equations simplify considerably when the angles $\theta$ and $\beta$ are limited to special cases. Two common special cases are $\theta=0^{\circ}, \beta=0^{\circ}$ and $\theta=90^{\circ}, \beta=0^{\circ}$. Substituting the special values of sine and cosine for these angles produces the special cases. These values are:

$$
\begin{array}{cc}
\operatorname{SIN}(0)=0.0 & \cos (0)=1.0 \\
\operatorname{SIN}(90)=1.0 & \operatorname{COS}(90)=0.0
\end{array}
$$

Equations (16) and (17) are the transformation equations used during the plotting where only the values $(\mathrm{H})$ and $(\mathrm{V})$ change for each move. XO and YO are updated as we move to the next character in line to be plotted by using equations (18) and (19).

## Plotting Routine

The plotting routine is outlined in figures 6 and 7 as a flow chart. If you have BASIC, you should not have any problems imple-

## APPENDIX <br> AbRIDGING THE ASCII PLOTtING TABLES

## NOTES:

1. To abridge the plotting tables, do the following:
A. For upper case alphabetic, numerals, and punctuation, only use:
Primary Pointer Table - bytes (octal) 100 to 301
Subchain Pointer Table - none
Move Sequence Table - bytes (octal) 0 to 613
B. For all characters except ASCII control characters, use:
Primary Pointer Table - bytes (octal) 100 to 375 .
Subchain Pointer Table - bytes (octal) 0 to 5
Move Sequence Table - bytes (octal) 0 to 1070
2. If you want abridged Set $A$ above, note that you do not need to include the traps for special subchain option in your program.
3. Note that the move sequence table is set up so that no sequence of moves crosses the boundary of a 256 byte page of memory. This eases the programming of micros such as the Intel 8008 or 8080.
4. Note that the blank or space character was included at the end of the move sequence table. If you abridge the table, move the code to the end of your abridged table and correct the location code in the primary pointer table. Better yet, include in your program a trap to catch any spacing, as there is no actual plotting for this character. Just shift the relative origin to the next character to be plotted.
menting this routine, as BASIC has the required floating point arithmetic and the transcendental functions, sine and cosine. If you plan to implement the routine in machine language, then I dare say you will have a little more work cut out for you. But, the advantage of going this route is that you will take full advantage of your microcomputer's design in order to minimize the use of memory and increase the speed of plotting. Speed is very important if you have a CRT graphics terminal, because of the refreshing problem.

## Summary

In summary, I think you have here a start in creating your own vector character generation package on whatever graphic media you have or plan to use. You can implement the package as I have given it to you or abbreviate, expand, or abridge the package to suit your needs.

I encourage you, though, to expand the


Figure 5: Character Orientation. To add an element of finesse to the plotting function, provision for general purpose rotation and slanting is a desirable feature. There are two angles to specify: angle $\theta$ is the orientation angle of the baseline for a character string; angle $\beta$ is the frame slant relative to a perpendicular through the base line.

basic character set 1 have given you to include foreign language alphabets, a music symbol set, a mathematical symbol set, or a centered symbol set for line graphs. The horizon in character plotting is only limited to your own efforts or imagination. Climb the highest mountain, rather than be content with a mole hill..


## GLOSSARY

Absolute origin: In a typical plotter or display device, there is an absolute origin for all possible positions of the writing mechanism. A common location of this origin is the lower left hand corner of the plotting field, so that points to the right and above can be specified by positive integer displacements.

Byte: A cell in memory which can store 8 bits of information.

Chain: A set of vector moves to be performed sequentially.

Chain plotting: The technique of specifying a movement of the plotting or display mechanism by a series of small movements.

Character frame: A small region of the plotting medium in which motions will take place while plotting a single character. See figure 1.

Coordinates: A point in a two dimensional space can be specified by a pair of numbers. These numbers are the coordinates of the point with respect to a reference point called the origin.

Masking: The technique of selecting bits for inspection using the AND operation and a mask. The word which is being tested is combined with the mask using the AND operation. Every logical 1 bit in the mask will select a corresponding bit in the word being tested; every logical zero bit in the mask forces a zero in the result independent of the word being tested.

Medium: A plot or a display is usually performed on a two dimensional object which can be viewed by a human being. In the context of this article, the medium is the piece of paper or display tube on which you see the resulting characters.

Page: In many microcomputers it is convenient to divide memory into blocks of multiple bytes, called pages. In the context of this article, the Intel 8080 and 8008 definition is intended: a block of 256 bytes whose high order address byte is identical.

Plotting frame: The range of possible positions for the plotting or display mechanism. In most equipment, this is a grid of points specified by two integer coordinates for horizontal and vertical position.

Relative origin: A local origin which is used for convenience of programming. The relative origin is specified by a coordinate pair with respect to an absolute origin of the mechanism used; movements involved in plotting a character are specified with respect to the relative origin to simplify placement of character patterns.

Resolution: A degree of detail involved in the plot. Ultimately this is limited by the resolution of hardware, which is specified as the number of points per linear inch (or centimeter) of display in each coordinate direction.

Subchain: In a chained plotting table, a subchain is like a subroutine of a computer program. It is a fragment of a plot which is often referenced, so use of the subchain economizes the memory requirements of the data tables.

Vector move: In the context of this article, a vector is a line segment which connects two points in the plotting frame. A vector move is the act of moving the plotting mechanism (pen or electron beam) from one of the points to the second point. In a chained approach, as used in inis article, the starting point is implied by the last position of the mechanism and the ending point of the move is specified by the coordinates of the position.

## Get a

## mouthful to

talk about,

## Subscribe to

## :iII

## IMSAI 8080

Call or write for details \& catalog

SYNCHRO-SOUND ENTERPRISES
193-25 JAMAICA AVE.
HOLLIS, N. Y. 11432
(212) $468-7067$ ( 9 AM - 5 PM)

Computer Mart of
New York Inc.
Authorized dealer for: Sphere - IMSAI Processor Technology National Semiconductor Scamp Featuring a complete line of technical publications and books and a full line of chips and parts at MAIL ORDER PRICES. Repair service for all 8080 and 6800 microprocessors

Friendly Service Advice Problem Solving
Open Monday through Saturday 10-6, Thursday until 9:30

314 Fifth Avenue
New York NY 10001
212 279-1048

## SURPLUS $\mu$ P \& TERMINAL

- Not a KIT! Ready to use
- 2 ea; 400 line ( 80 char/line) Cassette Tapedecks
- Kbd, $\mu \mathrm{P}, 9^{\prime \prime}$ da Video Display, \& Power Supply.
- Modular 'Card' Construction
- MOS-LSI, ASCII, RS-232C
- NEW or like-New, GUARANTED
- Size: $20^{\prime \prime} \mathrm{W} \times 24^{\prime \prime} \mathrm{L} \times 33^{\prime \prime} \mathrm{D}$
- Everything INCLUDED

This Data-Management-Station (DMS) is a business-oriented, highly versatile, machine especially suited for mailing, addressing, filing, labelling, data storage, and other time \& personnel consuming tasks.
DIS SPECIAL PRICE: $\$ 699.00$ CF R

We also offer many ELEC. IRONIC MUSIC SYNTHESIS kits, plans, parts, and more... mail us a S.A.S.E. or POSTAGE for FREE INFORMATION

## Good Grief!



Dave Brockman, 11648 Military Rd S, Seattle WA 98168, has a fairly sophisticated home brew system which includes a vector graphics output capability implemented with a surplus CRT display, heavily modified and interfaced to his PDP-8/S. One result is this vector graphics representation of a famous personality which was taken using a 35 mm camera with multiple exposures to compensate for the slowness of the film used. A program was written to paint a oneshot picture on the screen in response to an input keystroke. The program and camera were set up, the computer room was darkened with blackout curtains, the camera shutter was opened, the oneshot picture-drawing program was cycled several times to burn the image onto the film, then the shutter was closed.
should accept the character from the keyboard, then force lower case letters to be upper case letters. In order to avoid interference with the special meanings of the values octal 200 to octal 377 , the high order bit should be forced into a zero state by an AND operation with octal 177. After this processing, the input value should be left in the accumulator prior to return.

The character display routine, CHRPR, accepts a seven bit ASCII code which is passed from the calling routine in the accumulator. The characteristics of CHRPR will vary from device to device and it is the responsibility of the person using this program to provide an appropriate CHRPR. The original Educator-8080 program was designed to be used with a Digital Group TV-Monitor output device which displays 16 lines of 32 characters. This device automatically begins a new line whenever the previous line has been filled, so the user of a Teletype or a video display with line length greater than 32 should generate the equivalent of a carriage return and line feed at the end of each 32 character line. This can be done by setting a software counter to 32 at the start of a line, and decrementing every time a displayable character is reached. When the counter reaches zero, the line feed and carriage return codes are sent and the counter is reset to the starting value.

There is one special character value, octal 177, which must be detected by CHRPR and interpreted if necessary. This is the character used for clearing the screen of a video display, or doing a form feed operation on a printing device. For a Teletype device, the

Table 5: Work Areas. The Educator-8080 program employs several work areas in programmable memory. These work areas are listed in this table, along with initial values where applicable, and commentary. Not shown in this table is one very important work area which is part of listing 1: the location XQTOP (address <2>1046) identifies a two byte area which must be programmable to allow an 8080 operation code to be "dummied in" and executed based upon the user's input command.

| Address | Name | Length | Octal <br> Initial Value |
| :--- | :--- | :---: | :---: |
| $<3>/ 346$ | PSWA | 2 | 000000 |
| $<3>/ 350$ | BANDC | 2 | 000000 |
| $<3>/ 350$ | CREG | 1 | - |
| $<3>/ 351$ | BREG | 1 | - |
| $<3>/ 352$ | CMDAR | 22 | - |

## Commentary

Initial PSW and A values.
Reference to B and C symbolically
Equivalent to BANDC, C value,
Equivalent to BANDC+1, B value.
Command work area, uninitialized
form feed operation can be simulated by spacing out the listing. Simply emit a carriage return followed by several line feeds. For a video display communications terminal, the ASCII form feed character, octal 014, is often interpreted as a command to clear the screen. For the Digital Group TV-Monitor display the screen is cleared then the screen address is set for the upper left hand corner.

The Educator-8080 is very useful to the novice. On the other hand, its implementation requires a skill level beyond that of the beginner, who might fare well to obtain the help of an advanced hobbyist in implementing the program.

For The Digital Group 8080 system, a version of this program is available on a cassette including documentation - contact The Digital Group Software Systems Inc, POB 1086, Arvada CO 80001.

The Data Domain is a computer store serving the Central U.S. Many of you have already contributed to our success, and we wish to thank you. If you have not yet entered The Data Domain, we'd like to hear from you. We are here to serve your needs with a full line of products and services, and to offer assistance in solving both hardware and software problems.

The retail computer store is more than a business to us. The Data Domain is a personal commitment of over 18 years of computer experience toward the growth of an industry which we are convinced will be of major significance to our nation's future in science, education and business.


Offering IMSAI Processors and all major brands of accessories

## Itty Bitty Computers Tiny BASIC for the 6800

Tom Pittman is a professional software person who has developed a debugged and running version of Tiny BASIC for the 6800. Paper tape and documentation are available for $\$ 5$. The object of Mr Pittman's enterprise, Itty Bitty Computers, is to perform "an experiment in favor of the hobbyist." In a letter to BYTE, Mr Pittman points out that his regular customers are commercial custom software purchasers, who pay the usual high costs of custom software; at $\$ 5$ for a high level language, users will get quite a bargain since you will receive a fairly thick paper tape and a 24 page computer printed instruction manual. To order Tiny BASIC for the 6800, send $\$ 5$ to Tom Pittman, PO Box 23189, San Jose CA $95153 . ■$


Photo 1: The Solid State Sales Video Camera Kit, as assembled. The kit includes the two printed circuit boards and all electronic parts, but does not include the packaging into a neat box with lenses, etc. The charge coupled device sensor is shown mounted between the two boards in this photograph.


Photo 2: The video signal out of the Solid State Sales Video Camera Kit is shown here on an oscilloscope. This sort of camera could be used as an inexpensive starting point for software experimentation in pattern recognition using personal systems.

## Will You Look at That?

Solid State Sales, PO Box 74B, Somerville MA, has introduced a video camera kit which uses a $100 \times 100$ resolution charge coupled device to produce an EIA video signal which can be used to drive a standard video monitor. The product comes in the form of a kit of parts including the imaging array, printed circuitry and electronic parts. Since the solid state imaging array does not require any high voltages, the camera is an ideal product for experimentation with video and as a source of video signals to be converted to digital form in the context of pattern recognition experiments. The spectral range of the charge coupled device sensor array starts in the infrared region and includes the visible wavelengths of light, so in security applications it could be used with an invisible infrared light source. The price of the kit is $\$ 225$ including the circuit boards and sensor array as shown in photo 1. For pattern recognition applications a simple converter circuit is needed to convert the video into a binary (black and white) TTL signal with an adjustable threshold. Also, depending upon the application, the user must supply a case and the optics needed to form an image plane at the surface of the sensor array. An example of a digitized signal displayed with the help of an oscilloscope and some external timing circuitry is illustrated in photo 2.■

## FREE!

New CAMBION XQ catalog. Scores of famous CAMBION electronic components for hobbyist or experimenter in small packs or individually at modest cost. Parts available at retail electronic shops or by mail from factory. Coils - chokes - IC sockets - tools - patch cords, and lots more. Circle Readers Service number now.

## CAMBION

## Cambridge Thermionic Corp.

 445 Concord Avenue, Cambridge, Massachusetts 02138 Phone: 617-491-5400
## WANTED

Microprocessors Kit or Assembled
ALTAIR • SPHERE • MOTOROLA INTEL \& OTHERS

Memories, I/O boards, chassis, power supplies, modems, peripherals, specials, prom programmers, A.DCS, DACS, GPIOS, RS232 faces, and any associated software.

FORWARD INFORMATION AND/OR INQUIRIES TO: AMERICAN USED MICROPROCESSORS, EQUIPMENT \& SUPPLY CORP. BOX 515 ,
PRAIRIEVIEW, ILLINOIS 60069 TELEPHONE (312) 634-0076

For complete monthly catalog of items available and in stock Enclose $\$ 10$. for six issues. Many new unused items. Amazing values for everyone.

A-70

## TIRED OF LOADING YOUR COMPUTER BY HAND?

TURN AN AUDIO CASSETTE RECORDER INTO A FASY reliable program loader and mass data storage unit with the perCom Cassette interface for $\$ 59.95$

- Fuliv assembied and tested
- byte Standard lsee Fet '76 BYTE
- Load a IK program in 40 seconds
- May be upgraded to 218 bytes/sec
- Store 100 K bytes on C 60 cassette
- Standard 8 brt parallel intertace
- Essily connected to any computer
- Little or no solisware sequised
- Provision for custom interfaces
- Stgnal level sensar
- Sidetone crecuir
- Operate 2 tape unis simulianeousiy
- Standard wazed circuit card ( $4 \% \%^{*} \times 6 \%^{\prime \prime}$ )
- Described in BYTE (March April "76)
- Comprehenswe Instruction Manual


TO OADEA Send cherk or money ordem lat 55995 Texas renderns add 54. MasierChaton Bank Americard welcome, inclucte all numbert and bigh
 Full relund of not satslioc
PerCom - 'peripherats for persomat computing'
Custom Design Services . 4021 Windsor
Garland. Texas 75042 - (214) 276-1968

## 6 TV GAMES ONA CHIP

Single MOS/LSI Chip \$39.95 each - supply limited Kit $\$ 69.95$ (includes chip) or send $\$ 1.00$ for kit info and data sheet
21 L 02 ( 500 ns ) 1.95
2102 (500 ns) 1.75 8080A 24.95 680027.95 AY51013A 5.00
Add $\$ 1.00$ to cover handling on orders under $\$ 25.00$. \$10.00 Min. Order. Calif. residents add 6\% tax.

Advanced Microcomputer
Products
P. O. Box 17329

Irvine, CA 92713



Pictured above is the new OP-80A High Speed Paper Tape Reader from OAE. This unit has no moving parts, will read punched tape as fast as you can pull it through ( $0-5,000$ c.p.s.), and costs only $\$ 74.50$ KIT, $\$ 95.00$ ASSEMBLED \& TESTED. It includes a precision optical sensor array, high speed data buffers, and all required handshake logic to interface with any uP parallel I/O port.
To order. send check or money order (include $\$ 2.50$ shipping/handling) to Oliver Audio Engineering, 7330 Laurel Canyon Blud., No. Hollywoiod. CA 91605, or call our 24 hr . M/C-B/A order line: (213) 874-6463.


## PRIZES

Grand prize winner will receive, in addition to the fame of having his entry appear on the cover of the December 1976 issue of BYTE,

- \$100
- A lifetime subscription to BYTE
- A bound volume of BYTE Volume 1.
Two runners up will receive $\$ 100$ and a five year subscription to BYTE.
Five honorable mention winners will receive a tee shirt wardrobe (one from Creative Computing, one from Computer Lib, and one from BYTE) and a one year subscription to BYTE.



## COMPUTER ART CONTEST

BYTE announces a contest for artwork suitable for use on the cover of BYTE Magazine. Entries can either be artwork created by a computer or human produced artwork on a computer theme.

All entries become the property of BYTE Publications but will be returned if accompanied by a self addressed stamped envelope. Though only one grand prize winner will be chosen to appear on the cover of the

December 1976 issue of BYTE, any entries which the BYTE staff like well enough to use as covers will be purchased at our usual rates.
All entries must be received by August 31 1976.

Mail to:
Ms Janice Black
Contest Editor
BYTE Publications Incorporated
70 Main Street
Peterborough NH 03458

## Some Electronics Industry News

For those who are not familiar with the electronic engineering marketplaces, Electronic News is one of the best places to find out about late breaking developments in the hardware end of computing. We refer readers to the May 10th edition of that publication, which had the following interesting tidbits: Page 65: Motorola Semiconductor will shortly (early 1977) be sampling the new M6900 processor. This is described as a 16 bit processor aimed at the computer industry. In short it should be like the central processor of a minicomputer.
Page 66: Mostek and Zilog have made a second sourcing agreement for the new Zilog Z-80 super 8080 processor. [BYTE has scheduled a Microprocessor Update article on the processor for August 1976, to be followed by Dr Robert Suding's detailed central processor hardware design in the September 1976 issue.]
Page 68: Motorola will have a 16 K dynamic programmable random access memory in a 16 pin package, to be sampled later this year. The pinout is compatible with similar devices from TI and Mostek. [It is said that the 16 pin dynamics are pin compatible with the present 16 pin 4 K devices.]

For the details, reference Electronic News. The address of the publication is Fairchild Publications, 7 E 12th St, New York NY. Subscriptions to this weekly paper are $\$ 12$ per annum. ${ }^{\text {- }}$

## SOFTWARE AVAILABILITY NOTE:

 Perspective Plot Package (in ANSI FORTRAN Language)The subroutine package contains 22 calls needed to generate a perspective plot of a three dimensional object or scene described by coordinate points and line segments between such points. The package uses two external calls to the Cal Comp's basic drum plotter system. These calls can be converted very easily to any other system.

This package is being made available, free of restrictions and cost (except copying and mailing charge), to the individual hobbyist, clubs, schools, colleges, and universities for their own computer systems.

Any commercial use of the package in any part, form, or media for sale as Software or given as an inducement to selling other Software or Hardware is strictly forbidden without the expressed permission of the author.

Copying and mailing cost for each copy is $\$ 5$, which includes:

1. Users Manual - 37 pages.
2. Hard copy listing of Subroutines - 23 pages.

> Richard J Lerseth 8245 Mediterranean Way Sacramento CA 95826 (916) 381-0335 (after 5:00 PM)


101 BASIC Computer Games is the most popular book of computer games in the world. Every program in the book has been thoroughly tested and appears with a complete listing, sample run, and descriptive write-up. All you need add is a BASICspeaking computer and you're set to go.

101 BASIC Computer Games. Edited by David H. Ahl. 248 pages. $81 / 2 \times 11$ paperbound. $\$ 7.50$ plus $75 d$ postage and handling ( $\$ 8.25$ total) from Creative Computing, P.O. Box 789-M, Morristown, NJ 07960.





Paul Terrell, founder of the BYTE Shop, the world's first computer store franchise, puts lock and key to "the family jewels." Each BYTE Shop Computer Store has on display, as a distinguishing mark of its identity, BYTE magazines under glass. The BYTEs were placed under glass as a result of an unfortunate rip off of issue number one after a BYTE subscription ad entitled "A Lesson in Economics" mentioned that number ones were selling for $\$ 15$.

Though current issues of BYTE are not yet as rare as the first few, they soon will be, as we do not print a large overrun. Don't miss any. Subscribe today. Use the pull out card in this issue or the coupon below.

Subscribe to BYTE and Creative Computing and save $\$ 2$.
Since we feel that you really need both magazines to be fully informed in the small systems field, the publishers are offering a special package. Receive one year of both BYTE and Creative Computing, normally $\$ 20$, for only $\$ 18$. A ten percent saving.
Please enter my subscription to:
$\square$ BYTE 1 year $\$ 12,2$ years $\$ 22,3$ years $\$ 30$ (circle one)BYTE and Creative Computing
1 year \$18, 3 years $\$ 46$Check enclosedBill meCharge to MC\# BAC\#

Exp. date $\qquad$
Name

Address $\qquad$

City $\qquad$ State $\qquad$ Zip

Please allow six weeks for processing.

## Software Bug of the Month 2

A programmer was trying to learn how to write DO loops without using DO statements, and was getting exceedingly frustrated. Every other time he wrote a loop, there would be either one case too many, or one case too few.

As an example, he was trying to print out a table of sines, cosines, and tangents, just as one would find in trigonometry tables. He was doing it by degrees, from one degree up to 90 , and somehow the program always ended up printing an extra case, for 91 degrees.

Our programmer knew that the sine, cosine, and tangent functions work with radians, not with degrees. He therefore started with the number that is equivalent to one degree in radians, and worked upward from there. Thus his FORTRAN program was as follows:
$P I=3.1415926$
DEGREE $=$ PI/180.0
$X$ = DEGREE
$1 Y=X / D E G R E E$
WRITE $(6,91) Y, \operatorname{SIN}(X), \operatorname{COS}(X), \operatorname{TAN}(X)$
91 FORMAT(F6.2,3F15.8)
$X=X+$ DEGREE
IF (X.LT.PI/2+DEGREE) GO TO 1
STOP
END
The statement $\mathrm{Y}=\mathrm{X} / \mathrm{DEGREE}$ converts X from radians back to degrees, in order to make the table come out with a value in degrees.

Can you find the bug?
Solution in next month's BYTE

## SOLUTION TO BUG OF THE MONTH 1

The problem is not in the statement $\mathrm{J}=(\mathrm{l}$ $+K) / 2$. It is true that $I+K$ might be an odd number, and thus $(I+K) / 2$ might not be exactly half of I + K; but this, as it turns out, doesn't matter.

The problem is in the figure, which was drawn in a misleading manner. At the right is the figure as it should have been drawn.

The element with index $J$ is the last element of the first half of the table. The first element of the second half of the table is therefore the element with index $J+1$, not J . Therefore the new value of I , if $\mathrm{X}>$ $A(J)$, is $J+1$, and not $J$. We should change statement number 2 to read $\mathrm{I}=\mathrm{J}+1$.

What happens if the bug is not fixed can be seen, for example, if $X=A(2)$. Eventually

we will get down to $I=1$ and $K=2$. Now $J$ is set equal to 1 , which is all right in itself; we test $X$ against $A(1)$, and find that it is smaller than $A(1)$. But then $I$ is set equal to J, which is 1 , exactly the same as it was before. This causes the endless loop we mentioned. You can check that it does not occur for any possible value of $I$ for which $X$ $=A(1)$ (which ones does it occur for?). .

[^2]

## You'll Want to Nybble at these Byte Books

Where does the editor of a computer magazine turn to when he must verify some author's hardware design? Information on a 75450 interface gate, or a 74147 priority encoder circuit does not spring forth by magic. Checking the information supplied by authors is part of BYTE's quality control program.
When you build a project, you need this same sort of information. All you find in the advertisements for parts are mysterious numbers identifying the little beasties ... hardly the sort of information which can be used to design a custom logic circuit. You can find out about many of the numbers by using the information found in these books. No laboratory bench is complete without an accompanying library shelf filled with references - and this set of Texas Instruments engineering manuals plus Don Lancaster's TTL Cookbook will provide an excellent starting point or addition to your personal library.

- The TTL Cookbook by Don Lancaster, published by Howard W. Sams, Indianapolis, Indiana. Start your quest for data here with Don's tutorial explanations of what makes a TTL logic design tick. 335 pages, $\$ 8.95$.
- The TTL Data Book for Design Engineers, by Texas Instruments Incorporated. How does an engineer find out about the TTL circuits? He reads the manufacturer's literature. This 640 page beauty covers the detailed specs of most of the 7400 series TTL logic devices. No experimenter working with TTL has a complete library without The TTL Data Book for Design Engineers. Order yours today, only $\$ 3.95$.
- The Supplement to The TTL Data Book for Design Engineers, by Texas Instruments Incorporated. What happens when you can't find a 7400 series device listed in The Data Book for Design Engineers? Before you start screaming and tearing your hair out in frustration, turn to the Supplement. The Supplement has 400 pages of additional information including a comprehensive index to both TTL Data Book volumes. To be complete (and keep your hair in place and vocal cords intact) you'd best order the supplement at $\$ 1.95$ to accompany the main volume.
- The Linear and Interface Circuits Data Book for Design Engineers, by Texas Instruments Incorporated. When you run across one of those weird numbers like 75365 the immediate frustration problem occurs again. What kind of gate could that be? We won't tell in this ad, but you can find out by reading the specifications in The Linear and Interface Circuits Data Book for Design Engineers. You can interface your brain to the 72xxx (linear) and $75 \times x x$ (interface) series of functions by ordering your copy of this 688 page manual at only $\$ 3.95$.
- The Semiconductor Memory Data Book for Design Engineers, by Texas Instruments, Incorporated. Don't forget the importance of memories to your systems. Refer to this 272 page manual to find out about the T.I. versions of many of the popular random access memories and read only memories. Order your personal copy today, only $\$ 2.95$.


#### Abstract

- The Transistor and Diode Data Book for Design Engineers, by Texas Instruments, Incorporated. You'd expect a big fat data book and a wide line of diodes and transistors from a company which has been around from the start of semiconductors. Well, its available in the form of this 1248 page manual from T.I. which describes the characteristics of over 800 types of transistors and over 500 types of silicon diodes. This book covers the T.I. line of low power semiconductors (1 Watt or less). You won't find every type of transistor or diode in existence here, but you'll find most of the numbers used in switching and amplifying circuits. Order your copy today, only $\$ 4.95$. - The Power Semiconductor Handbook for Design Engineers by Texas Instruments, Incorporated. To complement the low power transistor handbook, T.I. supplies this 800 page tome on high power transistors and related switching devices. Here is where you find data on the brute force monsters which are used to control many Watts electronically. Fill out your library with this book, available for only $\$ 3.95$.


- Understanding Solid State Electronics by Texas Instruments, Incorporated. This is an excellent tutorial introduction to the subject of transistor and diode circuitry. The book was created for the reader who wants or needs to understand electronics, but can't devote years to the study. This 242 page softbound book is a must addition to the beginner's library at only $\$ 2.95$.
- The Optoelectronics Data Book for Design Engineers by Texas Instruments, incorporated. This 366 page book is a compendium of information on T.1. phototransistors, LEDs and related devices. Order yours at $\$ 2.95$.

Buyers of these books should be cautioned: heavy reading will be required. These books are so filled with information that they weigh in at a total of about 190 ounces ( 5387 grams). On the basis of sheer mass, these books have got to be the bargain of the century. Make sure that you use a structurally sound book shelf and above all avoid dropping one of these books on your foot.

TTL Cookbook @ \$8.95
TTL Data Book @ \$3.95
Supplement to TTL Data Book @ \$1.95
Linear and Interface Circuits @ \$3.95
Semiconductor Memory Data @ \$2.95
Transistor and Diode Data Book @ \$4.95
Understanding Solid State Electronics @ \$2.95
Optoelectronics Data Book @ \$2.95
Power Semiconductor Handbook @ \$3.95

Piease add $\overline{75}$ cents for postage and handing. 一一 Please allow six weeks for delivery. Send to: Name

Address
City $\quad$ State $\quad$ Zip

- Check enclosed
- Bill MC \#

Exp. Date

- Bill BA \# Exp. Date

On May 2 1976, the Amateur Computer Group of New Jersey held the first Trenton Computer Festival, a combination of manufacturers displays, flea market activities, and technical presentations. The site of the affair was Trenton State College, in Ewing Township NJ. The photos give several vignettes of the activities.

There were a total of 45 exhibitors present, and 25 speakers in the technical program. The attendance was counted by that effective method of selling tickets which registered guests for the door prize drawing. Of the 1500 people who purchased tickets at the door, 180 received door prizes at the end of the day. People from 16 different states attended.

Despite a torrential downpour (which set local records) the day before, the skies cleared for a beautiful sunny festival day. This made the outdoor flea market activities a big success. Items on sale at the flea market moved quite briskly, and ranged in price from 25 cents or less to the $\$ 10,000$ asking price one individual had on his original Samuel FB Morse telegraph key.

Exhibitors booths included firms ranging from computer stores to such firms as

Photos by Marj Kirk


Photo 1: Claude Kagan, a researcher at Western Electric, gives a talk on "Computers in the Home, Present and Future."

Motorola Semiconductor Products, Digital Equipment Corporation, and RCA Solid State Division. ${ }^{-}$


## Still More BYTE's Books



- DESIGNING WITH TTL INTEGRATED CIRCUITS by the Components Group, Texas Instruments Inc. Edited by Robert L Morris and John R Miller.

People often ask questions like "Where do 1 get basic information on hardware design?" One answer is in "Designing With TTL Integrated Circuits."
This book, published by McGraw Hill in 1971, is a fundamental starting point for any person designing peripherals and custom logic employing TTL integrated circuits. While its publication date precludes any reference to the later additions to the TTL 7400 series of components found in the Data Books, it is nevertheless the source of a wealth of ideas on TTL integrated circuits and design of logic with this family of circuits.

What is fanout? You may have heard this term mentioned at computer club meetings or in advertisements for circuitry, or in articles in BYTE. You can find out background information on the calculation of fanouts by reading the chapter on Circuit Analysis and Characteristics of Series 54/74.

Worried about noise, shielding, grounding, decoupling, cross talk and transmission line effect? (Or, more properly, did you know you should worry about these effects in certain circumstances?) Find out about general precautions and background information by reading the chapter on Noise Considerations.

The chapter on Combinatorial Logic

Design gives 53 pages of background information on Boolean algebra and practical representations of logic in the form of SSI gates. The chapter includes a description of Karnaugh mapping techniques and the minimization of logic. From combinatorial design, the book progresses into Flip Flops, including background information on the workings of these devices, and fairly detailed descriptions of the uses and applications of these devices including synchronization of asynchronous signals, shift registers, flip flop one shots, etc. Then the book returns to static combinatorial logic with its description of the Decoders available in the 7400 line as it stood in 1970-1971.
A chapter on Arithmetic Elements gives fundamental descriptions of binary arithmetic, diagrams of the basic gate configurations for combinatorial logic adders, and a section on number representations for use in computers. Much of the material in this section is dated, due to the fact that the later 74181 series of multiple function arithmetic units had not yet appeared when the book was written. But for a background on arithmetic operations implemented with the simpler 7483 circuits, this chapter is ideal. A chapter on Counters and a chapter on Shift Registers complete the detail logic sections. The book is closed out by a chapter on miscellaneous Dther Applications including a simple binary multiplier, a 12 hour digital clock and a modulo-360 adder.
__Designing With TTL Integrated Circuits \$24
__Microcomputer Design $\$ 25$
Send to: Name

| Address |  |  |
| :--- | :--- | :--- |
| City | State | Zip |
| $\square$ Check enclosed |  |  |
| $\square$ Bill MC No. | Exp. Date |  |
| $\square$ Bill BA No. |  |  |
| Signature |  |  |



The most important use of this book is its value as an introduction to TTL logic. By reading and studying it, you will begin to understand the ways in which SSI and MSI TTL gates can be utilized in your own experimental logic designs. After studying this text, you should be able to make much more sense out of the technical information summaries typically published as specifications sheets and data catalogs.

Order your copy today from BYTE's Books, $\$ 24$ postpaid.

- MICROCOMPUTER DESIGN by Donald P Martin, Martin Research. Edited and Published by Kerry $S$ Berland, Martin Research.

Purchase your copy of the definitive source for circuitry and hardware design information on the 8008 and 8080 computers today.

Even Intel, the originator of the microprocessor revolution, is hard put to compete with the wealth of information found in Martin Research's new second edition of Microcomputer Design. This is the book which was originally published as an expensive (but quite practical) engineering report in loose leaf form, at about the time the microprocessor technology was first catching on in the form of the 8008. This 388 page second edition of the manual is loaded with detailed information on how to build and use computers based on the 8008 and 8080.

But even if you do not intend to use the 8008 or 8080 , the practical pointers on digital logic design, peripherals and applications of hardware techniques will more than justify the new low price of $\$ 25$ for this handbook. Microcomputer Design is a must for 8008 owners and 8080 owners who want to truly understand how their processors process.

Microcomputer Design is complete with numerous illustrations, tables and diagrams, plus reprints of the specifications sheets for the Intel processors. There are numerous practical examples of circuitry and many complete computer designs ranging from "minimal microcomputers" to a full blown 8080 processor.

Order your copy today, \$25 postpaid from BYTE's Books.


Photo 2: Herb Nelinger of Miami FL gets a demonstration of one of the Delta Data CRT terminals which were sold at the festival. To his complete surprise, Kevin Moran of Delta Data sold out all his terminals practically before the doors had opened as the ACGNJ people running the show snapped them up.


Photo 3: For some (namely the lucky winners) the drawing at the end of the day's activities was the high point. The young lady holding the cardboard box has won an MOS Technology KIM-1 processor board which was one of the 180 door prizes given out at the festival.

## MIKRA-D 16K STATIC RAM

## ALTAIR/IMSAI Plug-in Compatible

- 16K BYTES of static 500 ns . memory for your ALTAIR/IMSAI Microcomputer.
- YOUR processor runs at full speed.
- PLUGS directly into ALTAIR/IMSAI machine.
- START with 4K - expand to 16K IN ONE SLOT using our expansion kit.
- ALLOWS maximum possible 8080 memory ( 64 K bytes) in 4 slots.
- MIL-SPEC tested memory chips use low-power technology.
- MEMORY protect feature. Buffered inputs.
- SOLVES DMA problem caused by Dynamic memories.


## PRICES

## MD-2046-4 \$175.00 MD-2046-12 \$455.00 <br> MD-2046-8 \$315.00 MD-2046-16 \$595.00 Expansion Kit $\$ 140.00$

## AUTOMATIC 1702A PROM PROGRAMMER

## MD-2044

RS 232 compatible interface. Use with any computer serial output. Programs 1702A in 2 minutes. Complete self-contained unit. Simple software routine allows you to read or program 1702A completely under software control.


Mikra-D•P.O. Box 403•Holliston, Mass. 01746•Tel. 617-881-3111

## ALTAIR 8800 OWNERS

We recently received the following letter:

RUN
APRIL 26.1976
GENTLEMEN:
1 dust wanted to tell you that 1 think your cleck fix-it kit is REALLY GREATI I WAS HAVING trotiale runnimg basic and after INSTALLING YOUR KIT FOUR OF MY MITS BCAADS TAAT WEREN'T REGNVING CAME BACK TO LIFE AND NOW AGE HELPING NE TO WRITE THIS LETTER ON THE COMPUTER. ENCLOSED IS ANOTHER ORDER FCR A CLOCK KIT. THIS IS FOR THE SECOND ALTAIR THAT I'AM NOW IN THE PROCESS of building.
AGAIN HiANY tianks for slch a rime fagoduct.
SINCERELY
SINGERELY
LLCYE L. SAITH
on

How well does your Altair run? A Clock Fix Kit is only $\$ 15$ postpaid.

## PARASITIC ENGINEERING

PO BOX 6314
ALBANY CA 94706

## Clubs and Newsletters



Don Stevens, 545 Meadow Ln, Sheboygan Falls WI 53085, is the contact person for a computer club which has been formed in Wisconsin. An alternate contact for this club is James S White, 1202 Riverview Ln, Watertown WI 53094.

## Kansas City

Harold Schwarz, 7631 Broadmoor Ln, Overland Park KS 66204, would like anyone in the Kansas City MO and Kansas greater metropolitan area to contact him if interested in forming a computer club. Phone (913) 371-2616 from 9 AM to $4: 30$ PM or (913) 648-5410 after 6 PM.

## Language Processor's Committee?

Here is an idea for a club or interest group defined by a logical classification rather than a geographical one: Robert Heller, Box 51A Star Route, Wendell MA 01379, requests "anyone interested in starting a committee to study the possibility of writing language processors for $\mathrm{PL} / \mathrm{M}$ or SNOBOL 4 for 8008,8080 or 6800 systems, please write to me." The idea presumably is to achieve the necessary software by writing processors for the fun of it , and incidentally achieve a time and money sharing effect.

## San Diego Computer Society

The latest Personal Systems (the SDCS newsletter) to pass our desk was Volume II Issue 1 for March 1976. In this issue was an

SCCS Personal Computing ' 76 trip.
The Southern California Computer Society has made arrangements with Travel Coordinators, 8317 W Third St, Los Angeles CA 90048, for a group tour travel arrangement for individuals attending Personal Computing '76, August 28 and 29. The show looks like a big event, and will be held in Atlantic City NJ with representatives of most manufacturers in the field purchasing booth space and contributing door prizes.
The tour basis group rate air fare will be available through Travel Coordinators from any major city in the US. For information write Personal Computing ' 76 Trip, Travel Coordinators, 8317 W Third St, Los Angeles CA 90048, or call (213) 655-0650. To make a reservation by phone call (213) 655-0650 collect.
excellent potpourri of technical and miscellaneous information including: Ralph McElroy's article on how to convert a television into a video monitor as exemplified by a Hitachi portable TV chassis conversion made with the aid of a Sams Photofact File; "Understanding the VDM" by Stan Skoglund, wherein you'll find information on how the Processor Technology VDM module can be programmed and utilized; a reprint of part one of Alan Hastings' "A Software UART" found in The Analytical Engine of the Chesapeake Micro Computer Club; a BASIC game called AMAZIN by Frank Maclachlan, "A Microprocessor Survey" by Lance A Leventhal, "Power Supplies and Regulators" by David G Llenaresas, and a Super Nim program by Jim Farschon. The technical content of Personal Systems is excellent, and will benefit many readers. This newsletter can be obtained for an SDCS membership at $\$ 10$ per annum. Contact SDCS, PO Box 9988, San Diego CA 92109, attn: Newsletter.

## The Analytical Engine, Continued

Each issue of The Analytical Engine put out by the Chesapeake micro Computer Club, 236 St David Ct, X4, Cockeysville MD 21030 (phone (301) 667-9690) is an excellent addition to the literature of personal computing. Here is a club newsletter that is well worth a subscription even if you live in Prudhoe Bay AK. The March issue contained several views about the issue of software piracy and proprietary rights. The major technical contribution of this issue was Carl Hallberg's background article on use of the common garden variety programmable

## HIT THE BEACH <br> 檑 㖇 㖇

# Atlantic City，N．J． August 28th－29th 

\(\left.\left.$$
\begin{array}{r}\text { What its all about！} \\
\text { Software Development } \\
\text { Micro Computers } \\
\text { Hardware Development } \\
\text { Disc Memories } \\
\text { Computer Comparisons } \\
\text { Interfacing } \\
\text { Program Implementation } \\
\text { AMSAT }\end{array}
$$\right\} \begin{array}{r}Computerized Music <br>
Video Terminals <br>
Kit Construction <br>

Printers\end{array}\right\}\)| Computer Games |
| ---: |
| Digital Tapes |

－Seminars and Technical talks by leading electronic equipment manufacturers
－Major Exhibits from all over the country
－Demonstrations in many areas including Home and Personal Computing
－Door Prizes，Free Literature and Free Mementos
－All this plus Sun and Surf－Fun and Excitement－Relaxation and Leisure
Weekend Fair admission $\$ 5.00$ advanced．$\$ 7.50$ at door
Admission includes Exhibits，Seminars
Write for FREE TRIP－KIT to Personal Computing 76 Fair Headquarters Shelburne Hotel－Motel Box 1138 Boardwalk and Michigan Ave．Atlantic City，New Jersey 08404 EXHIBITION BOOTHS STILL AVAILABLE－CALL（609）927－6950


UART integrated circuits, and a continuation of John A Hastings' article on programming a software UART algorithm. Besides the technical "meat" found in the two asynchronous communications articles, the 16 page Analytical Engine includes a flourish of professionalism with halftone photos.

## Washington Amateur Computer Society

The nation's capital computer club sent a February edition of their newsletter, Volume 1 Number 1. It is one of the few such papers prepared and reproduced using a line printer. (The dot matrix (5 by 7) upper and lower case listing suggests it is the product of a DECWriter.) The newsletter included the club constitution (brief and to the point), plus a questionnaire. Contact WACS c/o Robert ! Jones, 4201 Massachusetts Av, Washington DC 20016.

## Are There Any Others in Idaho?

Edwin S Hill, 1900 W Quinn \#117, Pocatello ID 83201, would like to meet others who want to form a club to service southeastern Idaho, southwestern Wyoming, northwestern Colorado and northeastern Utah. We know there are a bunch of BYTE readers in that area, so how about a club?

## Westchester Amateur Computer Society

The kickoff meeting of the Westchester Amateur Computer Society was April 1 1976 at the Community Room of the Greenburgh Public Library in Elmsford NY. For residents of Westchester County NY interested in further information, call Harold Shair at (914) 967-7853, or write him at 41 Colby Av, Rye NY 10580.

## Quest: Philadelphia

Walter White and Everett Holland are interested in making contacts for computer club activities in the Greater Philadelphia Area. Anyone interested, please drop a line to Everett Holland, PO Box 462, Wayne PA 19087, or call evenings at (215) 647-8460.

## Re-Quest: Philadelphia

2005AD Inc, Philadelphia PA, wishes to establish a computer club in the metropolitan area. Write 2005AD Inc, 2005 Naudain St, Philadelphia PA 19146. [See also page 70 of February 1976 BYTE for a listing of the SCCS Delaware Valley chapter in Philadelphia/Camden area.]

## CACHEing IN?

A quite unofficial contact from Charles Douds reported on some of the activities of the Chicago Area Computer Hobbyist Exchange. In his letter, he mentioned that the March 28 meeting included a speaker from Motorola talking about the 6800 family (what else?) and an IBM salesperson giving "a full blown IBM dog and pony show" about their top of the line (for those with limited budget) IBM 5100 table top APL and BASIC machine. [The 5100 is an interesting beast; word on the technological grapevine has it that the APL and BASIC were implemented by having a microcontroller emulate the System/360 (370) architecture. Then IBM engineers stored the machine language object code of an appropriately patched version of the "big system" APL and BASIC in the monstrously large ROM chips IBM uses.] Also reported in Charles' letter is a suggestion attributed to Ruth Lowe that the name of the CACHE Newsletter be changed to "CACHE Register" on the grounds that nobody ever has enough registers to work with in low level code.

Individuals interested in finding out about CACHE should write CACHE, PO Box 36, Vernon Hills IL 60061.

## Unofficial Purdue University Club Philanthropists Take Notice

"A group of (insane) Computer Bums at Purdue University is looking for a PDP-11 series mini, along with associated equipment. We are interested in any and all hardware which might be useful for same. In trade we can offer a limited amount of cash, and/or custom software work. We are very familiar
with: CDC- 6000 series, $C D C-1700$ series, all PDP series, HP-2100 series, Imlac graphics, PDP graphics, 8008, 8080, 6800, and others, all major languages. If you have spare equipment you don't use and need custom software, we welcome your inquiries. Contact: JGM/TGI/GHG at PO Box 2345, West Lafayette IN 47906 or (317) 463-7167. Philanthropists very, very welcome."

## North Alabama Club Activity

Persons in the Huntsville - North Alabama area interested in any aspect of amateur or hobby computing are invited to join the North Alabama Computer Club (NACC). Contact Jack Crenshaw, 1409 Blevins Gap Rd SE, Huntsville AL 35802. Phone (205) 859-7344 or 883-7973.

Mid Michigan Microcomputer Group $\left(\mathrm{M}^{\mathbf{3}} \mathrm{G}\right)$
Activities in the central portion of Michigan have resulted in the "Mid Michigan Microcomputer Group." For further information, contact William Serviss, president, at 13121 Tucker Dr, DeWitt MI 48820, phone (517) 669-3179, or Daniel Herrick, vice president, at 1214 Frederick St, Box 513, Owosso MI 48867, phone (517) 723-3264.

ALTAIR OWNERS! . . WANT TO COMMUNICATE WITH THE ANALOG WORLD??


Available now is an analog to digital converter, ( $A D C$ ) and a digital to analog converter, (DAC); both on the same card which plugs directly into your 8800 and is low power!

| Features | DAC | ADC |
| :---: | :---: | :---: |
| Resolution | 8 Bits, Binary | 8 Bits, Binary |
| Linearity | . $2 \%$ Fullscale | . $2 \%$ Fullscele |
| Temp Coef | 50 ppm | 50 ppm |
| Ranging | $0-10 \mathrm{~V}, \pm 10 \mathrm{~V}, \pm 5 \mathrm{~V}$ | $0.10 \mathrm{~V}, \pm 10 \mathrm{~V}, \pm 5 \mathrm{~V}$ |
| Settling Time | 5 Micro-Sec |  |
| Conversion Time | N/A | 15 micro-sac |
| External Inputs | N/A | Convert start*, Analog |
| External Outputs | Analog out | Serial Out, Strobe for |

- Both $A D C$ and DAC have full software and interrupt capabilities -
* NOTE: Convert start input may be either software or hardware initiated, or both


## KIT: $\mathbf{\$ 2 5 0 . 0 0}$ Assembled and tested: $\mathbf{\$ 3 0 0 . 0 0}$ <br> BankAmericard, Master Charge accepted

AVAILABLE SOON: An eight channel multiplexer and demultiplexer board which bolts directly to the converter board and expands input-output capabilities.
Write or call for more information:
3199 Trinity PI
MACRODATA MICRD San Jose, CA 95124


A bi-monthly 88 -page magazine for students, hobbyists, and anyone curious about computers. Fiction, articles, humor about computers, cybernetics, careers, building info., etc. Emphasis on games, puzzles, and projects. Contemporary, non-technical approach. Subscription. $\$ 8.00 \mathrm{pp}$.

## Games \& Puzzles Issue of Creative Computing

88 pages of games and puzzles for pocket calculators, computers, and humans. "Beating the Game," "Computer Chess," "Hunting a Wumpus in a Cave," building your own computer, reviews of 24 games, books, and much more! $\$ 1.50 \mathrm{pp}$.
Futures Issue of Creative Computing Artificial Intelligence (Bertram Raphael, Herbert Dryfus, etc.), Extraterrestrial Intelligence (Isaac Asimov, Martin Harwit, etc.), microprocessors, videodiscs as an ultimate computer input device, 4 new games, and more. 88 big pages! $\$ 1.50 \mathrm{pp}$.

## Artist and Computer

A high-quality, 4-color book edited by Ruth Leavitt which displays the work of 35 internationallyknown computer artists. Each artist describes his or her work in non-technical terms. 140 illustrations. \$4.95.


# Letters 



THE POLAROID HARD COPY UNIT?
1 want to suggest a cheap, portable, multiuse, easy to interface, and inexpensive hard copy unit for BYTErs with video terminals - the camera. I suggest:

1. A multipurpose camera loaded with cheap black and white film for photographic listings.
2. That the camera and film should be compatible with the home production of microfiche (a viewer can be built for $\$ 15$ - reference on request). A high resolution film will be desirable to make eventual computer scanning of your fiche easier (assuming that you will make fiche of things not in the computer system already - quite likely since costs can be less than $1 \phi$ per page).
3. Personally, I would like to add a Polaroid SX70 for instant color (\$1 per picture) copies of Dazzler displays.
Last - since I know nil about photography, would some BYTEr write an article about photographic copy systems for home computers, preferably with all the information a novice needs to get a system up and running, plus cost information for the cameras, film and developing needed?

## Martin Buchanan 2040 Lord Fairfax Rd <br> Vienna VA 22180

## AN OPINION ON SOFTWARE MARKETING

Let me first say that 1 write as a concerned hobbyist. My comments are my own and do not necessarily reflect the opinions of any organization of which I may be a member.

Bill Gates" "An Open Letter to Hobbyists" very clearly explained one of the chief problems of the hobby computer industry, the low Return on Investment (ROI) on the
software component of the system sold [for a copy of Bill Gates' views, see page 14 of Radio Electronics, May 1976, page 24 of March-April 1976 PCC (Box 310, Menlo Park CA 94025), page 3 of February 1976 Computer Notes (published by MITS Inc) ].

I have no solution to this problem. I only wish to express my personal views on three related aspects of the computer hobbyist market: recovery cost per unit, software delivery methods, and target market expansion.

## Recovery Cost per Unit

When was the last time you saw the word "profit" written on a rest room wall? Profit is not a four letter word. It is the incentive that brings vendors into this insane world of home computing. It is no more evil than indirect addressing or fig newtons. Assume that a vendor adopts the following profit philosophy.

Profit will be generated from the hardware sold. Since hardware is of no use without good software, a high initial capital investment will be made in software. This cost will be partially recovered with each unit sold. Recovery cost per unit must be as low as possible to discourage unauthorized duplication of software.

You can see from this that
(Recovery cost per unit) is equal to (Total software cost) divided by (Number of units sold)
Therefore, to minimize recovery costs, vendors must view the hobby market not in terms of hundreds or thousands, but tens of thousands of units.

One problem with a volume of this size is that present software delivery methods of cassettes and paper tape will not be economical. New delivery methods must be used.

## Software Delivery Methods

A high volume software delivery method should be based on proven technology, use agreed upon standards (such as Kansas City), be easily adaptable by all hobbyists, and have a low per unit cost in high volume.

One method that may work, although it's probably not the best, is the use of standard $331 / 3 \mathrm{rpm}$ long playing phonograph records. It meets all of the requirements from above. Interfacing would be done from the turntable to the cassette input. A transformer may be needed for impedance matching. I'm sure other hobbyists will devise better methods.

Now that the need for a high volume of software and a possible delivery method
have been established, the question remains of how to create a market to justify this.

## Target Market Expansion

The best approach to increasing the number of computer hobbyists is to publicize our activities in a medium that is accessible to the general public. Once people see what we are doing with a few $K$ of memory, a couple of LEDs, and a lot of sweat, I'm sure it won't be long before they too are "hooked."

I therefore propose that hobbyists, individually and through their local clubs, start writing to their local Public Broadcasting television station suggesting that they join with other stations to produce a special on th is "world's greatest toy."

Who knows, we may one day be able to enjoy a weekly half hour program on home computing while our micros are storing the latest game of the week, by frame grabbing, off the same television signal.

To summarize, I believe software costs can be reduced considerably by increasing the size of the market by at least an order of magnitude, and devising new delivery methods to meet that market. It remains for hobbyists to inform the general public of our activities so that all this may happen.

What are you waiting for? Fire up your pencils, pens, 33 ASRs, and matrix printers. Send that letter off now. Address it to the director of programming in care of your local PBS station.

You might even begin it with, "As one programmer to another . . . ."

Robert H Wada<br>Garden Grove CA

## LET'S HAVE A SNOBOL FIGHT (BUT ICEBOLs ARE NO FAIR)

[^3]
## TH:S IS DUR CDPYDAT CDPYCAT

 HER NAME :S TCHEBYCHEU ...

This month we are offering a PRIE to the first to correctly guess what our Copydat Copycat is doing. The prize: your choice of one of our P.C. cards listed below. (Hint: it's related to our newest product line, the COPYDAT blueprint copiers.) Next month we print the answer and the prizewinner. Send in your guess today!

The "old standby" line of Hardware Assemblers:
CDA 1.1 4KX8 Memory Matrix Prototyping Card for 2102-2602-9102 Static RAMS. 7" $\times 9^{\prime \prime}$, predrilled wire wrap custom decoding area and plated-through holes in RAM area. \$34.95.
CDA 2.1 General Purpose Prototyping Card, predrilled for wire-wrap, pictured in BYTE Dec. '75, 7" $\times 9^{\prime \prime}$, $\$ 34.95$.
CDA 3.1 Digital Graphic Display for Oscilloscopes designed by Jim Hogenson and featured in BYTE Oct. '75, double-sided card with plated-through holes. $\$ 29.95$.
CONNECTOR - For CDA 1.1, 2.1-70-pin wire wrap, Viking NORSMAN ${ }^{\circledR}$ - $\$ 5.00$. (Postage and packaging per card: $\$ 2.00$.)
Our newest line - the COPYDATS, starting at \$149.95 FOB Amherst, N.H.

COPYDAT $11-$ for copies up to $24^{\prime \prime} \times 36^{\prime \prime}$
Copies for less than 36 each (Copydat I) - without sacrificing quality.
Send for a brochure and sample print today!

## CELDAT DESIGN ASSDCIATES <br> BOK 352 AMHERST,N.H. O3O31

very little computing, operating mainly by table lookup and short library routines. Only the Top Processor seems capable of real time operation. The biggest glitch in the whole system is the Value Judgement Table. The Main Processor was never designed to handle value judgements, but insists on trying. No wonder an external Diagnostic Processor is needed!

Great article on the TMS9900 [April 1976 BYTE, page 64], but a bit too short. How about publishing a similar one on the RCA CDP 1800?

## J Gordon Los Angeles CA

## A STAR TREK PRODUCT

Several letters have recently appeared in BYTE concerning the history of the Star Trek game. A version of the game is now available for the Altair 8800 . It is written in Altair 4 K BASIC and is available from International Data Systems Inc. The purchase price of $\$ 10$ (checks OK) includes a complete program source listing, operational instructions, tips on how to "patch" the program to add your own features, a one year limited warranty against "bugs," and postage and handling. The limited warranty states that a corrected copy will be provided if any errors are identified. The limited warranty does not provide replacement for time or resources lost as a result of such errors. Orders should be addressed to Star Trek Offer, International Data Systems Inc, PO Box 593001-AMF, Miami FL 33159.

## D E Hipps <br> Vice President <br> International Data Systems Inc

## SPACE WAR ORIGINS UNVEILED

l'd like to answer the comment in $A B$ Bonds' letter (page 8 of April 1976 BYTE) on the "mysterious" origins of Space War. I believe that the first Space War was implemented on a PDP-1 at MIT. This machine was one of the first models of one of the first minicomputers ever made. It was donated by DEC to MIT where it created a new type of student - the "computer hacker." Several members of the Tech Model Railroad Club attached themselves to this machine, developing, besides the original Space War, one of the first timesharing systems, predecessors to DDT and TECO, and I believe the first LISP implementation. This Space War had a sun with gravity, torpedos, and hyperspace. It ran in 4 K words of memory.

Several of the PDP- 1 hackers went on to the MIT Artificial Intelligence Laboratory when it started up with a PDP-6. Space War came with them. My introduction to the game at MIT was on the AI lab's PDP-6. That version was similar to the PDP-1 version. One of its new features was an accurate star field in the background (with proper magnitudes) that moved slowly as the game progressed and the observer moved around the sun.

I have since seen a Space War implementation on a PDP-9, a PDP-7 and even on an IBM 7040/7094 DCS. The latter was a large scale second generation batch system at Yale. The only way we could take over the machine to play was with the cooperation of some third shift operators who were also Space War addicts. We would check the control cards of the night's runs to see how much time we could spend playing and still get all the users' jobs done by morning. We would then send out for pizza and start playing.

The best version of Space War I have yet seen was developed on the MIT AI Lab's PDP-6 by a person with the login ID KLH. This is a Star Trek-y game, with one ship looking like the Enterprise and optional use of phasers and deflector shields instead of torpedos and hyperspace. This game also has space mines and invisibility shields.

## Sidney Markowitz Cambridge MA

Then, of course, there are the MIT undergraduate courses in digital electronics which come up with a new variation of Space War as a lab project each year.

## MORSE POWER TO YOU

I have just received my April issue of BYTE and regarding Richard Fall's request for information on Morse translation algorithms: I wish him luck. There is a remarkable dearth of information available to the average person. I have been working on my master's thesis and have been researching this area for over a year. I will try to save Richard and others like him the trouble of looking for information that is not really there. Virtually all of the information available today is either military related (that generally means classified!) or is closely guarded by civilian companies that are hoping for government contracts (that generally means classified too, or even worse, proprietary to the company).

There are a few articles on hard wired Morse decoders in amateur radio magazines
from which a simple mathematical model can be derived, but I caution the overzealous experimenter for there is probably no ONE satisfactory algorithm to translate HANDsent Morse code. Du to the high degree of variability of mark/space characteristics (for those who don't care for the higher math terms, this means that the lengths of dots, dashes, and the various spacing intervals are constantly changing) ANY fixed algorithm will eventually fall apart and fail to decode properly mainly because we do have some pretty rotten fists on the air today!

Obviously, the more complex the algorithm, the more memory it will require. Some professionally designed and programmed algorithms I have seen on PDP-11-type machines run in the vicinity of 8 K for a BASIC decoder. This figure can go as high as 250 K or more for the more sophisticated models, but in most cases the simpler models will work fine for jisting copy (getting the general trend) and are generally unsatisfactory if your interest is in solid copy.

Machine-sent Morse and also RTTY are completely different problems, since they represent machine-to-machine interface rather than man-to-machine interface. I would advise interested readers to look up: (1) November 1971 Ham Radio magazine (2) January 1971 QST magazine (3) Oc-tober-December 1975 QST and either or both of the following papers if they are available at your library (try college libraries too):
(1)"On Computer Transcription of Manual Morse," Journal of the Association for Computing Machinery, Volume 6 \#3, July 1959 (?), pp 429-442, by Charles R Blair.
(2)"Machine Recognition of Hand Sent Morse Code," IRE Transactions on Information Theory, March 1959, pp 17-24, by Bernard Gold.
These articles are NOT simple, but should provide some insight to the magnitude of the problem of Morse Code translation.

W A Hickey
Naval Postgraduate School Monterey CA

TAKE A STAND?
Is the computer hobbyist getting state of the art components? How many years must pass before the hobbyist gets new developments?

EXAMPLE: Electronics magazine, January 8 1976, page 76: Matsushita Electrical

## MERRIMAC PROCESSOR SYSTEMS ANNOUNCES THE CRTV

A complete low cost Teletype ${ }^{\circledR}$ Model 33 KSR replacement using a standard TV as the monitor. Complete, whether purchased as a kit or assembled.
(1) Full cursor controls - clear, home, cursor up, down, left, right.
(2) Infinite scrolling.
(3) Selectable data rate to 600 baud.
(4) Half/full duplex; local mode.
(5) "Break" function allows entering many debug routines.
(6) Model 33 keyboard with added cursor functions, clear, and home keys.
(7) Housed in attractive enclosure.

Completely assembled and tested . . . $\$ 495.00$
Complete kit . . . . . . . . . . . . . . . . . $\$ 355.00$
For Complete Information write -
Merrimac Processor Systems
P.O. Box 76

Chester, New Hampshire 03036

## mp-40 THE ONLY PRINTER



- PARALLEL INTERFACE - $5 \times 7$ IMPACT DOT MATRIX - 64 CHARACTER ASCII - 40 COLUMNS AT 12 CHAR./IN. - 75 LINES PER MINUTE - ORDINARY ROLL PAPER COMPLETELY ASSEMBLED AND READY-TO-RUN SEND FOR FREE LITERATURE

BYTE's election year stand on motherhood, apple pie, computer power for the people, and state of the art components for hobbyists.

Industrial Company has developed a floppy disk memory at the price of a digital cassette.

EXAMPLE: Calculator and computer companies are using 4 K RAMs now. Some use 4 K RAMs in designs that permit an exchange for 16 K RAMs when they appear in quantity this summer.

EXAMPLE: Electronics magazine, January 22 1976, page 40: Axiom Corp is manufacturing a microprocessor controlled line printer at a cost of $\$ 1000$ in a single quantity, and $\$ 500$ in OEM quantities.

When do we get these components in our kit versions?

Where does BYTE magazine stand on this issue of state of the art components hobbyists?

Edward L Tottle Baltimore MD

To paraphrase Ted Nelson's introduction to his talk at the MITS World Altair Computer Convention: "Motherhood!"' "Apple Pie!" "computer power for the people!". Of course the computer hackers are getting state of the art components. State of the art means current technology which both works and is manufacturable and salable. As for timing of specific items based on what you see in electronics engineering magazines, that is up to the entrepeneurial designers who create actual products which can be marketed to the personal computing consumer. So, as with motherhood, apple pie and computer power for the people, BYTE is definitely in favor of advances in the state of the art.

## SUPER SERVICE TESTIMONIAL

I don't know how many computer hobbyists are at the hardware experimentation stage. If the computer club here in Columbus OH is any indication most are probably still at the thinking and planning level, but for those who are building I would just like to provide a little "reader feedback" on the service from a couple of your advertisers. All us hardware types are familiar with the "back order blues" that holds up experimentation for months, so the ultra fast turn around service I got from

Advanced Data Sciences (they sell keyboxes) and S D Sales (they sell parts) deserves special credit.

Since it was a letter just like this one in an amateur radio magazine that put me onto $S$ D, I just thought a little feedback among computer types might be of a little interest for those who are about ready to buy.

## B F Jacoby

88 W Frankfort
Columbus OH 43206

## WHERE IS THE INEXPENSIVE GRAPHICS DISPLAY CIRCUIT?

I can't understand it! There are so many computer hobbyists that I know without any 10 . And yet we can find no circuit that will take our 8 bit output and put it on our TV sets.

There are many TV typewriters, TV games, TV graphics (most in the range of $\$ 160-\$ 300$ ).

Where is the circuit (inexpensive) that will accept 8 bits of $x$, $y$, brightness (whatever the TV will need) and convert it to video for my portable TV?

My 6800 is waiting, as are many others!! This circuit would be more versatile than a dedicated TV typewriter or game.

Let the computer owner make it to fit into his system to accommodate his needs. Let him decide on where the memory shall reside.

There is a need! See what you can do!

> Paul Hyde Jr Milwaukee WI

Take a look at "Build a Television Display" by C W Gantt in the June 1976 issue of BYTE, page 16. If "will accept 8 bits" is loosely interpreted as an ASCI/ character, then Mr Gantt's circuit fills the bill.

However, the author-engineers in our audience should take a cue from this letter of Mr Hyde. Here is a functional specification challenge:

Design a 256 by 256 point display module which maps 216 bits of memory (organized as $213=8 \mathrm{~K}$ bytes) into a standard EIA composite video signal which will drive a TV monitor. We'll all look forward to the first such high resolution plotting module since it will enable us to do detailed graphics for a "reasonable" price. But TANSTAAFL (There Ain't No Such Thing As A Free Lunch). The minimum price you can expect on such a graphics display is the cost of the memory and the cost of the video monitor combined.

## Ask BYTE

What computer systems are available that would store five thousand $(5,000)$ patient records? Each record would consist of name, address, telephone number, age, prescriptions and date of last examinations.

## Dr Sydney B Schrum Goldsboro NC

## BYTE Replies:

A first step is to calculate specifically what your data base requirements are. Then what you want in the way of access and manipulation of the data. In your case, picking numbers which look "typical," suppose each record is (worst case lengths):

| Name: |  | 30 bytes |
| :--- | ---: | ---: |
| Address: | line \#1 | 30 bytes |
|  | line \#2 | 30 bytes |
|  | zip | 5 bytes |
| Telephone: |  | 10 bytes |
| Age: | 2 bytes |  |
| Prescriptions text: |  | 30 bytes |
|  |  | (?length) |
| Date of last exam: |  | 6 bytes |

A fixed length record format based on this information would require 143 bytes per patient. Multiply this figure by 5,000 to find the total storage requirement of 715,000 bytes. If you use a varying length record format, this number can be shrunk considerably (perhaps 20-50\%); since the typical name, address, prescription, etc, will not use all of the worst case length allowed to it.

Now, what does this mean in terms of system hardware? The implications depend upon how you want to get at the data. If you only intend to access the data serially in a personal computer version of a "batch" tape oriented system, several channels of audio cassette would be an inexpensive but slow access method. Assuming a high performance cassette interface at 1200 baud, and a 12 bit UART data format (start, 8 data bits, 1 parity bit, 2 stop bits), it will take you a total of 7150 seconds or nearly two hours to read, process and write the data in an update run. It would thus take two hours worth of cassettes or 4 standard C-30 cassettes, and you would have to change tapes every half hour.

If you want "instantaneous" access, then a higher cost floppy disk system would be required. With one floppy drive, you can
store "on line" typically 200,000 to 250,000 bytes depending upon formatting. Thus you would have to spread your files out over about 3 to 4 diskettes in order to store all 715,000 records. "Instantaneous" would work out to be the time required to load and start the proper disk then access the desired patient record. If all the records are to be stored on line at one time, three to four floppy disk drives are required (or two dual disk drives).

Ancillaries you will need to make your software development convenient in either access method are a high level language like BASIC with string capability (to handle names, addresses, etc) and probably 4 K to 8 K or programmable memory work space in addition to the space required for the BASIC interpreter. You should also have a hard copy printer such as a Teletype if you want to program your patient billing and appointment reminder messages. $■$

Some letters to BYTE are technical in nature and are best served by some form of response. We will try to answer as many such letters as possible. If you have a puzzle concerning some aspect of the personal computing field, write down a clear statement of your question and send it to:

Ask BYTE
Byte Publications Inc
70 Main St
Peterborough NH 03458.
We will publish names and addresses of individuals making inquiries unless you specifically request us to omit the reference.

```
Introducing A New and Unique
Computer Graphics Terminal To Fit Your Altair
* B+W Matrix of 128 Horizontal x 192 Vertical Dots
* \(3 \mathrm{~K} \times 8\) No Wait, Static RAM, ON BOARD
* Light Pen And Control Panel Included
* Bandwidth is Compatible with standard TV Set
* Output is 2.25 V PP Video
GDT-1 is a 2 slot Altair plug-in graphics terminal that generates 24 K dots on a standard B+W TV set. Unit displays data stored in on board \(3 \mathrm{~K} \times 8\) memory, which is fully computer accessible. Control panel allows data to be entered by light pen or computer.
```

Kit \$185., Assembled \$235
Computer graphics, Associates
56 Sicker Road
Latham, N.Y. 12110
N.Y. Res. add Sales Tax - Add $\$ 2.00$ Shipping Send $\$ 1.00$ for Data Pkg.

To get further information on the products advertised in BYTE, fill out the reader service card with your name and address. Then circle the appropriate numbers for the advertisers you select from this list. Add a 9 cent stamp to the card, then drop it in the mail. Not only do you gain information, but our advertisers are encouraged to use the marketplace provided by BYTE. This helps us bring you a bigger BYTE.

| A75 | Advanced Microcomputer Products 77 | A15 | James 97, 105 |
| :--- | :--- | :--- | :--- |
| A70 | American Microprocessor 77 | A90 | Logical Services 73 |
|  | BYTE's Books 82, 84, 102 | A94 | Merrimac 93 |
|  | BYTE's Subscriptions 80 | A18 | Meshna 110, 111 |
|  | BYTE's Poster 96 | A93 | Micon 73 |
| A88 | Cambion 77 | A95 | Micro Data 89 |
| A37 | Celdat 91 | A71 | Micro Peripherals 93 |
| A89 | CFR Associates 73 | A57 | Mikos 29 |
| A91 | C. Itoh 83 | A20 | Mikra-D 85 |
|  | Computer Art 78 | $*$ | MITG CIV, 45 |
| A92 | Computer Graphics 95 | A22 | National Multiplex 43 |
| A81 | Computers and Stuff 73 | A64 | Oliver Audio Engineering 77 |
| A83 | Creative Computing Books 79 | A98 | On-Line 77 |
| A87 | Creative Computing Subscriptions 89 | A63 | Parasitic Engineering 85 |
| A41 | Cromemco 1 | A23 | Polymorphic Systems 49 |
| A85 | Custom Design Services 77 | A24 | Processor Tech 56, 57 |
| A97 | Data Domain 75 | A26 | Scelbi 17 |
| A7 | Delta 101 | A27 | SD Sales 108, 109 |
| A78 | Digital Group 5 | A59 | Solid State Sales 103 |
| A8 | Dutronics 81 | A29 | Southwest Tech CII |
| A9 | Godbout 99 | A30 | Sphere 21 |
| A69 | HAL Communications 51 | A96 | Synchronous Sound 73 |
| * IMS Associates, Inc CIII, 39, 112 | A32 | Tri Tek 107 |  |
| A38 |  |  |  |

* Reader service inquiries not solicited. Correspond directly with company.




## Classified Ads for Individuals and Clubs

Readers who have equipment, software or other items to buy, sell or swap should send in a clearly typed or printed notice to that effect. The notices are free of charge and will be printed one time only on a space available basis. Insertions should be limited to no more than 100 words. Notices can be accepted from individuals or bona fide computer users' clubs only. Commercial advertisers should contact Virginia Peschke at BYTE for the latest rate card and terms.

Before sending your classified ad to BYTE, read it over. Did you include your name, address, phone number (with area codel in the text of the ad? BYTE has received several ads with incomplete phone numbers or missing addresses.

WANTED: Any low cost computer terminal for private use. Please contact Robert Brodie, 15 Harwich Rd, Chestnut Hill MA 02167, (617) 332-8034.

The Beta Terminal owners group of the Computer Hobbyist Group of North Texas is interested in establishing communications with owners of terminals that use the Univac 0769 Series print mechanism. We are looking for Beta keyboards (Microswitch No. 53SW1-2). One of our members has some spare parts for Beta terminals for sale. Contact L G Walker, Rt 1 Box 272, Aledo TX 76008, (817) 244-1013.

WANTED: Complete working computer or one still in kit form. Need E\&L LR-outboards and E\&L mark 80 microcomputer system or its equivalent. State condition and give lowest price. John Waskowitz, 35-30 73rd St, Flushing NY 11372.

WANTED: New or used: computers; peripherals; components; plans; software; robots; laser weapons; science fair projects; Star Trek items; rockets; electronics equipment; amateur radio; antennas; satellites; whatever you have I will make a cash offer. Please send your specifications to Joe Halligan, 3331 Altamont Dr, Wilmington DE 19810.

WANTED TO MEET: Individuals using the Univac 769 printers for hobby purposes. I have knowledge of how to use these printers and will share same. FOR SALE or TRADE: 3 serial punch devices from Teletype Corp, new, very similar to the ASR 35 punch units; take 20 mA current loop input of ASCII and require a drive motor. Jim Beistle, 3728 Wilkie Way, Ft Worth TX 76133.

FOR SALE: IBM Selectric Printer/Keyboard unit (Serial \#4600785). Originally used as console for IBM 1410. Has 1407-type keyboard, mounts like 2741 in a table well. Convertible to computer console, terminal, or Selectric typewriter with proper user-supplied power supply and control logic. Needs minor repair. Sorry, no FE manuals or data sheets. Weight: approximately 40 lb . Best offer above $\$ 75$ takes it. M Scherer, RR 2, Box 116, Bedford NY 10506.

1702s PROGRAMMED - I can program your 1702 s for $\$ 3$ or duplicate for $\$ 1.50$. Changes in a duplication $10 d$ each + duplicating charge. Also, I have a limited number of 4 KE PROM kits for the Altair/Imsai that hold up to 16 1702s for $\$ 65$ each (1702s are not included). Please send material to be programmed in HEX (Base 16). Include return postage. John B Jay, 165 River Valley Rd NW, Atlanta GA 30328.

FOR SALE: Wurlitzer Electric Piano. Uses standard piano action to strike tuning forks with electronic pickup. Sounds like an amplified clavichord. Has controls for volume and optional vibrato; headphone and auxiliary amplifier jack. Best offer over $\$ 550$. Call Carl Helmers, BYTE Publications, (603) 924-7217, or write to 70 Main St, Peterborough NH 03458.

FOR SALE OR TRADE: Friden Flexowriter Model FPC-8, 8 level code printer, keyboard, paper tape punch and paper tape reader with table. Excellent condition \$500. Also disk memory, 27 track fixed head, 5454 bits per track $\$ 100$, HP120A oscilloscope with 2 Tektronix probes \$150. J G Hansen, PO Box 1337, Fort Davis TX 79734 (915) 426-3331 (evenings).

WANTED: Humane individual to provide back issues of BYTE from "GO' to number 5 inclusive. Will pay all costs and bribery. P L Christie, 20 James St, Adelaide, 5000, Australia. II'm not necessarily interested in original copies, only the information therein. [BYTE won't mind if some individual makes an international good will gesture by loaning Mr Christie copies of BYTE \#1 to \#5 for the purpose of limited Xerox reproduction.])

FOR SALE: DEC PDP-8/S, one working, good condition, one for parts (both \$2000), IBM 024 keypunch for experimenter $\$ 150$. NCR card reader, new-unused, EM-D2, $\$ 100$. Call or write Reg Conkling (516) 744-9475, Box 310, Shoreham NY 11786.

ALTAIR 8800 computer, fully assembled and operative, with 4 K static RAM, \$575. MITS VLCT, assembled and working, with non-MITS keyboard and no case, $\$ 85$. You pay shipping, I pay crating charges. James R Einolf, 12149 N Piney Lake Rd, Parker CO 80134. (303) 841-2105.
"Hints for troubleshooting and suggestions for application of the Viatron 21 Data Management Station." Write Greg Ludwig, Box 408, Rice Lake WI 54868.

Will swap my time to help you set up, program, develop system(s), or to give you training in the above, in return for using your system. Or would be interested, also, in swapping my time to secure a system of my own. Have an extensive background in systems/programming, many systems, many languages, including assembly and higher level. Please contact Toby Maki, 102 Minott Rd, Westminster MA 01473, (617) 874-5410.

Will trade comics, bicycle magazines, or pre-1950 radio magazines for microcomputer magazines or hardware. Ian MacMillan, PO Box 128, Mount Royal, Quebec, Canada H3P 3B9.

FOR SALE OR TRADE: 5 channel TTY; IBM SMS logic; CDC 2816 office typewriters with optical paper tape reader and high speed paper tape punch and controller (paper tape version of MTST). W Palya, St Joseph's College, Rensselaer IN 47978, (219) 866-7111.

FABULOUS SALE! Spectrum Analyzer unit 10 cps to 100 kHz expand to MHz . Displays spectrum on any scope with documentation, cost \$1475 new, sell for $\$ 550$ or trade. VIATRON 21-PARTSI Parts, documentation and all the goodies including RS-232 interface. Can provide repair service. $R$ Brown, PO Box 83, Marcy NY 13403.

Bell 103 compatible Acoustic Couplers. TTL compatible data lines, single +12 volt supply. Optional signals provided: Clear to Send, DAA input. Switch selectable Up/Inverted code. We have used these with HP, DEC and Burroughs timesharing. Includes schematic and interfacing information, \$50. George H Enos 2400 W Valley Pkwy, Sp 42 Escondido CA 92025.

## From the same people who brought you Vectorboard, the time-saving Wire Pencil, and the Altair Universal Protptype Board comes a truly revolutionary tool. The "Slit-N-Wrap" tool provides---for the first time---a fast manual wrapping device which supplies insulated wire from a spool directly to wrap posts WITHOUT any external precutting or prestripping. It can do this in a "daisy chain", so more than one level is rarely required...thus saving space when you want to stack cards. Wires are cut off only at the end of lines--no premeasuring of wires is necessary. The secret? An almost undetectable slit is cut into the wire as An almost undetectable slit is cut into the wire as it leaves the tool's tip; as the wire pulls past the wrap post, the copper is exposed to the post. Quality of wrap is extremely good---another advantage.



The Slit-N-Wrap package comes complete with tool, 200 feet of wire, and instructions. Add 1 pound shipping, please.

## A FULL FEATURE



Now that other companies have picked up on our buffering and static operation...and on adding hysteresis to the data in lines (for unambiguous
$5-9.95$ data transfer)...we decided to continue staying ahead of the pack by giving more features at
a lower price. Like our previous memory kits, this one is FULIY PLUG-IN COMPATIBLE WITH BOTH IMSAI AND ALTAIR 8800 MICROCOMPUTERS. The board features buffered address, data out, and data in lines...zero wait states... 1A typical power consumption...low power Schottky for all TTL support
ICs.. on board regulation with 3 regulators, situated for ICs... on board regulation with 3 regulators, situated for optimum thermal stability... lots of bypassing ( 39 disc ceramics, and 7 tantalums for power supply decoupling). ceramics, and 7 tantalums for power supply decoupling)... ty, double-sided plated through PC board (epoxy glass base like our other kits)...sockets for all ICs... plus ease of like our otner kits)...sockets for all ICs... plus ease of
assembly, thanks to our simple-to-follow instructions.
Documentation also includes a logic print of the memory. Documentation also includes a logic print of the memory. Others can incorporate the engineering, but we've got the recipe on keeping the price down.
$\square$
 0000 SH. 95 EACH. EXCELLENT FOR FREQUENCY CALIBRATORS, hams, or computer bugs. Series mode, fundamental operation, hire leads. Choose from the following freauencies:
4 MHz (PaCE clock frequency), $5 \mathrm{MHz}, 8 \mathrm{MHz}, 10 \mathrm{MHz}$, $12 \mathrm{MHz}, 15 \mathrm{MHz}, 18 \mathrm{MHz}$ ( 8080 clock frequency), and 20 MHz .

8080 CHIP SET
THE 8080 IS A HOBBYIST STANDARD IC. BUT A MICROPROCESSOR BY ITSELF DOESN'T MEAN MUCH WITHOUT SOME MEMORY, AND PREFERABLY SOME KIND OF MONITOR ROUTINE SO YOU CAN AT LEAST Fill and examine the memory. So, here's WHAT YOU GET IN OUR CHIP SET:

## 1 - TMS8080 MICROPROCESSOR

8 - 2102 1K STATIC RAMS
1 - 5204 4K EROM (eraseable, programable)
The cost of this chip set is $\$ 59.95$. HowEVER, BY ADDING $\$ 10$ WE'LL PROGRAM YOUR 5204 WITH THE 8080 MICRO-MONITOR. THE SOFTWARE IN THE 5204 ALLOWS YOU TO EXAMINE AND FILL MEMORY, PUNCH AND READ PAPER TAPE, BRANCH TO---AND RECOVER FROM---PROGRAMS. A COMPLETE SOFTWARE LISTING COMES WITH THE
MICRO-MONITOR. THE MICRO-MONITOR IC IS AL-
SO AVAILABLE SEPARATELY FOR $\$ 29.95$.



8008 CHIP SET $\$ 32.50$
(1) 8008 CPU AND
(8) 2102 s
No data other
THAN PINOUT INCLUDED.

PACE BASIC CHIP SET (16 BITS) $\$ 125.00$
(1) PACE CPU,
(4) DS3608..
(1) DS0026...
(1) DM8837.,

PACE DELUXE CHIP SET $\$ 195.00$
(1) PACE CPU...(4) DS3608...(1) DSO026...
(1) DM8837...(32) $2102 \mathrm{~s} . .(4) 5204$ EROMS,
AND OUR PACE DATA PACKET.
(1) DM8837...(32) 2102s...(4) 5204 EROMs, and our Pace data packet.

The best thing in life is free, so they say...but so are some of the second and third best things, like our flyer. Send us your name and address if you're in a hurry, or circle the reader service card if you can afford to wait a little longer.

Here lies documentation of known bugs detected in previous editions of BYTE. . .

## Patching the Biorhythm Program

Don Libes, 995 Chimney Ridge, Springfield NJ, reports that the biorhythm program in the April issue of BYTE (page 22) incorrectly calculates a person's life span. He suggests the following patch to correct the problem. Delete lines 130 to 470 of the listing as printed, and replace them with the following set of lines:

```
130 FOR X=1, M-1
140 D=T(X)+D
150 NEXT X
160 IF M<=2 THEN 180 or IF Y/4-INT(Y/4)
    # O THEN }18
1 7 0 \mathrm { D } = \mathrm { D } + 1
180 D=(Y-1)*365+1NT((Y-1)/4)+D
190 D4=D1
200 FOR X=1,M1-1
210 D4=T(X) + D4
220 NEXT X
230 IF M1=<2 THEN 250 or IF Y/4-INT(Y/4)
        # O THEN 250
240 D4=D4+1
250 D4=(Y1-1)*365+INT((Y1-1)/4)+D4
260 D3=D4-D
```

With this patch, all cycles should now start in phase at one's birthdate and calculate the biorhythm correctly according to the biorhythm hypothesis as stated. However, this patch does not affect the question of whether the biorhythm hypothesis as stated and programmed is valid.■

## How to Do It Better

Michael M Dodd, 291 Waples Estates, Fairfax VA 22030, sent us a lengthy note commenting upon the design of Robert Bosen's interface, page 42 of April BYTE. It is reprinted here with a couple of editorial notes interjected. The key to the notes is as follows: [Oops] indicates a design flaw that should have been corrected editorially; [Point] indicates a possible better way to have done something; [Good point] indicates a place where an engineer with considerable experience would have done things differently, and the reasons why. Michael's letter is printed as much for the two real [Oops] class errors as it is for the tutorial value of his comments concerning design practice.

In reading the article, "Controlling External Devices With Hobbyist Computers" by Robert Bosen in the April issue of BYTE, 1 noticed several cases of what I consider to be poor hardware design. I would like to point them out to possibly help someone wanting to build a similar circuit.

1. The IC inverting bit D4 in figure 1 should be a 7404 , not a 7400 (the pin numbers are wrong for a 7400 ). [Oops] In addition, it is generally not good practice to parallel TTL gates; if you need more drive, split the load. It is possible that one of the gates could produce a slightly lower zero voltage and take all the current while the other one did nothing. [Good point]
2. Pin 10 (the preset input) and pin 4 (same) of the 7474 flip flops should be tied to +5 V ; leaving it open is inviting false triggering due to noise pickup. /Good point/
3. Driving the base of transistors with TTL outputs can cause overheating of the device. A three volt drop across a 1 k resistor causes 3 milliamps to be drawn from the TTL output in the high state; this is far in excess of the 400 microamps specified. [Good point]
4. The jumpering scheme in figure 2 is wrong. It will cause the outputs of the inverters to go low when the inputs are left floating; this will short the address bus lines to ground. The jumpers should be between the inputs of the 7430 and either the outputs of the inverters or the address bus directly. [Oops]
5. Mr Bosen's point about saving address loading by using low power TTL is well taken. The second input of any 7400 used as an inverter should be tied to +5 V to reduce loading, not to the active input. This will save one input load on the address line. [Good point]
6. In figure 3a the diode across the relay coil does not have to be a Zener. When the magnetic field collapses, a spike is indeed generated, but it is of the opposite polarity; thus, a common silicon diode (polarity as shown) will shunt the spike, protecting the transistor. [Good point]
7. On the software side, I might point out that with the "clear" line on bit D4" (figure 1) connected as shown, a zero on D4 will cause a "clear" condition. Many people like to work with active high (a one performs the operation) and one inverter is required to achieve this. It sometimes helps in programming if you use active high. [Point]
8. Also on software, the description about clocking another output channel to get the pulse necessary to toggle the flip flop is not really necessary. The chip enable inputs to the 75154 (figure 1) are designed to put all the outputs in a high state when the chip is disabled. All that is necessary to get a pulse is to tie either pin 18 or 19 to an unused bit on the output port (there are three left) and wiggle that bit while maintaining the channel selection number on the lower four bits. This may or may not be easier than selecting another channel, but at least you don't have to worry about what channel to select to get your pulse. [Point]

I just wanted to point out these items. Since Mr Bosen has obviously built the device, I can't argue with success; but some errors must have crept in between his wire wrap tool and your pages.-


## TRANSFORMERS

We carry one of the largest \& most diversified stocks of transformers in the country. Listed below is a sampling of our many different types. Our free catalog lists many more.

PRIMARY
SECONDARY NO. 1
115 v or 230 v
$115 v$
115v
$115 v$ or $230 v$
115v
$115 v$


AUTOTRANSFORMERS: 2 KVA rating, voltage taps every 5 volts from $85 v$ to $125 v$. Valuable in areas of low or fluctuating line voltage. $41 / 2 \times 51 / 2 \times 41 / 2,15 \mathrm{lbs}$. STOCK NO. B9721 $\$ 9.95 \mathrm{ea}, 2 / 18.00$
2 KVA rating, step up/step down autotransformer. Taps at $100 \mathrm{v}, 115,125,200,220$, and 230v. $33 / 4 \times 41 / 2 \times 41 / 2,11$ lbs.

# Viatron System 21 PC BOARDS 



B6362


Recently, Viatron System 21 Data Management Stations have become available from several surplus sources. While this is an excellent machine, the availability of spare parts is practically nil. We have these three PC boards available for this system. They are in new, unused condition, but do NOT have the LSI chips. We are offering them for their use as replacements or for parts.

KEYBOARD TYPE 2, number 109207.
DISPLAY ARRAY PARTITIONING, number 109300 Rev. B
STOCK NO. B6362
$\$ 9.95$ each

COMMUNICATIONS ADAPTOR BOARD, number 1099 Rev. B. Used in the 6003 and 6004 Communication Adapter,
which provides an ASCII interface to other equipment.
STOCK NO. B6364
$\$ 14.95$ each


Current factory price is $\$ 12$.

## Philbrick/Nexus OPERATIONAL AMPLIFIER

Philbrick/Nexus Model SQ-10a is a high gain, high stability general purpose op amp. It operates over the temp. range of $-25^{\circ} \mathrm{C}$ to $+85^{\circ} \mathrm{C}$. Both inputs \& output are fully protected against shorts to ground or the power supply. It may be used open loop as a voltage crossing detector. 1 Mhz unity gain bandwidth. Open loop gain $60,000 \mathrm{~min}$. Slew rate $1.1 \mathrm{~V} / \mathrm{sec}$. min. 100 nA . max input bias current. Nominal power supply $\pm 15 \mathrm{v}$. $11 / 8^{\prime \prime} \times 11 / 8^{\prime \prime} \times 1 / 2^{\prime \prime}$ stock No. B4539

With data
$\$ 1.75$ each, 4/6.00

[^4] be refunded. BANKAMERICARD \& MASTERCHARGE welcome, ALL numbers needed for processing. Minimum charge $\$ 15$.

# What's New? 

## But It's New to Me?

Yours truly [CH] had a little argument with photographer Ed Crabtree over this item, the "Hexadat" mechanical calculator available from Radix Precision Co, Atlanta GA. The "Hexadat" is a boon to all lovers of 4 bit representations of binary data as hexadecimal digits. Hexadecimal notation is the natural way to represent memory contents in lieu of other formatting conventions for any machine with a word length that is a multiple of 4 bits (but not divisible by 3 or else octal lovers will have a good argument).


But hexadecimal arithmetic is another matter, a technique which takes practice to make perfect or reliable as anyone with experience puzzling out an OS 360 or 370 core dump can attest. So, several years ago, the Radix Precision people created this neat little mechanical adding machine which is operated by a stylus and uses base 16 notation. Now that the microprocessor era of 8 bit processors is upon us, the same problem of hexadecimal addition and subtraction which afflicts users of large scale IBM machines is repeated in miniature for those of us who use hexadecimal for 8 bit processors. And this same mechanical calculator will provide excellent service for hexadecimal arithmetic in your own computer laboratory. It comes in a carrying case with a reference card for base conversions and the stylus used to perform operations.

Which brings the subject back to a certain argument. When Ed Crabtree returned with the pictures, he protested that "there's nothing new here!" Whereupon he produced from his pocket a neat little mechanical calculator circa 1900 which had been used by his grandfather. In "the old days" people apparently got along without the convenience of the $\$ 8.95$ hand cal-

Stamp Out Cybercrud

## COMPUTER



Have you every been victimized by one of a myriad computer based interpersonal putdowns? In Ted Nelson's book, Computer Lib/Dream Machines, you'll find an excellent essay on the nature of this "cybercrud."

Have you ever wondered where to go for a basic starting point in your quest for information about computer applications and uses? Ted Nelson's book, Computer Lib/Dream Machines, is the place for you to begin.

Computer Lib/Dream Machines is for the layman - the person who is intelligent and inquisitive about computers. It is written and self published by a philosopher who is also a self confessed computer fan and an excellent teacher of basic concepts. (For those who have not yet heard, ivory towers are constructed out of real and substantial white bricks.) The most important aspect of this book
is its inspirational data content. The machines we're all busy working on are deep personal expressions, and not the cold and inhuman monsters of the traditional stereotype. The book defines many of the terms and explains many of the techniques which can be used in the personal computer systems we're all busy constructing and programming. It performs this service in a way which adds color and excitement to this newest of art forms, the computer application.

Computer Lib/Dream Machines is must reading for the beginner, and is also a refreshing self examination for the old hand at programming and systems work.

You can order your copy of Computer Lib/Dream Machines from BYTE's Books for \$7. Send your order today to BYTE's Books, 70 Main St, Peterborough NH 03458. Help stamp out cybercrud.

\footnotetext{
Please add 50 cents for postage and handling. Please allow six weeks for delivery.
Send copies of Computer Lib/Dream Machines to:


## VIDEO CAMERA KIT

SOLID STATE SALES

## Features Value

SPECIAL PURPOSE TRANSISTORS

## INTEGRATED CIRCUITS

HARDWARE

## P.C. SOCKETS

SILICON POWER
RECTIFIERS

## ZENERS

TRIACS

## SCRS

## LIGHT EMITTING DIODES

SPECIAL PURCHASE ITEMS

A UNIQUE ALL SOLID STATE CAMERA KIT FEATURING A. . . $100 \times 100$ BIT SELF SCANNING CHARGED

COUPLED DEVICE


SCHEMATIC FOR COMPUTER
INTERFACE INTERFACE
SUPPLIED MAY be USED WITH AMATEUR
RADIO FOR VIDEO USED FOA CHARAC-
TER RECOGNITION TER RECOGNITION
FOR COMPUTERS FITH EXTERNAL CIRCUITS
MAY BE USED IN
A VACUUM, UNDER WATER, HIGH ALTITUDE, AND IN MAGNETIC ENVIRONMENT
BECAUSE THERE IS NO HIGH
VOLTAGE OR VOLTAGE OR MAGNETIC

WE SUPPLY ALL
SEMICONDUCTORS BOARDS, DATA SHEETS, DIAGRAMS, RESISTORS
AND CAPACITORS
SORRY, WE DO NOT SUPPLY
THE CASE, BATTERIES IOR SUPPLY) THE LENS
(NOT SUPPLIEDI DEPENDS UPON THE USE



\section*{| $2 N$ |
| :--- |
| $2 N$ |
| $2 N$ |
| R1 |
| R1 |
| En |
| $2 N$ |}

VERAS-8 MICROPROCESSOR KIT
IS ONE OF THE MOST ADVANCED MCU SYSTEMS ON THE MARKET TODAY FOR ONLY $\$ 159.00$

2) Twit






He wiplec:
 .

4K MEMORY KIT $\mathbf{S t 5 9 . 0 0}$
Mannory Board: Our unique mermory card lealures





expander board

ull butierng. Individuat powet lerminals tor aeeth memmary card at
memuty by adding
connactior options.
EXPANDER KIT WITH ONE EDGE CONNECTOR 559.00
To the bust of our knowledge, oursst the onity kit on the markel dessgned lor aisy



TANTULUM CAPACITORS | $.22 U F$ | $35 V$ | $5 / \$ 1.00$ | $6.8 U F$ |
| :--- | :--- | :--- | :--- | 4.7UF 35V $4 / \$ 1.00$

CT 7001 ALARM CLOCK CHIP . . . $\$ 5.75$

## NATION

## MM1403-3.20

MM1404-2.50
MM5013-7.75
MM5016-3.50
MM5SO17-4.75
MM5055-4.00
MM5055-4.00
MM5056-4.00

CIMOS (DIODE CLAMPED)

| $74 \mathrm{CO2-} 2-.30$ | $4016-.50$ | $4035-1.25$ |
| ---: | :--- | :--- |
| $74 \mathrm{CO1-}$ | .30 | $4017-1.10$ |
| $4001-.24$ | $4042-.75$ |  |
| $4002-.24$ | $4019-1.20$ | $4046-2.20$ |
| $4006-1.20$ | $4022-1.00$ | $4047-2.20$ |
| $4069-.50$ |  |  |
| $4007-.30$ | $4023-.25$ | $4050-.50$ |
| $4009-.50$ | $4024-.85$ | $4055-1.95$ |
| $4010-.50$ | $4025-.25$ | $4066-1.10$ |
| $4011-.25$ | $4027-.55$ | $4071-.25$ |
| $4012-.25$ | $4028-.95$ | $4077-.25$ |
| $4013-.35$ | $4029-1.25$ | $4081-.2$ |
| $4015-1.10$ | $4030-.95$ | $4076-1.20$ | LED READOUTS

HP 7740..3"C.C

MAN-4.25" C.C
MAN-7.3"C.A.

Fill Wave Bridges


## BOOK REVIEWS

Microprocessors and Microcomputers by Branko Soucek. Published by John Wiley and Sons, New York.

This over 600 page book contains important information for the experienced hardware and software person but is still very readable for the neophyte. The information presented is both valuable in learning the basics of microprocessors and in selecting a microprocessor or microcomputer.

The first 100 pages cover number systems, digital circuits, and basic microcomputer instructions in enough depth to get the beginner up to speed. In the chapter on basic microcomputer instructions and in much of the material which follows, the DEC PDP-8 minicomputer (or Intersil 6100 microprocessor) instruction set is used as an example. Although this is a somewhat primitive instruction set, it is sophisticated enough to get the concepts across without confusing the beginner with a complicated machine language.

This introduction to basic hardware is followed by about 50 pages of discussion of programming, both in assembly language and in FORTRAN. Throughout this chapter FORTRAN source code is presented with the assembly code which performs the same function.

The last chapter in Part I presents information on microcomputer 10 and interfacing. The necessary hardware and software for both programmed 10 and direct memory access are discussed.

Part II is a detailed discussion (almost 250 pages) of representative microprocessors. The architecture, instruction set, timing, and interfacing are discussed and programming examples are presented for the following microprocessors:

```
4004/4040
8008/8080 and MCOM-8
M6800
PPS-4
8008/8080 and MCOM-8
PPS-4
```

PPS-8
IMP 4/8/16 and PACE
Part III, called "New Microprocessors and Special-Purpose Microsystems," is much like Part II. It presents detailed discussions of microprocessors and hardware which is similar to microprocessors. The subjects include:
PDP-11 Minicomputer and
LSI-11 Microcomputer
F8 Microprocessor
SMS Microcontroller
IM6100 Microprocessor and
PDP-8 Minicomputer

Phil Hughes
PO Box 43
Richland WA 99352

## Amateur Radio Publications

The world of amateur radio has been one of the long time homes for individuals practicing state of the art electronics techniques. In the present state of the art, the combination of microprocessors with amateur radio communications practices leads to some exciting applications such as packet switching communications networks, automated amateur radio stations, etc. (The first prize winner in the World Altair Computer Convention which MITS put on in March was an amateur radio operator.) To provide a pointer into amateur radio circles, the following is a set of short reviews of some amateur radio publications which appear monthly.

## QST

QST was the first amateur radio magazine ever. It is published monthly by the American Radio Relay League, 225 Main St, Newington CT 06111. It is professionally edited and currently is run in a full size magazine format (before January 1976, all the ham magazines were in small magazine format). Each issue typically has a mixture of technical articles, tutorial articles, and a


large amount of information about amateur radio happenings. As an example of technical content, the April 1976 issue featured a cover article entitled "One KW - Solid State Style, Part 1" which shows how H O Granberg, WB2BHX/OH2ZE/7 built a high power radio frequency amplifier using solid state electronics exclusively. Other April articles included the fourth part of a series on "Learning to Work With Integrated Circuits," and articles on several technical aspects of amateur radio. A major portion of the magazine is devoted to general interest items. The April issue also reported amateurs' emergency work in the Guatemalan earthquakes, and operating activities such as message networks, civil defense organizations, etc. QST is also chock full of advertising about amateur radio products. This journal is a "must" for radio amateurs, and can be obtained for $\$ 9$ per annum.

## Ham Radio

Ham Radio is another excellent amateur radio publication, which has been published since the late 1960s. Its primary thrust is technical, and this is evident in the list of articles found (for example) in the March 1976 issue:

> Crystal Controlled Oscillators DT-500 RTTY Demodulator WWVB Signal Processor High Speed Divide by $N$ Counters Off the Air Transmitter Tuneup

VHF/UHF Receivers - How to Improve Them
5/8 Wavelength Vertical Antenna for Two Meter FM
Microprocessors - 8080 Output
Instructions
High Performance Bench Power Supply
The non-technical editorial content is a good approximation to $0.0 \%$; the quality of the technical articles is on a par with many engineering publications of the professional world. As is the case with QST, Ham Radio has many interesting and informative advertisements for amateur radio products.

You can subscribe for $\$ 10$ per annum by writing to Ham Radio, Greenville NH 03048.

## CQ

Another amateur radio publication is $C Q$, named after the greeting code used by amateur radio operators when looking for a contact. It is also a monthly publication, and its February 1976 issue was filled with technical and general interest information for amateur radio operators. On the technical side, this issue had an extended commentary on slow scan TV methods, short columns on technical pointers, an article on hardware modification of a commercial ham transceiver, etc. $C Q$ also carries many advertisements for amateur radio products.

You can subscribe to CQ for one year @ \$7.50. Write CQ, 14 Vanderventer Av, Port Washington NY 11050.■

MC14412 UNIVERSAL MODEM CHIP
MC14412 contains a complete FSK modulator and de-modulator compatible with foreign and USA communications. ( $0-600 \mathrm{BPS}$ )
FEATURES:
. On chip crystal oscillator

- Echo suppressor disable tone gerierator
- Originate and answer modes
. Simplex, half-duplex, and full duplex operation
. On chip sine wove
- Modem self test mode
. Selectable data rates: $0-200$
$0-300$
0-600
. Single supply
VDD $=4.75$ to $15 V D C-F L$ suffix $V D D=4.75$ to 6 VDC - VL suffix
TYPICAL APPLICATIONS:
.Stand alone - low speed modems
.Built - in low speed modems
- Remote terminals, occoustic couplers

MC14412FL. . . . . . . . . . . . . . . . . . . . . . . . . . . . . . . . . . . . $\$ 28.99$
MCI4412VL. $\$ 21.74$
6 poges of data.. above. ...... $\$ 4.95$ .60
Crystal for the above.......\$4.95


## MCI4411 BIT RATE GENERATOR.

Single chip for generating selectable frequencies for equip mánt in dota communications such as TTY, printers, CRTs or microprocessors. Generates 14 different stondard bit rates which are multiplied under external control to $1 X$, $8 \mathrm{X}, 16 \mathrm{X}$ or 64 X initial value. Operates from single +5 volt supply. MC14411................................ \$11.98 4 pages of data. . . 40 Crystal for the above.
$4 \times 4 \times 2$ CROSSPOINT SWITCH. MC3416L is a poi of $4 \times 4$ Matrices of isulated SCR's triggered by a common selection matrix. Use for low loss analog switching such as relephone, data selection, etc.

9 Pages of specs for 3416 . . . . . . . . . . . . . . . . . . . . . . . $\$ 1.00$

4 DIGIT COUNTER. MM74C926 is a 4 digit counter with 7 segment output. Carry output for coscading and intermal display select allows outputting of counter or set of internal latches. 3 to 6 V operation. Great for clocks, event ond frequency counters.
MM74C926 - with spec sheet.
.
FOUR QUADRANT MULTIPLIER. MC1495L provides output as a linear product of two analog input. Use for frequency doubler, balanced modutar/demodulator electronics gains control.
MC1495 L.
55.50

6 poges of specs
9 pages of opplications. . .90

## TELETYPE CODE CONVERSION CHIP

MM5220BL converts 5 level Boudat into 8 level ASCII. Use this chip to make your ald TTY talk to your new computer. MM5220BL. . . .......................................... $\$ 18.00$


## MOS TIME BASE KIT

Only 1 " $\times 1.5^{\prime \prime}$. Input 5 to 15 VDC , output is 60 HZ square wave for partable or mobile clocks. PC board is


$8 \mathrm{HO5}$ Volroge regulator, Fairchild $5 \mathrm{~V}, 5 \mathrm{~A}$ TO-3 regulotor. Take care of those heovy current requirments without separate regulatar/pass transistor combinations. Use it with the same ease of instalation as the 309 K (same pin arrangement.). ............ with specs.................... . $\$ 11.25$


LM3I7 Voltage Regulator. 1.5A, 3 terminal adjustable regulator in $\mathrm{TO}-3$ case. Adjusts from +1.2 V to +37 V . Complete overload protection. . $1 \%$ load regulation, $.01 \% / \mathrm{V}$ line regulation. No need to stock assorted regulators - just stock resistors. . . . . . . . . . . . . . . . . . . . . . . 54.99 Specs for the obove..

BAIA BOOKS BY NATIONAL SEMICONOUCTOR INEAR Cose $\$ 3.95$ LINEAR APPLICATIONS, Dozens of opplication notes and rechnical briafs covering the use of op-amps, regulators, phase locked loops and oudio amps................... $\$ 3.25$ CMOS Gates, Flip Flops, registers, functional blocks \$3 VOLTAGE REGULATORS. A must for anyone making a power supply. Complete theory including transformers, filters, heat sinks, regulators, etc. ................... \$ MEMORY, Infarmation on MOS and Bipalar memories' RAMS, ROMS, PROMS and decoders/encoders. . ... \$3.95 NTERFACE: Covers peripheral drivers, level translators, ine driver/receivers, memory and clock drivers, sense amps display driver and opto-couplers............
(Outside U.S., add postoge for 1.51 )

DATA BOOKS FROM FAIRCHILD.
UA Linear. 776 pages of data and applications for fairchild linear IC. Great value......................... $\$ 4.25$ MOS/CMOS/nMOS/PMOS/CCD. Data and applications on MOS and charge coupled devices including preliminary data on new and future offerings. Wont to know about 16 K charge coupled line oddressoble mamories?.......... \$3.95

# ADPRECDII Says 

This Year, Tri-Tek is Celebroting too!
It's our 5th Anniversary of supplying you,
our good customers, with quality parts and products.

## RРл

## GOLD CHIP

## Linear Integrated Circuits

Brand new process by RCA in which the aluminum metalization has been replaced by gold. The chip is then hermetically seal. ed. What this means to you is unprecedented reliability and uniformity. Plastic parts thar meet mil specs!!
Tri-Tek is proud to be the first to bring this new level of performance to you of SURPLUS PRICES. Why buy regrades??
CA301A. Improved, general purpose op-amp, 8 pin dip. . 59 C CA307... Super 741 op-amp. 8 pin dip................... 52 C CA324... Compensated quad op-amp, 14 pin dip...... 51.80 CA339A.. Low offset quad comparator. 14 pin dip... $\$ 1.59$ CA741C. . Fomous general purpose op-amp, 8 pin dip.. 45 c CA747C.. General purpose dual op-omp, 14 pin dip... 82c CA748C. Externally compensated 741, 8 pin dip........49 C CA1458.. General purpose dual op-amp. 8 pin dip..... 69 C CA3401. .Quad single supply ( $5-18 \mathrm{~V}$ ) op amp. 14 pin.. 89 c

Another super buy from RCA. CA555 timer. 8 pin dip. 59 C

INTEL Data Catalog. Contains latest information on all the famous INTEL micro-processor and memories. . . . . . . . . . . . 54.00

MPROVED Performance version of the famous 8080, 8 bit
nicro processor, 8080A.................................... . . . $\$ 34.95$

peripherols such as teletypes \& video terminols. ... 59.45

DC to DC CONVERTER MODULE. Tiny potted module is a complete regulated output up-converter. $4.5 V \mathrm{D} . \mathrm{C}$. input provides approximotely 14 V @ 12 mA . Will operote down to $2.5 \vee$ with reduced outpur. Size only $I^{\prime \prime} \times \frac{1}{2}{ }^{\prime \prime} \times \frac{1}{1}$. With connection diagram. . . . . . . . . . . . . . . . . . . . . . . . . 51.25

NSL4944 LED. Current regulated, universal diffused-lens red LED lamp. A GaAsP solid-state high intensity LED encopsulated in a plastic packoge containing o current regulating IC that provides constont intensity over a wide voltage range. 2 to $18 \mathrm{~V}, \mathrm{AC}$ or $D C$. Use for indicator lamps, optical coupling, battery charging circuits, logic probes, almost any place you need a lamp. Long life, wide angle. No series resistor needed. Typical 13mA forward current. NSL4944. .......with ponel mounting clip............. 89 c

D-A CONVERTER BY ZELTEX
8 bit precision hybrid circuit for use in controllers, timers, volt meters, etc. Molded plastic pockage with P.C. pins. Super buy on this better than usual subsystem. ZELTEX model 2D430. DAC-430............................ 54.95

NEW BOOK1!! "An Introduction to Microcomputers" This is the book which Fairchild Somiconductor Company called ". .....the best darned introduction to the industry to date." Covers overything from basic concepts to a review of real microcomputers. $1 \mathrm{MC}-001 \ldots . . . . . . .$.
tRI-tek, Inc.
6522 north 43RO avenue.
Glenoale, arizona 85301
phone 602-931-6949

## ALARM CLOCK KIT SIX DIGIT LED

Thousands of hobbyists have bought and built our original clock kit and were completely satisfied. But we have received many requests for an alarm clock kit with the same value and quality that you have come to expect from S.D. So, here it is!

## THE KIT INCLUDES:

1 Mostek 50252 Alarm Clock Chip
6 Hewlett Packard . 30 in. common cathode readouts.
15 NPN Driver Transistors
2 Switches for time set
2 Slide Switches for alarm set and enable
1 Filter Cap
4 IN4002 Rectifiers
1 IN914 Diode
1 . 01 Disc Cap
15 Resistors
1 Speaker for alarm

PCB -3.00
1 LED lamp for PM indicator.
60 Hz . Crystal Time Base FOR DIGITAL CLOCKS S. D. SALES EXCLUSIVE!

2102 1K RAM's - 8 FOR $\$ 12.95$
New units We bought a load on a super deal, hence this fantastic price.

Units tested for 500NS Speed.
MOTOROLA RTL IC'S
Brand new, factory prime. Hard to find, but still used in a variety of projects. (See the RTL Cookbook by Howard W. Sams.)

| MC724P-59c | MC780P-89c | MC791P-69c |
| :---: | :---: | :---: |
| MC725P-59c | MC785P-49c | MC792P-59c |
| MC764P-49c | MC787P-89c | MC799P-59c |
| MC767P-69c | MC788P-49c | MC9704P-89c |
| MC771P-49c | MC789P-59c | MC9709P-69c |
| MC775P-89c | MC790P-89c | MC9760P-68c |
| MV-50 TYPE LED's by LITRONIX | 3 DIGIT LED ARRAY - 75c <br> by LITRONIX <br> DL33MMB. 3 MAN-3 Size Readouts in one package. Thase are factory prime, not retested rejects as sold by others. compare this price! 75 c 3 for $\$ 2$. |  |
| 10 for $\$ 1$ <br> Factory Prime! |  |  |

SALE ON CUT LEAD SEMICONDUCTORS
Leads were cut for PCB insertion. Still very useable.
1N914/1N4148
1N4002 1 Amp 100 PIV 1N4745A 16V 1 W Zener EN2222 NPN Transistor


All in one 28 PIN DIP. 4 Decade counters, latches, MUX circuits, display decoders, etc. Features: 5VDC operation, 25 MW power consumption, BOTH 7 segment and BCD outputs. Perfect for DVM's, frequency meters, tach's, etc. Can be cascaded for more digits. \#5002 $\$ 8.95$.

## 8008-1 MICRO PROCESSOR

New Units. High speed 8008. Almost twice as fast as units sold by others. Still a very versatile and widely used MPU. No data book, only pinout data included at this price, $\$ 12.50$ LIMITED QTY

1101 A RAM IC
$256 \times 1$ BITS. Perfect for 8008 or small systems such as telephone dials that do not require a lot of memory. Special 59c. 8 FOR $\$ 4$.

C \& K MINI TOGGLE SWITCH
\#7103 SUB MINI SPDT Center OFF. Special - 99c
TTL INTEGRATED CIRCUITS

| 7400 | - 19c | 7420 | 19c | 7483 | 95c | 74154 | 1.00 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 7402 | - 190 | 7432 | 34c | 7485 | 95c | 74157 | 75c |
| 7404 | - 19c | 7437 | - 39c | 7486 | - 45c | 74161 | 95c |
| 74S04 | 29c | 7438 | 39c | 7490 | 65c | 74164 | - 1.10 |
| 7406 | 29c | 7447 | 85c | 7492 | 75c | 74165 | 1.10 |
| 7408 | 190 | 7448 | 85c | 7495 | 75c | 74174 | 95c |
| 74504 | 440 | 7473 | 39c | 74121 | 38c | 74181 | - 2.50 |
| 7410 | 19c | 7474 | 35c | 74123 | 65c | 74191 | - 1.25 |
| 7411 | 29c | 7475 | 69c | 74141 | 75c | 74192 | - 1.25 |
| 7413 | 50c | 7476 | 35c | 74153 | 75 c | 74193 | - 1.00 |

S. D. SALES CO.
P.O. BOX 28810

DALLAS, TEXAS 75228
$T-\infty-\infty$
$-\infty-\infty$

## ORDERS OVER \$15 CHOOSE \$1 FREE MERCHANDISE

## S. D. SALES CO.

## UP YOUR COMPUTER!

## 21L02-1 1K LOW POWER

500NS STATIC RAM

## TIME IS OF THE ESSENCE

And so is power. Not only are our RAM's faster than a speeding bullet but they are now very low power. We are pleased to offer prime, new 21LO2-1 low power and super fast RAM's. Allows you to STRETCH your power supply farther and at the same time keep the WAIT light off.

## TOUCHTONE KEYBOARD SWITCH SET

By Controls Research. High quality long life switches with keytops. For encoders, combination locks, etc.


12 switches and tops, including 0 thru 9. Switch contacts are independent, allows hook-up to any matrix. Keytops easily removed.

## 8 for $\$ 17.50$

## \$2.95 Set 2 for \$5.00

## 4K LOW POWER RAM BOARD KIT

Imsai and Altair 8080 plug in compatible. Uses low power static 21L02-1 500 ns. RAM's. Fully buffered, drastically reduced power consumption, on board regulated, all sockets and parts included. Premium quality plated thru PC Board.

THE WHOLE WORKS
\$89.95


Call your BANK AMERICARD or MASTER CHARGE order in on our CONTINENTAL UNITED STATES TOLL FREE WATTS:


## 1-800-527-3460

TEXAS RESIDENTS CALL COLLECT:

## 214/271-0022

Please call between 8:30 AM and 6:00 PM C.S.T. - Monday through Friday. You may also call to check stock or just ask a question. However, only B.A.C. and M.C. orders will be accepted. We do not ship C.O.D. (See terms of sale on other page.)
at a later date permitting later versions of the same board to have 32 K bytes of memory. Further, the firm in question stated that it would introduce a new twist to the marketplace by selling this product exclusively through retail computer stores in much the same way as high fidelity equipment is sold. The selection of power supply, keyboard, cassette recorder and TV monitor vendors is left to the purchaser as aided by the personal consultation of the retailer, just as selection of components of a high fidelity stereo system is often aided by an interactive session with a retailer.

- Item: As if to confirm the trend heralded by the phone call from "brand A," "brand B" calledrone week later to mention its version of the completely packaged and ready to go computer. The "brand $B$ " product is similar to "brand $A$ " in its confirmation of the trend to "no hassle" computing, but the nature of the system is a bit different. The "brand B" computer is a modular product
with a traditional backplane design, 6800 processor and packaging assembly with power supply. The "brand B" computer looks inherently like a traditional minicomputer system available at very low prices which include (so it is said) a BASIC package. This "brand B" computer also comes in one form: assembled and completely burned in component cards which are mixed and matched to form a completed system. Again note the emphasis on "no hassle" hardware in the form of assembled and tested boards, with systems software bundled into the price, also under $\$ 1000$ for a minimal system ( 8 K bytes) exclusive of RS-232 or 20 mA current loop 10 terminal.

Both the "brand A" and "brand B" computers are not yet widely marketed, but the trend is clear: the products in the personal computing market place are becoming more refined and oriented to user satisfaction, yet still very reasonable in price. The "little IBMs" of the micro computer world will become the corporate giants of tomorrow by perceiving this trend and wrapping up a package of "customer satisfaction" and convenience.



## POWER SUPPLY MODULE

New, plug-in module. Plugs into $A C$ outlet provides 12 volts $A C$ at $1 / 2$ amp by two screw terminals. Great for various clocks, chargers, adding machines, etc. New
\$2.50 ea. 5/\$11.00

## POWER SUPPLY

LAMBDA 5VDC 74 AMP
LV-EE-5-OV
$\$ 125.00$

NJE 5/OUP-D5
5 VDC 32 AMP
$\$ 75.00$

## CLOCK KIT \$14.00

Includes all parts with MM5316 chip, etched \& drilled PC board, transformer, everything except case.

SP-284 \$14.00 each $2 / \$ 25.00$


ASCII KEYBOARD, brand new w/TI ASCII chip inplace \& data $\$ 45.00$
COMPUTER GRADE LOGIC SUPPLY CAPS, BRAND NEW

| 47,000 Uf | 25V | $\$ 2.00$ | ST | 1,000 | 50 | .90 | AL |
| :--- | :--- | ---: | :--- | :--- | :--- | ---: | :--- |
| 32,000 | 25 | 1.75 | ST | 3,300 | 35 | 1.25 | AL |
| 160,000 | 10 | 2.00 | ST | 1,600 | 20 | .60 | AL |
| 66,000 | 10 | 2.00 | ST | 8,000 | 16 | 1.25 | AL |
| 1,000 | 60 | .90 | AL | 500 | 6 | .35 | AL |
| 2,000 | 55 | 1.00 | AL | "ST" screw top ...."AL" axial |  |  |  |

5 VOLT 1 AMP REGULATED power supply kit for logic work. All parts including LM-309K
$\$ 7.50$
DUMMY LOAD RESISTOR, non inductive, 50 ohm 5 watt $\$ 1.00$
"AA" NICAD CELLS brand new, fine biz for handy talkies
\$1.25 ea. 9/\$9.00

LINEAR by RCA, brand new, gold bond process

| 301 | $\$ .60$ | 747 | $\$ .82$ | MM5314 | $\$ 3.00$ |
| :--- | ---: | :--- | ---: | :--- | ---: |
| 307 | .52 | 748 | .50 | MM5316 | 3.00 |
| 324 | 1.80 | 1458 | .96 | 7001 | 8.00 |
| 339A | 1.60 | 3401 | .80 |  |  |
| 741 | .50 | 555 timer | .60 |  |  |

MEMORY SYSTEM \$125.00
New memory system by Honeywell, small ... measures only $9 \times 4 \times 1$ inches. 1024 core memory, 1024 words with $8,9,10$ bits/word. Random access, with all logic, register, timing, control, core select and sense functions in one package. New, booklet of schematics and data. Looks like a good beginning for a mini-computer. Limited supply on hand.
Ship wgt 3 lbs. \#SP-79
$\$ 125.00$


## CORE MEMORY

Another brand new memory, ultra small. Measures only $4 \times 4$ inches with format on one plane of $32 \times 32 \times 16(16,384)$. Only about 35 units of this on hand.
\#SP-81
$\$ 20.00$

FREE CATALOG SP-8 NOW READY


Please add shipping cost on above.

PHONE 617-595-2275
MESHNA PO Bx 62 E. Lynn Mass. 01904
culator by either actually doing arithmetic (what?) or using an earlier version of the mechanical hand held adding machine. However, Ed forgot about the difference between base 16 and base 10 , and the fact that readers without an exposure to traditional magazines for data processing (ie: Iarge scale computers) probably had not heard of the "Hexadat." Since no one has come out with a way to manufacture and market several tens of millions of base 16 four function calculators at $\$ 8.95$, we hexadecimal lovers are prevented from economically using a hand held base 16 electronic calculator. But in lieu of electronics, the Radix Precision "Hexadat" will serve excellently at times when you can't rely on your computer for accurate calculation.

You can acquire a Hexadat for $\$ 35.95$ complete with the leather case and instruction manual. In case you are prejudiced against hexadecimal, the same company can help you out with the "Octadat" octal calculator at $\$ 14.95$. Either way, contact Radix Precision at PO Box 13861, Atlanta GA 30324.■

## April BOMB Results

Winner of the $\$ 50$ for most appreciated article in the April 1976 issue of BYTE is Theodor Nelson's "The Magic of Computer Languages." Runners up in the voting were Robert Wier and James Brown, who wrote "Design an On Line Debugger," and Don Lancaster's "How To Build a Memory With One Layer Printed Circuits."■

## BDMB: bYTE's Ongoing Monitor Box


#### Abstract

BYTE would like to know how readers evaluate the efforts of the authors whose blood, sweat, twisted typewriter keys, smoking ICs and esoteric software abstractions are reflected in these pages. BYTE will pay a $\$ 50$ bonus to the author who receives the most points in this survey each month.


PAGE
NUMBER
ARTICLE
6 Jones: Coincident Current Ferrite Core Memories
18 Anderson: Assembling a Sphere
22 Howerton: Explore an 8080 with Educator-8080
30 Wadsworth: Machine Language Programming for the "8008"-Chapter 1
40 Baker: Put the "Do Everything" Chip into Your Next Design
46 Suding: Why Wait?
54 Hayes: Surplus Electronics in Tokyo and Manila
58 Hogenson: Make Your Own Printed Circuits
64 Lerseth: A Plot is Incomplete Without Characters

## LIKED

LEAS
BEST
012345678910
012345678910
012345678910

012345678910

012345678910
012345678910
012345678910
012345678910
012345678910

Call Your
IMS Dealer
Today.
CALIFORNIA: Berkeley, Applied
Computer Technology, (415) 527-6760;
Campbell, Byte Shop No. 3, (408) 377-4685;
Chico, Micro-Byte, (916) 345-9396;
Cupertino, Computer Media, Inc.,
(408) 294-7004; Fremont, Casual Bedding,
(415) 796-4555; Fullerton, Bits $n$ Bytes,
(714) 525-9613; Morgan Hill, Metatic

Corporation, (408) 779-8150; Mountain
View, Byte Shop No. 1, (415) 969-5464;
Orange, Computer Mart of Los Angeles,
(714) 633-1222; Oxnard, Pete's Electronics,
(805) 485-6467; Rancho Palos Verdes,

Microprocessor Marketing, (213) 541-6344;
San Diego, The Computer Center,
(714) 292-1657; San Gabriel, Sunny Sounds,
(213) 287-1811; San Leandro, Robert H.

Edmonds, (415) 357-5837; San Lorenzo,
The Computer Den, (415) 278-4720;
San Luis Obispo, Proko Electronics,
(805) 544-5441; Santa Clara, Byte Shop No. 2, (408) 249-4221; Van Nuys, Computer Components, Inc., (213) 782-7924;
COLORADO: Boulder, Inter-Mountain
Digital, (303) 442-3960; CONNECTICUT:
Ellington, Heuristic Systems, (203) 871-1181;
FLORIDA: Jacksonville, Douglas
Computer Systems, (904) 724-8726;
Pinellas Park, Elecon Corporation,
(813) 541-3021; INDIANA: Bloomington,

Data Domain, (812) 337-1671; KENTUCKY:
Louisville, Cybertronics, Inc., (502) 459-0426,
LOUISIANA: Baton Rouge, Executone
Microcomputer Div., (504) 383-1371;
MASSACHUSETTS: Waltham, Computer
Mart, Inc., (617) 890-0677; NEW JERSEY:
Colonia, Computer Mart of New Jersey,
(201) 574-2173; Edison, William Electronic

Supply, (201) 985-3700; Hackensack, The
Computer Room, (201) 342-6667; Hoboken,
Hoboken Computer Works, (201) 420-1644;
Trenton, Technical Design Labs, Inc.,
(609) 599-1868; NEW YORK: Hollis,

Synchro Sound Studios, Inc., (212) 359-1489;
New York, Audio Design Electronics,
(212) 226-2038; Computer Mart of New York, (212) 787-4051; Oneonta, Collegiate
Audio, (607) 432-1930; Rye, The Computer
Corner, (914) 967-7853; OHIO: East
Cleveland, E.L.S. Systems Engineering,
(216) 321-8303; OREGON: Eugene, The

Real Oregon Computer Company, Inc.,
(503) 484-1040; PENNSYLVANIA: Frazer,

Personal Computer Corp., (215) 647-8460;
Philadelphia, J. B. Industries, (215) 224-9926;
Pittsburgh, Martin J. O'Boyle \& Assoc.,
(412) 361-1602; RHODE ISLAND: West

Warwick, Computer Power, Inc., (401)
822-0204; TEXAS: Dallas, Digitex,
(214) 744-1849; Houston, Electronic

Specialty Co., (713) 475-4402; Polaris
Computer Systems, (713) 527-0348;
Richardson, The Micro Store, (214) 690-7113;
UTAH: Orem, Computers \& Stuff, (801)
224-2066; VIRGINIA: Reston, Media
Reactions, Inc., (703) 471-9330; Richmond,
Computer Hobbies Unlimited, (804)
276-5056; WASHINGTON: Lacey,
Microcomputer Applications, (206) 456-3924;
Seattle, The Retail Computer Store,
(206) 524-4101; WISCONSIN: Marinette,

Lauerman Department Stores, (715)
735-3311; Milwaukee, Milwaukee Computer
Store, (414) 445-2403; CANADA:
Manitoba, Brandon, Canadian
Microcomputer, (204) 725-1600.
Join Our Growing Network of
World-Wide Independent Dealers.


IMS Associates, Inc.

# If you thought a rugged, professional yet affordable computer didn't exist, 

## think IMSAI 8080.

Sure there are other commercial, high-quality computers that can perform like the 8080. But their prices are 5 times as high. There is a rugged, reliable, industrial computer, with high commercial-type performance. And prices that are competitive with Altair's hobbyist kit. The IMSAI 8080. Fully assembled, it's \$931. Unassembled, it's $\$ 599$. And ours is available now. In our case, you can tell a computer by its cabinet. The IMSAI 8080 is made for commercial users. And it looks it. Inside and out! The cabinet is attractive, heavy-gauge aluminum. The heavy-duty lucite front panel has an extra 8 program controlled LED's. It plugs directly into the Mother Board without a wire harness. And rugged commercial grade paddle switches that are backed up by reliable debouncing circuits. But higher aesthetics on the outside is only the beginning. The guts of the IMSAI 8080 is where its true beauty lies.

The 8080 is optionally expandable to a substantial system with 22 card slots in a single printed circuit board. And the durable card cage is made of commercial-grade anodized aluminum. The Altair kit only provides 16 slots maximum in four separate sections, each section

requiring 200 solder connections.
The IMSAI 8080 power supply produces a true 28 amp current, enough to power a full system. The Altair produces only 8 amps.

You can expand to a powerful system with 64 K of memory, plus a floppy disk controller, with its own on board 8080 - and a DOS. An audio tape cassette input device, a printer, plus a video terminal and a teleprinter. These peripherals will function with an 8 -level priority interrupt system. IMSAI BASIC software is available in $4 \mathrm{~K}, 8 \mathrm{~K}$

and 12 K , that you can get in PROM. And a new \$1394K RAM board with software memory protect.

Find out more about the computer you thought didn't exist. Get a complete illustrated brochure describing the IMSAI 8080, options, peripherals, software, prices and specifications. Send one dollar to cover handling to IMS. The IMSAI 8080. From the same technology that developed the HYPERCUBE Computer architecture and Intelligent Disk systems.

Dealer inquiries invited.


IMS Associates, Inc. 1922 Republic Avenue San Leandro, CA 94577 (415) 483-2093

## The small wonder



## of the micro－World

Measuring just $11^{\prime \prime}$ wide $\times 11^{\prime \prime}$ deep $\times 5^{\prime \prime}$ high，and weighing a mere 7 pounds，the Altair ${ }^{\text {M }} 680 \mathrm{~b}$ is a complete，general－purpose computer．

The secret to this revolutionary，small computer is its CPU board．This double－sided board fits along the bottom of the Altair case and plugs directly into the front panel board．It contains the new 6800 microprocessor， 1,024 bytes of RAM memory，a 256 byte PROM monitor，provisions for 768 bytes of additional PROM or ROM，and a single Interface port with a Motorola ACIA serial interface adapter which can be configured either RS－232 or TTY． A five level Baudot interface option is also available．

The Altair 680b can be programmed from front panel switches， or it can be interfaced to a video display terminal，or teletype－ writer．Three additional circuit boards can be plugged inside the Altair 680b for further memory and interface expansion．The first of these boards now under development is an 8 K RAM memory board．

Software already developed includes a resident two pass assembler and 8K BASIC．The Altair 680b is also compatible with Motorola 6800 software．

The Altair 680 b is ideal for hobbyists who want a powerful computer system at an economic price．Altair 680b owners qualify for membership in the Altair Users Group，and like other Altair owners，they receive a complimentary subscription to Computer Notes and complete factory support．

## PRICES：

Altair 680b kit with complete，easy－to－understand assembly man－ ual，operator＇s manual，and programming manual． $\$ 466$
Assembled Altair 680b $\$ 625$
Baudot option．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．．． 42
8 K RAM memory board and software prices to be announced soon．

## MAIL THIS COUPON TODAY



Price，specifications subject to change．Please allow up to 60 days for delivery．


[^0]:    Articles Policy
    BYTE is continually seeking quality manuscripts written by individuals who are applying personal systems, or who have knowledge which will prove useful to our readers. Manuscripts should have double spaced typewritten texts with wide margins. Numbexing sequences should be maintained separately for figures, tables, photos and listings. Figures and tables should be provided on separate sheets of paper. Photos of technical subjects should be taken with uniform lighting, sharp focus and should be supplied in the form of clear glossy black and white prints (if you do not have access to quality photography, items to be photographed can be shipped to us in many cases). Computer listings should be supplied using the darkest ribbons possible on new (not recycled) blank white computer forms or bond paper. Where possible, we would like authors to supply a short statement about their background and experience.

    Articles which are accepted are typically acknowledged with a binder check 4 to 8 weeks after receipt. Honorariums for articles are based upon the technical quality and suitability for BYTE's readership and are typically $\$ 15$ to $\$ 30$ per typeset magazine page. We recommend that authors record their name and address information redundantly on materials submitted, and that a return envelope with postage be supplied in the event the article is not accepted.

[^1]:    Chapters 2 and 3 of MACHINE LANGUAGE PROGRAMMING FOR THE " 8008 " (and similar microcomputers) will appear in BYTE's August and September issues, respectively. ${ }^{\text {- }}$

[^2]:    W Douglas Maurer University Library Room 634 George Washington University Washington DC 20052

[^3]:    "The Magic of Computer Languages" [April 1976 BYTE, page 24] was a good article, but 1 wish people would stop writing about "compiling languages" and "interpretive languages." Any language can be either compiled or interpreted. /Quite true.] Small machines tend to use interpreters and large machines tend to use compilers. No doubt 4096 people have already pointed out that APL programs are executed from right to left, so 1 won't mention that. [?] SNOBOL has been simplified already; have you never heard of ICEBOL? [No.]
    "Frankenstein Emulation" [April 1976 BYTE, page 50] was interesting. I suspect that the Main Processor only appears to be faster than the Top Processor because it does

[^4]:    Send for our latest free catalog. Minimum order $\$ 5$, phone orders welcome. Include sufficient postage (2 ibs min.), excess will

